

n2n:
A Relational Graphing Language

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December 17, 2014

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Chapter 1

Introduction

In today's world there is an escalating interest in relationships. Be it the connection of people on a social networking platform, of family members in a family tree, of variables in a mathematical equation, or the connection of trains in a city subway system, we are constantly trying to find networks of people and things and analyze how they interrelate. Most often, relationships are implemented in a programming setting via the graph data structure containing a set of nodes and edges defining connections between the nodes. In standard programming languages, however, graphs can be tedious to create and manipulate, requiring the creation of separate immutable classes for nodes and edges, and they tend not to focus on the possibility of having various types of nodes and relationships within one graph's structure. `n2n` is a language designed with the connected graph as the primary data structure providing a higher level of abstraction to perform common operations on connected graph data structures with a specific focus on how nodes of a graph interrelate, and the types of relationships they have.

1.1 Language Goals

Graphing in many languages can be difficult and tedious with regards to creating nodes, relating them, adding or removing data from them, and finding existing relationships between nodes. `n2n` was designed to be an intuitive language focused on programming directed connected graph data structures, not only making it easier to perform basic operations on the graph, but also providing the ability to specify types of nodes and relationships that can

model information stored in graphs in real settings.

1.2 Overview

The primary components of an n2n program are Graph objects and the Nodes and Relationships they contain. Node and Relationship (Rel) constructors are named and can hold a set of primitive data fields which must be specified when an literal of that Node or Rel type is instantiated. Associations are made between nodes in a graph by adding Node1-Relationship-Node2 tuples to a graph, which simultaneously adds both the nodes and the relationship to the graph. Inserting and removing data from a Node or Rel, like inserting and removing Node-Rel-Node tuples from a graph, is done with ease, as is performing searches and other graphing algorithms with the instantiated graphs.

Chapter 2

Language Description

An n2n graph is a collection of Node-Relationship-Node tuples. Node and relationship constructors are defined when their name and the data types they contain are specified. Once their types are created, Node-Rel-Node tuples can be added to a graph representing a relationship from the first Node to the second. The Nodes and Rels of the tuple are either named instances of the newly created Node and Rel constructors (ids), or unnamed literals of some Node or Rel constructor with the required data fields specified as defined by their constructors. The two node literals in the tuple do not have to be of the same type. It is possible to have more than one Relationship from Node A to Node B, and to have a different Relationship connecting Node B to Node A than exists from A to B.

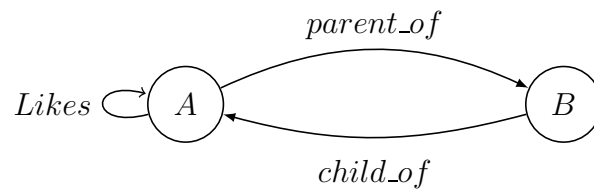


Figure 2.1: Simple Graph Example where one Relationship connects Node A to Node B

2.1 A More Descriptive Example

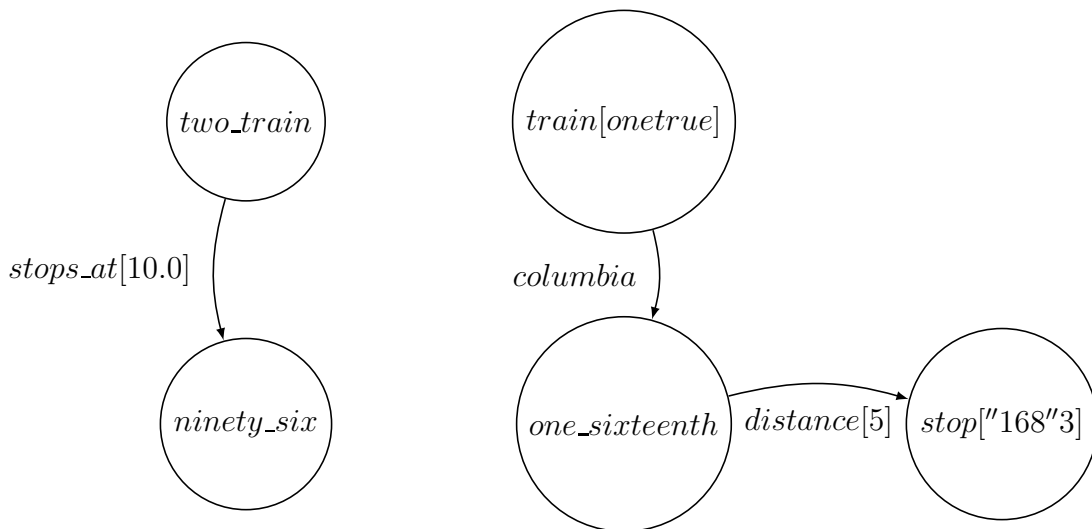
To make things clearer, consider the graph of the NYC subway system, specifically the red line (with 1, 2, and 3 trains). In the example below, we create a graph called *red_line*, Node constructors *subway_stop*, and *train*, and Rel constructors *stops_at*, and *distance*. The *subway_stop* Node type holds two data fields: a string that defines the name of a *subway_stop*, and an integer representing the number of different trains that stop there. When a new Node of type *subway_stop* is created such as the id named *one_sixteenth*, values for its string and integer data fields must be defined. When the *red_line* graph is created, the arguments of its constructor are Node-Rel-Node tuples which are either Node ids that have been instantiated already (as in the case with *one_sixteenth*), or a Node or Rel literal, defined within the Graph constructor by setting its primitive data fields within brackets right there! (This is done in the example when the *subway_stop*, *ninety_six* is inserted into the graph.)

```
subway_stop: Node = {name: String , num_trains_stop_here
  : Int}
train: Node = {line: String , local: Bool}
stops_at: Rel = {stop_rate: Double}
distance: Rel = {dist_by_num_stops: Int}

one_sixteenth: Node = subway_stop[ ''one_sixteenth'' , 1]
two_train: Node = train[ ''two'' , false]
columbia: Rel = stops_at[7.0]

red_line: Graph = <
  one_sixteenth distance[ 5 ] subway_stop[ ''
    one_sixty_eight'' , 3]
  two_train stops_at[ 10.0 ] subway_stop[ ''
    ninety_six'' , 4]
  train[ ''one'' , true] columbia one_sixteenth
>
```

Functions can then be defined to manipulate and perform graph algorithms on the defined Graphs, similar to the way functions are defined in languages such as C or Java. As *n2n* is specific to operating with Graphs, it supports basic binary and unary operations, but does not support features such as loops which are unnecessary in *n2n*. Graph algorithms that



would require looping in other languages, can be done more simply using n2ns specialized functions.

2.2 Built in Functions

One attribute of n2n that eases the manipulation of its Graph types is having built in functions that are native to the language. There are four built in functions:

1. A print function that takes in any valid expression or expression list as arguments and prints the value of the expression. For example,

```
print ( node_or_rel_instance )
```

prints the type of the node instance and the values of the data fields it holds.

```
print ( Graph )
```

prints the Node and Rel literals contained in the Graph
The print function will also print primitive data types as it does in Java.

2. A `find_many` function searches the graph for a specified pattern defined in the function arguments and returns a list of Nodes or Relationships. Arguments that the `find_many` function accepts are:
 - (a) `(node_literal)` returns a List of Nodes in the graph that match this `node_literal` with its data fields set to specified values.
 - (b) `(Node, Rel)` searches the graph for Nodes matching the passed Node id or literal that have a Rel (id or literal) relationship with one or more other Nodes in the graph. It returns a List of the Nodes on the other side of these relationships.
 - (c) `(Rel, Node)` does the inverse in that it searches the graph for Nodes being pointed at by the passed Rel and returns the Nodes that point to them.
 - (d) `(node, node)` returns a List of Rels that connect the first Node id or literal with the second.
3. A `map` function is called by a List or Graph of Node ids or constructors and accepts a set of statements as arguments. The function performs the statements on each Node in the List or Graph, and returns the resulting Graph or List (the Graph or List which is the result of performing the statements on the elements on the instances in the list).
4. A `neighbors` function is called by a Graph instance and accepts one Node instance as an argument. It returns a List of Nodes that argument has a Relationship towards within that caller graph (as in he has a relationship directed from himself to the these other nodes).

2.3 Compiling and Running an n2n Program

In the `src` file with the Makefile and the source code, run `make all` which compiles all n2n source files. Use the `./ n2n` operator to run your program as follows:

```
$ ./ n2n -[option] source\_files.n2n
```

This will compile the n2n code to Java and run the program.

Options for compiling n2n code:

1. `-a` for printing the AST

2. -s for running semantic analysis on the source code
3. -c is for compiling the whole program
4. -j is for compiling and running the java code of the program

2.4 Sample Code

This is a small program that tests the `find_many` function. It first creates two global Node types or constructors. Then within the main method it creates Node and Rel literals and inserts Node-Rel-Node tuples into the graph. It then calls the `find_many` function which finds all *actor* Nodes whose name data field is set to "Keanu" and whose age could be anything.

```
actor: Node = { name: String , age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String , year: Int };

fn main() -> Void {

    Keanu: Node = actor["Keanu", 35];
    Leo: Node = actor["Leo", 20];

    neo: Rel = actedIn["Neo"];
    jordan: Rel = actedIn["Jordan"];
    nelson: Rel = actedIn["Nelson"];

    matrix: Node = movie["Matrix", 1999];
    wolf: Node = movie["Wolf", 1994];
    sweet_nov: Node = movie["Sweet November",
        2000];

    Cast: Graph = <
        Keanu neo matrix ,
        Leo jordan wolf ,
        Keanu nelson sweet_nov
    >;
```

```
node_lit: List<Node>;
node_lit = Cast.find_many(actor["Keanu", -]);

print(node_lit);
print("\n");
```

This program outputs:

```
[actor{ name = Keanu, age = 35 }]
```

which is a list of the Nodes of the graph whose name is Keanu.

