Access Control
Security Begins on the Host

- Even without a network, hosts must enforce the CIA trilogy
- Something on the host — the operating system aided by the hardware — must provide those guarantees
Access Control

- Hardware
- Software
  - Operating systems
  - Databases
  - Other multi-access programs
- Distributed
Hardware

- What is the *minimum* necessary?
- What do other mechanisms buy us?
Minimum Functionality

- Protect the OS from applications
- Protect applications from each other
- Change state from application to OS
- Timer interrupt
Why a Timer?

- Availability is a security feature
- Must prevent uncooperative applications from hogging CPU
- Not going to discuss this more here, but it's a major topic in W4118 (Operating Systems)
**Historical Mechanisms**

- Single privileged mode bit — restrict ability to execute certain instructions
- Memory protection
- Interrupts — hardware and software — cause state transition
What Are Privileged Instructions?

- Ability to do I/O without the OS’s intervention — allowing that could bypass file permission checking
- Ability to manipulate timers
- Ability to access other programs’ memory without OS intervention
Example: IBM System/360 Mainframe

- 40 year old design
- Much of the architecture still in use...
- 4-bit protection key associated with each 4K block of memory, plus read-protect bit
- Single “supervisor mode” bit
- 4-bit state key of 0 can write to anything
- But — operating systems of that time didn’t use the hardware to its full capabilities
Memory-Mapped Control

- On some machines, privileged operations work by memory access
- If applications have no access to such memory, they can’t do sensitive things
- But — must have way to enter privileged state
Multics

- Virtual memory
- “Ring” structure — 8 different privilege levels (x86 has rings, too)
- OS could use rings 0-3; applications could use 4-7.
- (Original design had 64 rings!)
- Special form of subroutine call to cross rings
- Most of the OS didn’t run in Ring 0
What is the Advantage of Rings?
What is the Advantage of Rings?

- A single bit is theoretically sufficient
- Assurance!
- Don’t need to trust all parts of the system equally
- “Principle of Least Privilege”
Assurance

- How do you know something is secure
- Much harder to provide later than features
- A trustable secure system has to be designed that way from the beginning: designed, document, coded, and tested — and maybe proved
Underlying Principles of Privilege

- Two basic approaches to privilege: identity and attribute
- Hardware protection is *attribute*: the state of various registers controls what can and cannot be done
- Easier to manage in a single system
What is the role of the OS?

- Protect itself
- Separate different applications
- More?
Operating Systems and Hardware

- The hardware provides the minimum functionality
- The OS has to provide its own services on top of that
- Must manage access to I/O devices as well
What Protections do Operating Systems Provide?

- User authentication (why?)
- File protection
- Process protection
- Resource scheduling (CPU, RAM, disk space, etc)
User Authentication

- (Much more on this later)
- Why authenticate users?
- Most operating system privileges are granted by identity, not attributes

Procedure:
- Authenticate user
- Grant access based on userid
File Permissions

- Besides user authentication, the most visible aspect of OS security
- Read protection — provide confidentiality
- Write protection — provide integrity protection
- Other permissions as well
**Classical Unix File Permissions**

- All files have “owners”
- All files belong to a “group”
- Users, when logged in, have one userid and several groupids.
- 3 sets of 3 bits: read, write, execute, for user, group, other
- (512 possible settings. Do they all make sense?)
- Written `rwrxrwrx`
- `111 101 001`: User has read/write/exec; group has read/exec; other has exec-only
- Some counter-intuitive settings are very useful
Permission-Checking Algorithm

if curr_user.uid == file.uid
    check_owner_permissions();
else if curr_user.gid == file.gid
    check_group_permissions();
else
    check_other_permissions();
fi

Note the else clauses — if you own a file, “group” and “other” permissions aren’t checked
Execute Permission

- Why is it separate from “read”?  
- To permit only execution  
- Cannot copy the file  
- Readable only by the OS, for specific purposes
Directory Permissions

- “write”: create a file in the directory
- “read”: list the directory
- “execute”: trace a path through a directory
Example: Owner Permissions

$ id
uid=54047(smb) gid=54047(smb) groups=0(wheel),3(sys),54047(smb)