Chapter 3

Language Reference Manual

N2n enables the easy creation and manipulation of directed, connected graph types. Graphs, Nodes, and Relationships (Rel) are native data types in the language. Node and Rel types may enclose a set of primitive types. The language includes graph operators to add and remove Node-Rel-Node tuples from the graph, and element operators that allow the addition or subtraction of data fields from Node and Rel instances. Functions native to the language allow for easy searching and manipulating of the graph data structures. A programmer can create and manipulate the graphs while n2n handles their storage under the hood.

3.1 Lexical Conventions

An n2n program, written with the Unicode character set, consists of global variables and global functions that can be used in a main function from which the program is run. Its functions contain various types of tokens such as expression operators, keywords, built-in function names, identifiers, and other separators. All identifiers, keywords, and functions names are case sensitive. In general blanks, tabs, and comments are ignored except as they serve to separate tokens. Semi-colons indicate the end of a statement.

3.1.1 Comments

Comments are introduced and terminated by two semi-colons For example:

```
;; This is a comment ;;
```

3.1.2 Identifiers

An identifier is a sequence of letters and digits with no spaces between them. The first character must be alphabetic. Underscore characters may be included in the identifier.

3.1.3 Keywords

The following are reserved for use as keywords and may not be used otherwise:

Int	Graph	Void	true	find_many
Double	Node	if	false	main
String	Rel	else	map	return
Bool	List	fn	neighbors	

3.1.4 Constants

An integer constant is a sequence of digits that are assumed to be in base 10. A double constant represents a fractional number. It consists of an integer part, a decimal point, and a fractional part. Either the integer part or the fraction part but not both may be missing. Both are comprised of a series of digits. Boolean constants may hold only the values true or false. String constants are one or more characters surrounded by double quotes. Special characters may be escaped within a string the same as they can be escaped in Java.

3.2 Identifiers, Overloading, and Scoping. Oh my!

The type of the identifier determines the meaning of the values in the identifiers storage. The location and lifetime of an identifier are determined by the scope in which it is declared. A function sets a scope for a variable so that any variable declared in a function only exists within that function and are discarded on return. Any variable declared outside functions are global and visible by the entire program, independent of any function. If a local variable shares its name with a global variable, the local variable (within the function) has precedence.

3.3 Data Types

3.3.1 Primitive types

Primitive types include Int, Double, Bool, and Strings. The type Void can be used as a return type for functions.

3.3.2 Complex types

Complex types include Graph, Node, Rel, and List

1. Node literals are the elements stored in graphs. A Node constructor define a Node type and is created when its name and the primitive data fields it encapsulates are defined. When a Node literal is created (which may or may not have an id), it must provide values for the data fields defined in its Node constructor.
2. Rel is a directed, named relationship between Nodes. There can be a different Rel between Node1 and Node2 than exists between Node2 and Node1. A Rel type is declared and defined like Nodes with a name and primitive data fields associated with the Rel.
3. Graph is a collection of Node-Rel-Node tuples. It can be thought of as a list of relationships, and the nodes that the names relationships connect.
4. List maintains an ordered collection of elements

3.4 Declarations

3.4.1 Variable Declaration and Assignment

The type of the identifier determines the meaning of the values in the identifiers storage. Thus the type must be explicitly specified when declaring an identifier as follows:

Declaration

An identifier name, followed by a colon, its type and possibly followed by an assignment via the assignment operator (=).

```
var_name : type
```

Or

```
var_name : type = expr
```

(where expr here refers to some expression being assigned to the identifier).

3.4.2 Assignment

Once an identifier has been declared, to assign or reassign its value simply use its name followed by the assignment operator and an expression. For example:

```
var_name = expr
```

More specifically:

1. Nodes and Rel: There are two types of Node/Rel declarations:

- (a) The declaration of a Node constructor

To define a new Node or Rel node type, declare an identifier as above by specifying its name followed by a colon and its type, the assign character, = and the declaration of data fields within curly braces. For example:

```
node_constr1 : Node = {s: String , b: Bool}
```

- (b) The declaration of Node literal.

A Node literal can be declared explicitly by assigning it to an id, or instantiated within Graph objects. When declared explicitly, the expression following Node objects assignments is surrounded by square brackets in which the fields of data objects associated with that node are set. For example:

```
node_id_of_type1 : Node = node_type1 [ 'my_name' , true ]
```

2. Graph: To create a graph, declare its name and type (Graph) as above. Following the assignment operator there is a list of zero or more Node-Rel-Node tuples surrounded by <and >. The Nodes and Rels in the

tuples can either be ids representing pre-existing literals or definitions of new Node/Rel literals whose constructors are called and data fields defined within the creation of this graph. For example:

```
graph1: Graph = <
    keanu    acted_in [ 'John' ]    movie [ '
        Constantine' , 2003 ],
    actor [ 'Keanu' ] acted_in [ 'Neo' ]    movie
        [ 'Matrix' , 1999 ]
>
```

This creates a graph, *graph1*, with two Node-Rel-Node relationships. The first has has a Node id *keanu* which was defined before the declaration of *graph1*, as well as a Node literal of type *movie* whose String and Int data fields are given values within this instantiation of *graph1*.

3.4.3 Function Declaration

A function is declared by specifying the name, parameters, and return type of the function as follows:

```
fn func_name ( parameter-list ) -> return_type {
    statements }
```

The parameters are formatted as follows:

```
name : type
```

and are separated by commas. The body of the function is a list of statements that may include new variable declarations and function calls. Functions may be recursive. Note that functions create scope so that a variable declared within a function with the same name as a global variable defined outside of a function will take precedence over the global within that function. (More in section Scope) There must be a function called *main* in an *n2n* program from which the program is run

3.5 Expressions

1. Identifier: An identifier is an expression provided it has been suitably declared.

2. Primitive Literals of types Int, Double, Bool, and String are also expressions.
3. Complex Literals of types Graph, Node, Rel, or List
4. Binary Operators include Additive and Multiplicative operators as enumerated below
5. Unary Operators
6. Graph Operators
7. Node and Rel Operators
8. Access: To access a primitive data field of a Node or Rel instance use the dot operator as in:

```
node1.named_attribute
```

node1 must be an existing Node id that has a primitive data field with the name *named_attribute*. The access expression may be followed by the assignment operator (=) in order to change the value of the primitive data type held in the Node or Rel. For example:

```
keanu.visited = true
```

updates the "visited" field in the *keanu* node to be true

9. Call
To call an existing n2n function, type the name of the function followed by the function arguments (if any) within parentheses. For example:
- ```
dijkstra(graph1);
```
10. Built-in function calls including the *print*, *find\_many*, *neighbors*, *map* functions

## 3.6 Operators

### 3.6.1 Unary Operators

*-expression* results in the negative of the expression

*! expression* results in true if the value of the expression is false and true if the expression value is false

### 3.6.2 Multiplicative Operators

|                         |
|-------------------------|
| Expression * expression |
|-------------------------|

The \* operator indicates multiplication. If both operands are of type Int the result is Int. Same with Double. If one is Int and the other Double, the result is Double.

/indicates division. The same type considerations as for multiplication apply. The % operator indicates the remainder from the division of the first expression by the second. Both operands must be of type Int.

### 3.6.3 Additive Operators

*expression + expression*: The result is the sum of the expressions. The same type considerations as for multiplication apply.

*expression - expression*: The result is the difference of the operands.

### 3.6.4 Relational Operators

$e < e$  and  $e > e$  and  $e \leq e$  and  $e \geq e$  all yield false if the specified relation is false and true if it is true. These operators can only be applied for types Int and Double.

Same for the equality operators: == (equal to) and != (not equal to). *expression && expression* returns true if both expressions are true and false otherwise

*Expression '—'—' expression* returns true if either expression is true and false otherwise

### 3.6.5 Graph and Node Operators

Node-Rel-Node tuples may be added or removed from a Graph using the  $\hat{+}$  and  $\hat{-}$  operators. For example:

```
Graph_name $\hat{+}$ node1 rel1 node2
```

The format of the Node-Rel-Node tuple is the same as in Graph instantiation (above). To add or remove a data field from a Node or Rel instance, use the  $[+]$  and  $[-]$  operators respectively. The data field should have the form name: type

```
node_example $[+]$ visited : Bool
```

## 3.7 Built in Functions

One attribute of n2n that eases the manipulation of its Graph types is having built in functions that are native to the language. There are four built in functions:

1. A print function that takes in any valid expression or expression list as arguments and prints the value of the expression. For example,

```
print (node_or_rel_instance)
```

prints the type of the node instance and the values of the data fields it holds.

```
print (Graph)
```

prints the Node and Rel literals contained in the Graph  
The print function will also print primitive data types as it does in Java.

2. A find\_many function searches the graph for a specified pattern defined in the function arguments and returns a list of Nodes or Relationships. Arguments that the find\_many function accepts are:
  - (a) (node\_literal) returns a List of Nodes in the graph that match this node\_literal with its data fields set to some value. The underscore character is used to match again any value.

For example, a graph with an actor Node type or constructor may contain multiple actor literals whose data fields are set differently. In this example consider an actor that contains types name:String and age:Int. To find all actor Nodes whose first argument, *name*, is set to "John" and whose second argument *age* is set to 25:

```
john_actors: List<Node> = find_many(actor [' '
 John ' ' ,25]);
```

To pattern match against fewer fields of the Node literal, simply replace a data field with an underscore. For example, to find all *actors* named John regardless of their age:

```
john_actors: List<Node> = find_many(actor [' '
 John ' ' , -]);
```

Pattern matching against all underscores will return all Nodes of type *actor*.

- (b) (Node, Rel) searches the graph for Nodes matching the passed Node id or literal that have a Rel (id or literal) relationship with one or more other Nodes in the graph. It returns a List of the Nodes on the other side of these relationships. For example, in a graph of *movie* and *actor* constructors, to find movie Nodes for which *keanu* has an *acted\_in* relationship with:

```
keanu_movies: List<Node> = find_many(keanu
 acted_in)
```

This will search the graph for matches on a node *keanu* that has a relationship of type *acted\_in* and return a List of the *movie* Nodes on the other side of the relationship.

- (c) (Rel, Node) does the inverse in that it searches the graph for Nodes being pointed at by the passed Rel and returns the Nodes that point to them. For example:

```
neo_actors: List<Node> = find_many(acted_in [' '
 Neo ' '] matrix)
```

Finds Nodes that have a relationship of *acted\_in* with its data field set to "Neo" in Node matrix. Here too, data fields can be replaced

with underscores to match against zero or more data fields of the Node or Rel literals.

- (d) `(node, node)` returns a List of Rels that connect the first Node id or literal with the second. For example: Both

```
connection: List<Node> = find_many(actor['Neo']_matrix)
```

returns all Rels between *actors* whose name field is set to "Neo" and the Node id named *matrix* and

```
relationship: List<Node> = find_many(movie[_]actor[_])
```

returns all Rels between any *movie* Node and any *actor* Node.

3. A *map* function is called by a List or Graph of Node ids or constructors and accepts a set of statements as arguments. The function performs the statements on each Node in the List or Graph, and returns the resulting Graph or List (the Graph or List which is the result of performing the statements on the elements on the instances in the list). For example:

```
mutated_graph: Graph = Cast.map(node in {node [+]
visited: Bool;});
```

returns a Graph called *mutated\_graph* which adds a boolean field called *visited* to each Node in graph *Cast*.

4. A *neighbors* function is called by a Graph instance and accepts one Node instance as an argument. It returns a List of Nodes that argument has a Relationship towards within that caller graph (as in he has a relationship directed from himself to the these other nodes).

## 3.8 Statements

A statement is anything that can make up a line of n2n code.

### 3.8.1 Statement Blocks

A block of code is any list of statements surrounded by curly braces as shown:

```
{
 statement ;
 statement ;
 ...
}
```

### 3.8.2 Expressions

Any of the named expressions given above are expression statements if followed by a semicolon.

### 3.8.3 Conditional statements

Conditional statements are composed of two parts: a required if; and an optional else part.

1. if (expression) {  
    statements  
}
2. if (expression) {  
    statements  
} else {  
    statements  
}

In the first case, the compiler will execute *statements* if the result of evaluating *expression* is the boolean value of true.

The second case is like the first but if the result of evaluating *expression* is the boolean value of false, it will evaluate the *statements* inside the else part of the conditional.

### **3.8.4 Return Statements**

To return a function to its caller use the keyword *return* as follows:

```
return expression
```

This will return the result of evaluating expression.

### **3.8.5 Variable Declarations**

The declaration of a variable is a statement that can exist in an n2n program either globally or within functions definitions.



# Chapter 4

## From Plan to Finished Product

### 4.1 Planning

The planning of our project consisted of coming up with a language idea, determining team responsibilities, and setting a timeline for the execution of our language compiler. We had weekly, then biweekly in person meetings to plan, discuss, and execute our project. Over the course of the semester many aspects of our initial plan changed. Although we initially agreed upon a language and its specifications, our initial model was worked and reworked and expanded and reconfigured many times throughout the semester as the project progressed.

### 4.2 Development

At the beginning of each project stage, each team member took on a portion of the project and worked on it both individually and with other team members. For many parts of the project our designated roles overlapped; we worked together, we corrected each others mistakes, and redesignated tasks, though often our task assignments overlapped. It would be difficult to outline each team members responsibilities as most pieces of our compiler were touched by almost all of our hands. We met once, and, sometimes, twice every week to discuss issues that we had with the code, to formulate new decisions about the project, and to keep one another up to date with new developments. We also used git with Github for our version control so that each team member had access to our entire code base at all times.

### 4.3 Team Responsibilities

The initial responsibilities designated to each team member were generally not binding, as each team member took on various roles as our project progressed throughout the semester. We were each involved with multiple parts of the project, which are delineated below:

| Team Member    | Roles                                                 |
|----------------|-------------------------------------------------------|
| Nicholas Falba | Parser, Semantic Check, SAST                          |
| Johan Mena     | Compiler Executable, Java Backend, Testing, Git Ninja |
| Jialun Liu     | Scanner, Parser, AST, Testing                         |
| Elishева Aeder | AST, SAST, Code Gen, Final Report                     |

### 4.4 Project Timeline

The following was the initial timeline the we set out for our project: The

| Date         | Accomplishment                 |
|--------------|--------------------------------|
| September 10 | First Meeting                  |
| September 24 | Project Proposal               |
| October 20   | Scanner, Parser, AST Completed |
| October 27   | Language Reference Manual      |
| November 24  | Semantic Check Completed       |
| December 1   | Code Generation Completed      |
| December 10  | Testing Completed              |
| December 17  | Final Report and Presentation  |

actual flow of events was not as smooth as the above. While coding each new part of the compiler, previously completed sections had to be revised and updated based on new overall language updates and newly discovered errors. Most aspects of the compiler were not fully complete until the final stages of testing.

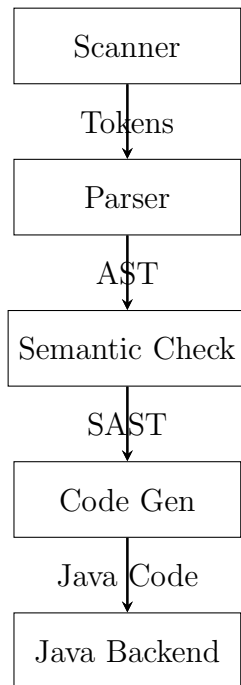
### 4.5 Development Environment

N2n was developed in the Mac OS X environment. The compiler was written in OCaml, using `ocamlc` to compile `.ml` files, `ocamllex` for the scanner, and

ocamlyacc for the parser. The generated Java code is compiled using the java compiler. The compiling process was automated with Makefiles. Testing and verification was run with shell scripts. Github was used to host our git repository.

# Chapter 5

## Architectural Design



The n2n compiler takes a .n2n file as input into the scanner which outputs tokens. It then gets passed to the parser which outputs an AST. It passes through the semantic check and produces an SAST and then the code generation outputs Java source code which is compiled and run. The code generation includes a Java library which allows for the functions required for Graph, Node, and Rel creation and manipulation.

## 5.1 Scanner

The scanner takes in a N2N source file and outputs tokens according to the lexical rules of our language.

## 5.2 Parser

The parser takes in the tokens from the scanner and uses our grammatical rules to decide if the input are in the language that is specified by our context free grammar. During parsing, it produces an abstract syntax tree that defines the structure of the given program.

## 5.3 Abstract Syntax Tree

The AST defines the structure of the N2N language, including various types of structures, for example, primitive and complex types, builtin function calls, user-defined functions and the program itself.

## 5.4 Semantic Check

The semantic check receives an abstract syntax tree as an argument and prints a semantically checked tree in a similar format. For the AST functions, their semantically checked counterparts (SAST functions), hold more information. In the case of expressions, this is useful for type checking. And in the case of Graph Elements (Nodes and Relationships), we keep track of element type, identifiers, and all the field information for each Node or Rel across the semantic check to assist in our back-end methods. To keep track of which identifiers exist and the types for them, we store this information in a list in the environment. The types are segregated in the environment based on what needs to be stored there. Nodes and relationships each have their own section, since the types are so similar in structure but used differently. This is to avoid confusion when checking, since theyre stored in separate lists. Graphs are stored in their own area since you store Node, Rel, Node triples in them, in addition to an Id. Lists are stored separately since they are the only type you cant mutate directly. And the rest of the types (Int, Bool, String, and Double) have their own table since they all need to keep similar

information. And there are two sets of tables, one for the global and one for the local scope, as different variables will be visible depending on the scope. Moreover, functions are defined and added to the environment as well as their return type, so you can check to make sure a function is returning what it is supposed to. Variables and functions are added to the variable tables when they are instantiated and defined, respectively. The semantic check makes sure that only certain expressions appear in places that are allowed by the syntax and semantics of the language, and makes sure the output of individual expressions and statements is what is expected by the language. The output is a set of semantically checked functions in tree format used by our code generating function to output java code.

## 5.5 Code Generation (`code_gen.ml`)

The `code_gen` takes the semantically checked abstract syntax tree generated from the semantic check and then converts it to the java code. Here we match types of expressions with primitive types(`Int`, `Double`, `String`, `Bool`) and complex types(`Node`, `Rel`, `Graph`, `List`); if we find a primitive type we just insert the java equivalent syntax, if a complex type we then need to insert the code according to methods we defined in our java backend.

## 5.6 Java Backend

(`Graph.java` `Node.java` `Relationship.java`) The backend has three infrastructure classes:

### 5.6.1 Relationship

Relationships are the glue that holds all the nodes that are connected to it.

### 5.6.2 Node

`Node` is the class that holds the data for each one of the members of the `Graph`. `Nodes` and `relationships` also keep track of the data they contain.