$ ls -l not_me
----r--r-- 1 smb wheel 29 Sep 12 01:35 not_me

$ cat not_me

cat: not_me: Permission denied

I own the file but don’t have read permission on it
Example: Directory Permissions

$ ls -ld oddball
dr--r--r-- 2 smb wheel 512 Sep 12 01:36 oddball
$ ls oddball
cannot_get_at
$ ls -l oddball
ls: cannot_get_at: Permission denied
$ cat oddball/cannot_get_at
cat: oddball/cannot_get_at: Permission denied

I can read the directory, but not trace a path through it to oddball/cannot_get_at
Deleting Files

• What permissions are needed to delete files?
• On Unix, you need write permission on the parent directory
• You can delete files that you can’t write. You can also write to files that you can neither create nor delete
• Other systems make this choice differently
Historical Note

• Unix has *never* been fond of asking “do you really mean that?”

• That said, at least as long ago as February 1973 the original Bell Labs Unix `rm` command prompted if you tried to delete a file you couldn’t write

• In other words, the Unix model is philosophically correct but perhaps incorrect from a human factors perspective
Access Control Lists

- 9-bit model not always flexible enough
- Many systems (Multics, Windows XP, Solaris, some Linux) have more general *Access Control Lists*
- ACLs are explicit lists of permissions for different parties
- Wildcards are often used
Sample ACL

smb.*   rwx
4187-ta.*  rwx
*.faculty  rx
**        x

Users “smb” and ‘4187-ta” have read/write/execute permission. Anyone in group “faculty” can read or execute the file. Others can only execute it.
Order is Significant

With this ACL:

*.*.faculty     rx
smb.*           rwx
4187-ta.*       rwx
*.*             x

I would not have write access to the file
Setting File Permissions

- Where do initial file permissions come from?
- Who can change file permissions?
Unix Initial File Permissions

- Unix uses “umask” — a set of bits to *turn off* when a program creates a file
- Example: if umask is 022 and a program tries to create a file with permissions 0666 (rw for user, group, and other), the actual permissions will be 0644.
- Default system umask setting has a great effect on system file security
- Set your own value in startup script; value inherited by child processes
Why Umask?

- Suppose files were always created with rw,r,r permissions
- What’s wrong with the application simply changing the file permissions after creating the file?
- Race conditions
Multics Initial File Permissions

- Directories contain “initial access control list” — values set by default for new files

- Common setting:
  
  smb.faculty rw
  *
  sysdaemon r
  *.* -

- If group “sysdaemon” doesn’t have read permission, the file can’t be backed up!
MAC versus DAC

• Who has the right to set file permissions?
  • Discretionary Access Control (DAC) — the file owner can set permissions
  • Mandatory Access Control (MAC) — only the security officer can set permissions
• Enforce site security rules
• Note: viruses and other malware change DAC permissions, but not MAC permissions
Implementing MAC

- Often side-by-side with DAC: system has both
- Processes need to pass both sets of permissions to access files
- Or — can have a special ACL-changing attribute in an ACL:
  
  `security_officer.wheel p`
- But — can `security_officer` give him/herself privileges?
- In reality, MAC is often used for classification levels (next class), rather than ACLs
Privileged Users

- Root or Administrator can override file permissions
- This is a serious security risk — there is no protection if a privileged account has been compromised
- There is also no protection against a rogue superuser...
- Secure operating systems do not have the concept of superusers
Database Access Control

- Often have their own security mechanisms
- Permit user logins, just like operating systems
- Some have groups as well
- Permissions are according to database concepts: protect rows and columns
- Different types of operations: select, insert, update, delete, and more
Databases versus OS Security

- The database has many objects in a single OS file
- The OS can control access to the file
- The DBMS has to control access to objects within the file
- The set of database users is not the same as the set of OS users
Access Control Formalisms

- Access control can be modeled formally. What does this buy us?
- There are theorems that can be proved.
- For example, if ACLs permit negation there are undecidable questions.
Access Control Formalisms (cont.)

- For the general case:
- Model using a Turing machine.
- Turing machine enters a special state if the access control is faulty.
- Contradiction!