### 5.6.3 Graph

Graph is the main class, and is used to store the set of relationships that connect all of the nodes in the graph. It also provides overloaded finder methods and functionality for adding, removing or printing out Relationships and Nodes. Most methods were implemented using Java 8's stream api, for readability and performance. There are also unit tests for most of the finder methods, to ensure reliability of the functionality.

## 5.7 Command Line Interface (n2n.ml)

The N2N command line is the file that puts all of these modules together in succession. It opens the source program that has been passed to it and runs it through the scanner, parser, semantic check and code generation to produce a java program and finally compiles the program using the java. Our compiler supports several different functionalities, for instance, printing the AST of the program, running semantic check over the source program, generating the java program that then be compiled.

# Chapter 6

## Test Plan

### 6.1 Development of the Test Plan

Our test plan consists of a test suite of comprehensive tests meant to check all of the functionalities of a n2n program, one script to test each stage of the compiler, e.g. printing the AST, running semantic check over the source program and generating the java code from the source program. We also have an automated test script that runs over all of the test cases and checks them against the expected output files. The test suite and scripts were primarily authored by Nicholas Ray Falba, Johan M. Mena and Jialun Liu.

### 6.2 Test Suites

Our test suite is listed in the table below. These test cases were chosen because they represent the smallest building blocks of a n2n program, making it possible for us to test each functionality individually and easily identify the location of bugs. There are also some more complicated tests which allow for testing multiple functionalities of the program integrated together.

### 6.3 Component Testing

```
./test.sh [option] source_files.n2n
```



| File            | Functionality Tested                                                     |
|-----------------|--------------------------------------------------------------------------|
| arithmetic1.n2n | Checks basic arithmetic, adding two numbers and printing the result      |
| arithmetic2.n2n | Checks basic arithmetic with a number of operations                      |
| comparison1.n2n | Checks comparison between two numbers and prints correct Bool output     |
| comparison2.n2n | Checks equality operations on numbers and prints the correct Bool output |
| declaration.n2n | Checks local and global variable declaration and scope visibility        |
| escapestr1.n2n  | Checks to make sure that new line escape strings work                    |
| escapestr2.n2n  | Checks to make sure that tab escape strings work                         |
| escapestr3.n2n  | Checks non-main function call to print escape string                     |
| escapestr5.n2n  | Same as above except different escape string                             |
| float.n2n       | Checks output of integer division                                        |
| float2.n2n      | Checks output of same expression above except for Doubles                |
| helloworld.n2n  | Checks print of "Hello World"                                            |
| init1.n2n       | Checks to see if we can declare and initialize at the same time          |
| functions1.n2n  | Checks function call                                                     |
| functions2.n2n  | Testing assignment to the output of a function call                      |
| functions3.n2n  | Testing nested function calls                                            |
| find_many1.n2n  | Testing find_many for (Node, Node) case                                  |
| find_many2.n2n  | Testing find_many for (Node, Rel) case                                   |
| find_many3.n2n  | Testing find_many for (Rel, Node) case                                   |
| find_many4.n2n  | Testing find_many for (Node.literal) case                                |
| graph.n2n       | Testing graph instantiation for Node and Rel variables                   |
| graph2.n2n      | Testing graph instantiation for Node and Rel literals                    |
| graphInsDel.n2n | Testing graph insert and remove after instantiation                      |
| if.n2n          | Testing if statements and dangling else problem                          |
| map.n2n         | Test built-in function map                                               |
| neighbor.n2n    | Testing basic neighbor function call                                     |
| node.n2n        | Testing node declaration and instantiation                               |
| nodeAccess.n2n  | Testing access on field of a Node                                        |
| nodeInsDel.n2n  | Testing Data.InsertRemove operators                                      |
| rel.n2n         | Testing rel instantiation and initialization                             |
| relAccess.n2n   | Testing access on field of a Rel                                         |
| simple.n2n      | Testing multiple global variable declaration                             |
| simple2.n2n     | Testing basic local variable assignment                                  |

test.sh - semantic check – outputs semantic check trace - all methods of semantic check that have been visited for a program and if there's error it

prints error message and if no error it prints "passed." The option is either `semantic_check` in which case the above program is run and the output trace sent to a file at a given location if there is an error. Or you could run it with `ast` as an argument in which case it just runs the print function of the `ast`.

```
#!/bin/bash

rm -f ../test/sem_output/*.out
echo

arg_array=($@)
test_instruction=${arg_array[0]}
test_message=""
test_output=""
for ((i = 1; i < $#; i++))
do
if [[$test_instruction == "ast"]]
then
 test_message="Printing ast for ${arg_array[$i]}: "
 test_output=$(./n2n -a ${arg_array[i]} 2>&1)
elif [[$test_instruction == "semantic_check"]]
then
 test_message="Checking semantic_check for ${arg_array[$i]}: "
 test_output=$(./n2n -s ${arg_array[$i]} 2>&1)
elif [[$test_instruction == "java"]]
then
 test_message="Generating java code for ${arg_array[$i]}: "
 if [[${arg_array[$i+1]} == *.java]]
 then
 test_output=$(./n2n -j ${arg_array[$i]}
 ${arg_array[$i+1]} 2>&1)
 else
 test_output=$(./n2n -j ${arg_array[$i]}
 2>&1)
 fi
fi
```

```

fi
if [[$test_output == *"Parse_error"*]]
then
 echo $test_message"Failed Parsing"
 echo
elif [[$test_instruction != "ast" && $test_output !=
 "Passed Semantic Analysis"]]
then
 filename=${arg_array[$i]}
 back_filename=${filename##*/}
 echo $test_message"Failed Semantic Check.
 Output in file ../test/sem_output/${
 back_filename%.*}.out"
 echo
 printf "$test_output" > ../test/sem_output/${
 back_filename%.*}.out
else
 if [[$test_instruction == "ast"]]
 then
 echo $test_message"Passed."
 echo
 printf "$test_output\n"
 echo
 else
 echo $test_message"Passed."
 echo
 fi
fi
done

```

## 6.4 Automation

We use an automated test script that takes .n2n files, compiles them into Java code, runs the programs, and compares their outputs with the expected outputs. It prints ok if the tests are correct, and returns the difference if not. The automated test script is run via the command:

```
./test_all.sh [source files]
```

And this is its code:

```
#!/bin/bash

COMPILER=" ./n2n_c"
TESTFILES=$@

rm errors.out

for f in $TESTFILES
do
 LEN=$((${#f} - 4)
 OUTFILENAME="${f:0:$LEN}.output"
 TESTFILENAME="${f:0:$LEN}.out"
 echo -n "Test_${f}:"
 $COMPILER $f > $OUTFILENAME 2>> errors.out

 DIFF=$(diff -w $OUTFILENAME $TESTFILENAME)
 if ["$DIFF" = ""]
 then
 echo "_OK"
 else
 echo "_FAIL"
 echo "====="
 echo "Expected:"
 echo "$(cat _${TESTFILENAME})"
 echo "====="
 echo "But_was:"
 echo "$(cat _${OUTFILENAME})"
 fi

 rm -f $OUTFILENAME
done
```

## 6.5 n2n to Java

### 6.5.1 DFS Search Algorithm

The following n2n program performs a DFS search on a graph and prints all Nodes in the graph:

```
member: Node = {Name: String, Age: Int, visited: Bool};
relation: Rel = {Relation: String};

fn DFS(N: Node, M: Graph)->Int {
 if(N.visited == false){
 print(N);
 print("\n");
 Neighbor: List<Node> = M.neighbors(N);
 N.visited = true;
 Neighbor.map(node in { DFS(node, M); })
 ;
 }
 return 0;
}

fn main()->Void {
 John: Node = member["John_Hamilton", 49, false
];
 Mary: Node = member["Mary_Lance", 47, false];
 Johan: Node = member["Johan_Hamilton", 21,
 false];
 Sara: Node = member["Sara_Hamilton", 20, false
];

 Mother: Rel = relation["Mother_of"];
 Father: Rel = relation["Father_of"];
 Child: Rel = relation["Child_of"];
 Brother: Rel = relation["Brother_of"];
 Sister: Rel = relation["Sister_of"];

 Family: Graph = <
 John Father Johan,
```

```

 John Father Sara ,
 Mary Mother Johan ,
 Mary Mother Sara ,
 Johan Brother Sara ,
 Johan Child John ,
 Johan Child Mary ,
 Sara Sister Johan ,
 Sara Child John ,
 Sara Child Mary
 >;

 DFS(Johan , Family);
}

```

The output of this n2n program looks like:

```

member{ visited = false , Age = 21, Name = Johan
 Hamilton }
member{ visited = false , Age = 47, Name = Mary Lance }
member{ visited = false , Age = 20, Name = Sara Hamilton
 }
member{ visited = false , Age = 49, Name = John Hamilton
 }

```

The Java code that this program compiles to is:

```

package com.n2n;

import java.util.*;

class Main {

 String relation = "relation";
 String member = "member";
 ;
 public static void main(String[] args) {

 Node John = new Node("member", new
 HashMap<String , Object>() {{
 put("Name" , "John.LHamilton");
 }}
 }
}

```

```

put("Age", 49);
put("visited", false);

}}
);;
 Node Mary = new Node("member", new HashMap<
 String, Object>() {{
 put("Name", "Mary_Lance");
put("Age", 47);
put("visited", false);

}}
);;
 Node Johan = new Node("member", new HashMap<
 String, Object>() {{
 put("Name", "Johan_Hamilton");
put("Age", 21);
put("visited", false);

}}
);;
 Node Sara = new Node("member", new HashMap<
 String, Object>() {{
 put("Name", "Sara_Hamilton");
put("Age", 20);
put("visited", false);

}}
);;
 Relationship Mother = new Relationship("
 relation", new HashMap<String, Object>() {{
 put("Relation", "Mother_of");

}}
);;

 Relationship Father = new Relationship("relation"
 , new HashMap<String, Object>() {{

```

```

 put("Relation", "Father_of");
 }}
);

 Relationship Child = new Relationship("relation", new HashMap<String, Object>() {{
 put("Relation", "Child_of");
 }}
);

 Relationship Brother = new Relationship("relation", new HashMap<String, Object>() {{
 put("Relation", "Brother_of");
 }}
);

 Relationship Sister = new Relationship("relation", new HashMap<String, Object>() {{
 put("Relation", "Sister_of");
 }}
);

 Graph Family = new Graph(Arrays.asList(new
 Graph.Member<>(John, Father, Johan), new
 Graph.Member<>(John, Father, Sara), new
 Graph.Member<>(Mary, Mother, Johan), new
 Graph.Member<>(Mary, Mother, Sara), new
 Graph.Member<>(Johan, Brother, Sara), new
 Graph.Member<>(Johan, Child, John), new
 Graph.Member<>(Johan, Child, Mary), new
 Graph.Member<>(Sara, Sister, Johan), new
 Graph.Member<>(Sara, Child, John), new Graph
 .Member<>(Sara, Child, Mary)));
 DFS(Johan, Family);
}
public static int DFS(Node N, Graph M) {
 if(N.getValueFor("visited").equals(false)) {
 System.out.print(N);
 }
}

```



```
 System.out.print("\n");
 Set<Node> Neighbor = M.neighbors(N);
 N.getData().put("visited", true);
 for(Node node : Neighbor){
DFS(node, M);
 }

 }
 else {
 }

return 0;
 }
}
```

This does not include the Java backend files which are included below in the appendix.

# Chapter 7

## Lessons Learned

### 7.1 Nicholas Falba

Its a cliché to say, start earlier, but its true. Specifically, get your front-end working as early as possible. And once you have a solid front-end (Parser, Scanner, AST), writing the semantic check will be way simpler and you should start on that too. This way you can get a feel for whats going to work and what isnt early on. This will give you a head start and crucial time management information in implementing the full check.

Also, practice Ocaml. Even a little bit. Learning by doing is how I wrote the semantic check. You dont want to do that. Learn beforehand.

Also resolve group differences early or actually make your group run like a dictatorship. Very early on, if you can agree on a syntax, itll go a long way. Dont wait until a month in to voice your concerns.

Also learn all about git. We only had one person in our group who knew git and it caused a bunch of minor roadblocks, but these could and should have been easily avoided.

### 7.2 Jialun Liu

This course is very interesting and I have definitely learned a lot. The most important thing is really to think early and start early, meaning that we should carefully consider what the syntax and rules we are going to use for our language. Finding out some errors in parser and insufficiency of the rule defined are going to a disaster in the later stage of the project, it makes us

change everything in the flow of the compiler and it makes us really suffered. Setting up the flow of the entire compiler is another important issue, we should have set it up at the very early stage so that it would thing clear about what we are really doing and figure out bugs a lot easier. Testing a very important issue of every project, as the tester, I didnt do well in this part, I should have conducted the tests as we moving along different parts of the compiler. Also, shell scripts for automated tests is a very useful thing it saved us a lot of effort of going through and debugging the compiler. Also, make some intuitive comment in the semantic check would be very helpful when the source program fails to go through the semantic check. Also, Ocaml is a very powerful language, all the pattern matching and function nesting make the language extremely useful. It would be a good start to get yourself familiar with Ocaml since that will save quite a lot a effort when writing semantic check in the later stage of the project. And it would definitely be a good idea to spend more time to play around with that after the project. In terms of project management, we had a great time working together, but we should have made a clear deadline for each part of the project so that it wouldnt become the case that we have to do a whole bunch of things in the last two weeks of the semester.

### **7.3 Johan Mena**

Don't pay attention to the lessons of othersit worked for them but it might not work for you. Pick your teammates carefully. Or maybe just at random.

### **7.4 Elisheva Aeder**

Learn Ocaml! And fast! One of my biggest struggles throughout the project was my lack of fluency in Ocaml. Spending more time on it in earlier stages would have been helpful. Also, presentations and reports take a lot longer than you'd think. Start earlier. Assign a leader and let them lead. Or have better defined roles.

# Chapter 8

## Appendix

### 8.1 Scanner

```
{
 open Parser
 exception LexError of string

 let verify_escape s =
 if String.length s = 1 then (String.get s 0)
 else
 match s with
 | "\\n" -> '\n'
 | "\\t" -> '\t'
 (*| "\\\" -> '\"'*)
 | "\\'" -> '\''
 (*| "\\`" -> '`'*)
 | c -> raise (Failure("unsupported character " ^ c))
}

(* Regular Definitions *)

let digit = ['0'-'9']
let decimal = ((digit+ '.' digit*) | ('.' digit+))

(* Regular Rules *)
```

```

rule token = parse
 | ['_ ' '\n' '\t '] { token lexbuf }
 | ';' { TERMINATION } (*
 terminate the code*)
 | ";;" { block_comment lexbuf }
 (* block comment*)

 | '(' { LPAREN }
 | ')' { RPAREN }
 | '{' { LBRACE }
 | '}' { RBRACE }
 | ']' { RBRACKET }
 | '[' { LBRACKET }

 | '+' { PLUS }
 | '-' { MINUS }
 | '*' { TIMES }
 | '/' { DIVIDE }
 | '%' { MOD }

 (* | ';' { SEMI }*)
 | ':' { COLON }
 | "," { COMMA }
 | '=' { ASSIGN }
 | "->" { ARROW }
 | "^" { CONCAT }
 | "." { ACCESS }

 | "==" { EQ }
 | "!=" { NEQ }
 | '<' { LT }
 | "<=" { LEQ }
 | ">" { GT }
 | ">=" { GEQ }
 | "!" { NOT }
 | "&&" { AND }
 | "||" { OR }

```

|    |                                       |                  |                   |
|----|---------------------------------------|------------------|-------------------|
|    | "if"                                  | { IF }           |                   |
|    | "else"                                | { ELSE }         |                   |
| (* | "elif"                                | { ELIF }*)       |                   |
|    | "map"                                 | { MAP }          |                   |
|    | "find_many"                           | { FINDMANY }     |                   |
|    | "fn"                                  | { FUNCTION }     |                   |
|    | "return"                              | { RETURN }       |                   |
|    | "in"                                  | { IN }           |                   |
|    | "^+"                                  | { GRAPH_INSERT } |                   |
|    | "^_"                                  | { GRAPH_REMOVE } |                   |
|    | "["+]"                                | { DATA_INSERT }  |                   |
|    | "[-]"                                 | { DATA_REMOVE }  |                   |
|    | "neighbors"                           | { NEIGHBORS }    |                   |
|    | "Graph"                               | { GRAPH }        |                   |
|    | "Rel"                                 | { REL }          |                   |
|    | "Node"                                | { NODE }         |                   |
|    | "List"                                | { LIST }         |                   |
|    | "Int"                                 | { INT }          |                   |
|    | "Double"                              | { DOUBLE }       |                   |
|    | "String"                              | { STRING }       |                   |
|    | "Bool"                                | { BOOL }         |                   |
|    | "Void"                                | { VOID }         |                   |
|    | '_'                                   | { ANY }          |                   |
|    | digit+ as lit                         |                  | {                 |
|    | INT_LITERAL(int_of_string lit) }      |                  |                   |
|    | decimal as lit                        |                  | {                 |
|    | DOUBLE_LITERAL(float_of_string lit) } |                  |                   |
|    | ''' ([^''']* as lit) '''              |                  | { STRING_LITERAL( |

```

 lit) }
| ("true" | "false") as lit
 { BOOLLITERAL(
 bool_of_string lit) }
| ['a'-'z' 'A'-'Z']['a'-'z' 'A'-'Z' '0'-'9' '_']* as
lit { ID(lit) } (*every ID should start with a
letter*)
| eof { EOF }
| - as char { raise (Failure("illegal character " ^
Char.escaped char)) }

and block_comment = parse
 ";" { token lexbuf }
| eof { raise (LexError("unterminated block_comment!"))
}
| - { block_comment lexbuf }

(*
| "new_node"

 { NEWNODE }
| "new_rel"

 { NEW_REL }
| "add_field"

 { ADD_FIELD}
*)

```

## 8.2 Parser

```

%{
 open Ast
%}

```

```

%token LPAREN RPAREN LBRACE RBRACE LBRACKET RBRACKET
 ANY
%token TERMINATION COMMA ASSIGN COLON ARROW CONCAT
 ACCESS
%token PLUS MINUS TIMES DIVIDE MOD
%token EQ NEQ LT LEQ GT GEQ AND OR NOT
%token IF ELSE
%token MAP IN FINDMANY
%token FUNCTION RETURN
%token GRAPHINSERT GRAPHREMOVE DATAINSERT
 DATAREMOVE NEIGHBORS
%token GRAPH REL NODE INT DOUBLE STRING BOOL VOID LIST
%token <int> INT_LITERAL
%token <string> STRING_LITERAL ID
%token <float> DOUBLE_LITERAL
%token <bool> BOOL_LITERAL
%token EOF

%nonassoc NOELSE
%nonassoc ELSE
%right GRAPHINSERT GRAPHREMOVE DATAINSERT
 DATAREMOVE
%left OR
%left AND
%left EQ NEQ
%left LT GT LEQ GEQ
%left CONCAT
%left PLUS MINUS
%left TIMES DIVIDE MOD
%left NEG NOT

%start program
%type <Ast.program> program

%%

program:
| /* nothing */ { ([, []) }

```



```

| program global_var_declaration { ($2 :: fst $1), snd
 $1 }
| program function_declaration { fst $1, ($2 :: snd
 $1) }

var_declaration:
| ID COLON n2n_type
 { Var($3,
 $1) } /* foo: String */
| ID COLON n2n_type ASSIGN LBRACE formal_list RBRACE
 { Constructor($3, $1, List.rev $6)} /* movie:
 Node = { title: String, year: Int } */
| expr ASSIGN expr
 {
 Access_Assign($1, $3) } /* foo: String = "lolomg",
 matrix: Node = movie[Matrix , 1999] */
| ID COLON n2n_type ASSIGN expr
 { Var_Decl_Assign($1, $3,
 $5)} /* Split the declaration and initializatio
 for complex */

global_var_declaration:
 var_declaration TERMINATION { $1 }

n2n_type:
| INT { Int }
| STRING { String }
| DOUBLE { Double }
| BOOL { Bool }
| GRAPH { Graph }
| NODE { Node }
| REL { Rel }
| LIST LT n2n_type GT { List($3)}
| VOID { Void }

function_declaration:
| FUNCTION ID LPAREN formal_parameters RPAREN ARROW
 n2n_type LBRACE statements RBRACE /* fn foo (bar:

```

```

Int) -> Bool { ... } */
{
 { fname = $2;
 formals = $4;
 body = List.rev $9;
 return_type = $7;
 }
}

formal_parameters:
| /* nothing */ { [] }
| formal_list { List.rev $1 }

formal_list:
| parameter { [$1] } /* foo:
 Int */
| formal_list COMMA parameter { $3 :: $1 } /* foo:
 Int, bar: String */

parameter:
| ID COLON n2n_type { Formal($3, $1) } /* foo: Int
 */

statements:
| /* nothing */ { [] }
| statements statement { $2 :: $1 }

statement:
| expr TERMINATION { Expr($1
) } /* 1 + 2 */
| RETURN expr TERMINATION { Return($2) }
 /* return 1 + 2 */
| LBRACE statements RBRACE { Block(List.rev
 $2) } /* { 1 + 2 \n 3 + 4 } */

```

```

| IF LPAREN expr RPAREN statement %prec NOELSE
 { If($3, $5, Block([])) }
| IF LPAREN expr RPAREN statement ELSE statement
 { If($3, $5, $7) }
| var_declaration TERMINATION
 { Var_Decl($1) } /*
 actor: Node, number: Int, graph_example: Graph */

expr:
| literal
 { Literal($1) } /* 42,
 "Jerry", 4.3, true */
| complex_literal
 { Complex($1) } /*
 constructor(literal, literal), { node rel node,
 node_literal rel_literal node_literal } */
| binary_operation
 { $1 } /* 4 + 3, "
 Johan" ^ "Mena" */
| graph_operation
 { $1 }
| graph_element_operation
 { $1 }
| unary_operation
 { $1 } /* -1 */
| ID
 { Id($1) } /* actor,
 number, graph_example */
| ID LPAREN actuals_opt RPAREN { Call($1, $3) } /*
 fuction_ID_String_param("Keanu") */
| built_in_function_call
 { Func($1) }
| LPAREN expr RPAREN
 { $2 } /* (4 + 6) */
| ID ACCESS ID
 { Access($1, $3) }

literal:
| INT_LITERAL
 { Int_Literal($1) } /*
 4, 3, 27 */
| STRING_LITERAL
 { String_Literal($1) }
 /* "Me", "You", "Bill Clinton" */
| DOUBLE_LITERAL
 { Double_Literal($1) }
 /* 4.2, 3.7, 7.4 */
| BOOL_LITERAL
 { Bool_Literal($1) }
 /* true, false */
| ANY
 { Any }

```

```

complex_literal :
 | ID LBRACKET literal_opt RBRACKET { Graph_Element
 ($1, $3) } /* actor("Keanu"), "true" */
 | LT graph_components GT { Graph_Literal
 ($2) } /* < graph_literals > */

graph_component :
 graph_type graph_type graph_type {
 Node_Rel_Node_Tup($1, $2, $3) } /* :: or commas?
 */

graph_components :
 | { [] }
 | graph_component_list { List.rev $1 }

graph_component_list :
 | graph_component { [$1] }
 | graph_component_list COMMA graph_component { $3 ::
 $1 }

graph_type :
 | ID { Graph_Id($1) }
 | complex_literal { Graph_Type($1) }

literal_opt :
 | { [] }
 | literal_list { List.rev $1 }

literal_list :
 | literal { [$1] }
 | literal_list COMMA literal { $3 :: $1 }

actuals_opt :
 | /* nothing */ { [] }
 | actuals_list { List.rev $1 }

actuals_list :
 | expr { [$1] }

```

```

| actuals_list COMMA expr { $3 :: $1 }

binary_operation :
| expr PLUS expr
 { Binop($1, Add, $3) }
| expr MINUS expr
 { Binop($1, Sub, $3) }
| expr TIMES expr
 { Binop($1, Mult, $3) }
| expr DIVIDE expr
 { Binop($1, Div, $3) }
| expr MOD expr
 { Binop($1, Mod, $3) }
| expr EQ expr
 { Binop($1, Equal, $3) }
| expr NEQ expr
 { Binop($1, Neq, $3) }
| expr LT expr
 { Binop($1, Less, $3) }
| expr LEQ expr
 { Binop($1, Leq, $3) }
| expr GT expr
 { Binop($1, Greater, $3) }
| expr GEQ expr
 { Binop($1, Geq, $3) }
| expr AND expr

```

```

 { Binop($1, And, $3) }
| expr OR expr

 { Binop($1, Or, $3) }
| expr CONCAT expr

 { Binop($1, Concat, $3) }

graph_operation :
| expr GRAPHINSERT LPAREN graph_component RPAREN
 { Grop($1, Graph_Insert, $4
) } /* ^+ */
| expr GRAPHREMOVE LPAREN graph_component RPAREN
 { Grop($1, Graph_Remove, $4
) } /* ^- */

graph_element_operation :
| expr DATAINSERT parameter
 { Geop($1,
 Data_Insert, $3) } /* [+] */
| expr DATAREMOVE parameter
 { Geop($1,
 Data_Remove, $3) } /* [-] */

unary_operation :
| NOT expr { Unop(Not, $2) }
| MINUS expr %prec NEG { Unop(Neg, $2) }

built_in_function_call :
| ID ACCESS find_many { Find_Many(
 $1, $3) } /* graph_example.find_many(...) */
| ID ACCESS map_function { Map($1, $3
) } /* graph_or_list_example.map(...) */
| ID ACCESS neighbors_function {
 Neighbors_Func($1, $3) } /* graph_example.
 neighbors(node_ID) */

map_function :

```

```

| MAP LPAREN ID IN LBRACE statements RBRACE RPAREN {
 Map_Func($3, $6) } /* map(node in {...}) */

neighbors_function:
| NEIGHBORS LPAREN ID RPAREN { $3 }

find_many:
| FINDMANY LPAREN complex_literal RPAREN
 { Find_Many_Node($3) } /* Find
 all nodes that match a literal, i.e. find_many(
 actor("Neo")) returns all nodes of type actor that
 have the name field equal to "Neo" */
| FINDMANY LPAREN graph_type COMMA graph_type RPAREN
 { Find_Many_Gen($3, $5) } /* Return what's
 missing, i.e. nodes pointed to, nodes pointed from
 , or rel between nodes */

```

## 8.3 AST

```

type op =
| Add
| Sub
| Mult
| Div
| Mod
| Equal
| Neq
| Less
| Leq
| Greater
| Geq
| And
| Or
| Concat

type grop =

```

```

| Graph_Insert
| Graph_Remove

type geop =
| Data_Insert
| Data_Remove

type uop =
| Not
| Neg

type n2n_type =
| Int
| String
| Bool
| Double
| Graph
| Node
| Rel
| List of n2n_type
| Void

type formal =
| Formal of n2n_type * string

type var_decl =
| Var of n2n_type * string
| Constructor of n2n_type * string * formal list
| Var_Decl_Assign of string * n2n_type * expr
| Access_Assign of expr * expr

and expr =
| Literal of literal
| Complex of complex_literal
| Id of string
| Binop of expr * op * expr
| Grop of expr * grop * graph_component
| Geop of expr * geop * formal

```



```

| Unop of uop * expr
| Access of string * string
| Call of string * expr list
| Func of built_in_function_call

and literal =
| Int_Literal of int
| Double_Literal of float
| String_Literal of string
| Bool_Literal of bool
| Any

and built_in_function_call =
| Find_Many of string * find_many
| Map of string * map_function
| Neighbors_Func of string * string

and complex_literal =
| Graph_Literal of graph_component list
| Graph_Element of string * literal list

and map_function =
| Map_Func of string * statement list

and find_many =
| Find_Many_Node of complex_literal
| Find_Many_Gen of graph_type * graph_type

and graph_component =
| Node_Rel_Node_Tup of graph_type * graph_type *
 graph_type

and graph_type =
| Graph_Id of string
| Graph_Type of complex_literal

and statement =
| Block of statement list

```

```

| Expr of expr
| Return of expr
| If of expr * statement * statement
| Var_Decl of var_decl

type func_decl = {
 fname : string;
 formals : formal list;
 body : statement list;
 return_type : n2n_type;
}

type program = var_decl list * func_decl list

let string_of_binop = function
| Add -> "+"
| Sub -> "-"
| Mult -> "*"
| Div -> "/"
| Mod -> "mod"
| Equal -> "=="
| Neq -> "!="
| Less -> "<"
| Leq -> "<="
| Greater -> ">"
| Geq -> ">="
| And -> "&&"
| Or -> "||"
| Concat -> "^"

let string_of_grop = function
| Graph_Insert -> "[+]"
| Graph_Remove -> "[-]"

let string_of_geop = function
| Data_Insert -> "^+"
| Data_Remove -> "^-"

```

```

let string_of_unop = function
 | Neg -> "-"
 | Not -> "!"

let rec string_of_n2n_type = function
 | Int -> "Int"
 | String -> "String"
 | Bool -> "Bool"
 | Double -> "Double"
 | Graph -> "Graph"
 | Node -> "Node"
 | Rel -> "Rel"
 | List(t) -> "List<" ^ string_of_n2n_type t ^ ">"
 | Void -> "Void"

let string_of_formal = function
 | Formal(the_type, id) -> id ^ " : " ^
 string_of_n2n_type the_type

let string_of_literal = function
 | Int_Literal(l) -> string_of_int l
 | Double_Literal(l) -> string_of_float l
 | String_Literal(l) -> "\"" ^ l ^ "\""
 | Bool_Literal(l) -> string_of_bool l
 | Any -> "_"

let rec string_of_expr = function
 | Literal(l) -> string_of_literal l
 | Complex(c) -> string_of_complex_literal c
 | Binop(e1, o, e2) ->
 string_of_expr e1 ^ " " ^
 string_of_binop o ^ " " ^
 string_of_expr e2
 | Grop(e1, grop, e2) ->
 string_of_expr e1 ^ " " ^
 string_of_grop grop ^ " " ^
 string_of_graph_component e2
 | Geop(e1, geop, e2) ->

```

```

 string_of_expr e1 ^ " " ^
 string_of_geop geop ^ " " ^
 string_of_formal e2
| Unop(o, e) ->
 string_of_unop o ^ " " ^
 string_of_expr e
| Id(s) -> s
| Access(id1, id2) -> id1 ^ "." ^ id2
| Call(func_id, actuals) ->
 func_id ^ "(" ^ String.concat ", " (List.map
 string_of_expr actuals) ^ ")"
| Func(k) -> string_of_built_in_fdecl k

and string_of_var_decl = function
| Var(the_type, id) -> id ^ " : " ^ string_of_n2n_type
 the_type
| Constructor(the_type, id, formal_list) ->
 id ^ " : " ^ string_of_n2n_type the_type ^ " = {
 " ^ String.concat ", " (List.map
 string_of_formal formal_list) ^ "}"
| Var_Decl_Assign(id, the_type, expr) ->
 id ^ " : " ^ string_of_n2n_type the_type ^ " = "
 ^ string_of_expr expr
| Access_Assign(e1, e2) ->
 string_of_expr e1 ^ " = " ^ string_of_expr e2

and string_of_complex_literal = function
| Graph_Literal(graph_type_l) -> "(" ^ String.concat
 ", " (List.map string_of_graph_component
 graph_type_l) ^ ")"
 (*Modification needed*)
| Graph_Element(id, el) ->
 id ^ "[" ^ String.concat ", " (List.map
 string_of_literal el) ^ "]"

and string_of_graph_component = function
| Node_Rel_Node_Tup(n1, r1, n2) ->

```

```

 string_of_graph_type n1 ^ " " ^
 string_of_graph_type r1 ^ " " ^
 string_of_graph_type n2

and string_of_builtin_fdecl = function
 | Find_Many(graph_id, s) -> graph_id ^ "." ^
 string_of_find_many s
 | Map(l, s) -> l ^ "." ^ string_of_map s
 | Neighbors_Func(l, node) ->
 l ^ ".neighbors (" ^ node ^ ")"

and string_of_find_many = function
 | Find_Many_Node(graph_type) ->
 "find_many (" ^ string_of_complex_literal
 graph_type ^ ")"
 | Find_Many_Gen(graph_type1, graph_type2) ->
 "find_many (" ^ string_of_graph_type graph_type1
 ^ ", " ^ string_of_graph_type graph_type2 ^ ")"

and string_of_map = function
 | Map_Func(element_id, s) ->
 "map (" ^ element_id ^ " in " ^ " { " ^ String.
 concat "\n" (List.map string_of_statement s) ^
 " })"

and string_of_graph_type = function
 | Graph_Id(id) -> id
 | Graph_Type(graph_element) ->
 string_of_complex_literal graph_element

and string_of_statement = function
 Block(l) -> "{" ^ String.concat "\n" (List.map
 string_of_statement l) ^ "}"
 | Expr(l) -> string_of_expr l ^ "\n"
 | Return(l) -> "return" ^ string_of_expr l ^ "\n"
 | If(e, l, p) ->
 (match p with

```

```

 Block([]) -> "if (" ^ string_of_expr e ^ ")\n"
 ^ string_of_statement l
 | _ -> "if (" ^ string_of_expr e ^ ")\n" ^
 string_of_statement l ^ "else\n" ^
 string_of_statement p)
| Var_Decl(v) -> string_of_var_decl v

(*****

let string_of_func_decl fdecl=
 "fn" ^ (fdecl.fname) ^ "(" ^ String.concat ", " (List
 .map string_of_formal fdecl.formals)
 ^ ")" -> " ^ (string_of_n2n_type fdecl.return_type) ^
 "{" ^ String.concat "\n" (List.map
 string_of_statement fdecl.body) ^ "}"

let string_of_program (vars, funcs) =
 String.concat "\n" (List.map string_of_var_decl (List
 .rev vars)) ^ (if (List.length vars) > 0 then "\n"
 else "") ^
 String.concat "\n" (List.map string_of_func_decl (
 List.rev funcs))

```

## 8.4 Semantic Check

```

open Ast
open Sast

exception Error of string;;

type environment = {
 functions: func_decl list;
 scope: string;
 node_types: (string * formal list) list;
 rel_types: (string * formal list) list;

```

```

locals: var_scope;
globals: var_scope;
has_return: bool;
return_val: expr;
return_type: n2n_type;
}

and var_scope = {
prims: (string * n2n_type * expr) list;
nodes: (string * string * (string * n2n_type * expr)
 list) list; (* Form (id, node_type, field storage
 list) *)
rels: (string * string * (string * n2n_type * expr)
 list) list; (* Form (id, node_type, field storage
 list) *)
graphs: (string * graph_component list) list;
lists: (string * n2n_type * expr list) list (*Form (
 list_id, type, list contents) *)
}

let beginning_scope = { prims = []; nodes = []; rels =
 []; graphs = []; lists = [] }

let beginning_environment = { functions = []; globals =
 beginning_scope; locals = beginning_scope; scope =
 "global"; node_types = []; rel_types = [];
 has_return = false; return_val = Id("None");
 return_type = Void}

(*For debugging *)

let rec print_type ty =
match ty with
Node -> "Node"
| Rel -> "Rel"
| Graph -> "Graph"
| String -> "String"
| Int -> "Int"

```

```

| Double -> "Double"
| Bool -> "Bool"
| Void -> "Void"
| List(t) -> "List<" ^ print_type t ^ ">"

(* End debugging *)

let get_id_from_expr ex =
match ex with
Id(v) -> v
| _ -> raise(Error("Trying to get id from a non-id
expression\n"))

let add_tuple_to_list env t =
match t with
(id,ty,fl) ->
let (list_to_update, location) = if List.exists (fun (
fid,_,_) -> fid=id) env.locals.nodes then (env.
locals.nodes, "nl")
else if List.exists (fun (fid,_,_) -> fid=id) env.
locals.rels then (env.locals.rels, "rl")
else if List.exists (fun (fid,_,_) -> fid=id) env.
globals.nodes then (env.globals.nodes, "ng")
else (env.globals.rels, "rg") in
let (new_list, location) =
let updated_list = List.fold_left (fun l (vid,vty,vfl)
-> if vid=id then (id,ty,fl)::l else (vid,vty,vfl)::
l) [] list_to_update in
(List.rev updated_list, location) in
match location with
"nl" -> let new_vt = {env.locals with nodes = new_list}
in
{env with locals = new_vt}
| "rl" -> let new_vt = {env.locals with rels = new_list
} in
{env with locals = new_vt}
| "ng" -> let new_vt = {env.globals with nodes =
new_list} in

```



```

 {env with globals = new_vt}
| "rg" -> let new_vt = {env.globals with rels =
 new_list} in
 {env with globals = new_vt}
| _ -> raise(Error("??"))

let check_arithmetic_binary_op t1 t2 =
match (t1, t2) with
| (Int, Int) -> Int
| (Int, Double) -> Double
| (Double, Int) -> Double
| (Double, Double) -> Double
| (_,_) -> raise(Error("Binary operation fails , wrong
 element type"))

let check_equality t1 t2 =
prerr_string(print_type t1 ^ " " ^ print_type t2 ^ "\n");
if t1 = t2 then Bool else
match (t1, t2) with
| (Int, Double) -> Bool
| (Double, Int) -> Bool
| (_, _) -> raise(Error("Equality operation fails ,
 arguments not same type"))

let check_logic t1 t2 =
match(t1, t2) with
| (Int, Int) -> Bool
| (Int, Double) -> Bool
| (Double, Int) -> Bool
| (Double, Double) -> Bool
| (String, String) -> Bool
| (_,_) -> raise(Error("Logical operation fails ,
 arguments not of correct types"))

let get_literal_type l = match l with
Int_Literal(i) -> Int
| Double_Literal(d) -> Double
| String_Literal(s) -> String

```

```

| Bool_Literal(b) -> Bool
| _ -> raise(Error("Should not be using Any here"))

let get_type_from_id var_table id =
if List.exists (fun (lid ,-, -) -> lid=id) var_table.
 lists then
 (let (_, ty, _) = List.find(fun (lid ,-, -) -> lid=id)
 var_table.lists in List(ty))
else if List.exists (fun (nid ,-, -) -> nid=id) var_table
 .nodes then Node
else if List.exists (fun (rid ,-, -) -> rid=id) var_table
 .rels then Rel
else if List.exists (fun (gid, _) -> gid=id) var_table.
 graphs then Graph
else let (_, ty, _) = try List.find (fun (vid, -, -) ->
 vid=id) var_table.prims with
 Not_found -> raise Not_found in ty

let check_for_var_existence var_table id =
(List.exists (fun (gid, _) -> gid=id) var_table.graphs
|| List.exists (fun (rid ,-, -) -> rid=id) var_table.rels
|| List.exists (fun (nid ,-, -) -> nid=id) var_table.nodes
|| List.exists (fun (vid ,-, -) -> vid=id) var_table.prims
|| List.exists (fun (lid ,-, -) -> lid=id) var_table.lists
)

let check_if_id_is_node env id =
if List.exists (fun (nid, -, -) -> prerr_string("
 Looking for: "^id^ " but finding " ^nid^ ".\n"); nid
 = id) env.locals.nodes then true
else if List.exists (fun (nid, -, -) -> prerr_string("
 Looking for: "^id^ " but finding " ^nid^ ".\n"); nid
 = id) env.globals.nodes then true
else false

let set_default_val ty =
match ty with
Int -> Literal(Int_Literal(0))

```

```

| Double -> Literal(Double_Literal(0.0))
| Bool -> Literal(Bool_Literal(false))
| String -> Literal(String_Literal(""))
| Graph -> Complex(Graph_Literal([]))
| Node | Rel -> Complex(Graph_Element("", []))
| _ -> raise(Error("Not a primitive type, YOU FOOL!"))

let get_new_env env func =
let new_locals = List.fold_left (fun l f -> let (t, id)
 = (match f with
Formal(ty, vid) -> (ty, vid)) in
if check_for_var_existence env.locals id then raise(
 Error("Variable: " ^ id ^ " already exists in local
scope"))
else (match t with
Node -> let new_nodes = (id, id, []) :: l.nodes in {l
 with nodes = new_nodes}
| Rel -> let new_rels = (id, id, []) :: l.rels in {l
 with rels = new_rels}
| Graph -> let new_graphs = (id, []) :: l.graphs in {l
 with graphs = new_graphs}
| List(ty) -> let new_lists = (id, ty, []) :: l.lists
 in {l with lists = new_lists}
| Void -> raise(Error("Can't have a variable of type
Void"))
| _ -> let new_prims = (id, t, set_default_val t) :: l
 prims in {l with prims = new_prims}) env.locals
 func.formals and
new_functions = func :: env.functions in
{env with functions = new_functions; locals =
 new_locals; scope = func.fname; return_type = func.
return_type}

let update_prim_table var_table id ty v =
let does_exist = check_for_var_existence var_table id
 in
match does_exist with

```

```

true -> raise (Error (" Variable to declare : " ^ id ^ " "
 already exists "))
| false ->
let new_prims = (id , ty , v) :: var_table . prims in
{ var_table with prims = new_prims }

let rec gen_tuple_list fl l tl =
match fl , l with
[] , _ [] -> List . rev tl
| _h1 :: t1 , _h2 :: t2 ->
(match _h1 , _h2 with
Formal (ty , id) , _lit -> let new_tl = (id , ty , Literal (
 lit)) :: tl in
 gen_tuple_list t1 t2 new_tl
)
| _ -> raise (Error (" You're such a failure "))

let update_node_or_rel_table env var_table id idt ty ex
=
let does_exist = check_for_var_existence var_table id
in
match does_exist with
true -> raise (Error (" Variable to declare already exists
 "))
| false ->
let l = (match ex with
Complex (Graph_Element (s , ll)) -> ll
| _ -> raise (Error (" Just cry . Stop what you're doing
 and cry "))) and
forml = (match ty with
Node -> let _ (, fl) = List . find (fun (fid , _) ->
 prerr_string (" Looking for " ^ idt ^ " constructor ,
 finding : " ^ fid ^ "\n"));
 gen_tuple_list fl (id = idt) env . node_types
in fl
| Rel -> let _ (, fl) = List . find (fun (fid , _) ->
 prerr_string (" Looking for " ^ idt ^ " constructor ,
 finding : " ^ fid ^ "\n"));

```

```

..... fid=idt) env.rel_types
 in fl
| _-> raise (Error("Wow. Do you even go here?")) in
let t1 = gen_tuple_list form1 l [] in
(match ty with
Node-> let new_nodes = (id, idt, t1) :: var_table.nodes
 in
..... { var_table with nodes = new_nodes }
| Rel-> let new_rels = (id, idt, t1) :: var_table.rels
 in
..... { var_table with rels = new_rels }
| _-> raise (Error(" You called the wrong var assign
 method you fool!")))

let update_graph_table var_table id v =
let does_exist = check_for_var_existence var_table id
 in
match does_exist with
true-> raise (Error(" Variable to declare already exists
 "))
| false->
let va = (match v with
| Complex(Graph_Literal(1))-> 1
| Func(Map(s, e1))-> let (_, e1) = List.find (fun (gid, _)
)-> gid=s) var_table.graphs in e1
| Call(id, _)-> []);
| _-> raise (Error(" Calling the wrong function , you
 fool!")) in
let new_graphs = (id, va) :: var_table.graphs in
..... { var_table with graphs = new_graphs }

let update_list_table var_table id ty v =
let does_exist = check_for_var_existence var_table id
 in
match does_exist with
true-> raise (Error(" Variable to declare already exists
 "))
| false->

```

```

let new_lists = (id, ty, v) :: var_table.lists in
{ var_table with lists = new_lists }

let rec check_expr env expr =
match_expr with
| Literal(l) -> get_literal_type l
| Complex(c) -> (match c with
Graph_Literal(nrn_list) -> check_nrn_list env nrn_list
| Graph_Element(id, lit_list) ->
 check_node_or_rel_literal env id lit_list)
| Id(v) -> prerr_string(" check_expr: " ^ v ^ " id
 called\n");
 (try get_type_from_id env.locals v with
 Not_found -> try get_type_from_id env.globals v with
 Not_found -> raise (Error(" Id does not appear in program
 "))))
| Unop(u, e) -> (match u with
 Not -> if check_expr env e = Bool then Bool else raise (
 Error(" Using NOT on a non-boolean expr"))
 | Neg -> if check_expr env e = Double then Double
 else if check_expr env e = Int then Int
 else raise (Error(" Using a neg on a non-int or float
 expr"))
 | Binop(e1, op, e2) ->
 let t1 = (match e1 with
 Access(_, _) -> check_expr env e2
 | _ -> check_expr env e1) and t2 = check_expr env e2 in
 let binop_t = (match op with
 Add -> check_arithmetic_binary_op t1 t2
 | Sub -> check_arithmetic_binary_op t1 t2
 | Mult -> check_arithmetic_binary_op t1 t2
 | Div -> check_arithmetic_binary_op t1 t2
 | Mod -> if (t1, t2) = (Int, Int) then Int else raise (
 Error(" Using MOD on a non-integer expr"))
 | Equal -> check_equality t1 t2
 | Neq -> check_equality t1 t2
 | Less -> check_logic t1 t2
 | Leq -> check_logic t1 t2

```



```

| _ -> raise (Error (" Can only insert field into a Node
 or Rel"))
| Access (idl, idr) ->
prerr_string (idl ^ "." ^ idr ^ " called\n");
let t = (try get_type_from_id env.locals idl with
Not_found -> try get_type_from_id env.globals idl with
Not_found -> raise (Error (" Can't find left identifier"))
) in
(match t with
| Node -> t
| Rel -> t
| _ -> raise (Error (" Trying to access something that is
 not a node or rel")))
| Call (" print", el) -> prerr_string (" Print function is
 being called\n"); List.iter (fun e -> ignore (
 check_expr env e)) el; Void
| Call (id, el) -> let func = (try List.find (fun f ->
 prerr_string (" Looking for function: " ^ id ^ " but
 finding function: " ^ f.fname ^ "\n"));

```

Not\_found -> raise (Error (" Function definition **not** found
 ))) in
(tr try List.iter2 (fun e f -> let ty = (match f with



|            |             |
|------------|-------------|
|            | Formal      |
|            | (           |
|            | t           |
|            | ,           |
|            | -           |
|            | )           |
|            | ->          |
|            | t           |
|            | )           |
|            | in          |
| let t =    |             |
| check_expr |             |
| env        |             |
| e          |             |
| in         | <b>if</b>   |
|            | t           |
|            | ◇           |
|            | ty          |
|            | <b>then</b> |
|            | raise       |
|            | (           |
|            | Error       |
|            | ("          |
|            | Argument    |
|            | does        |

```

Invalid_argument s -> raise(Error("Entered the wrong
 number of arguments into function")); func.
 return_type
| Func(fname) -> (match fname with
Find_Many(id, e1) -> prerr_string("Checking find_many on
 "^id^ ".\n");
if List.exists (fun (gid, _) -> gid = id) env.locals.
 graphs then
 let lt = check_find_many_arguments env e1 in
 List(lt)
else if List.exists (fun (gid, _) -> gid = id) env.
 globals.graphs then
 let lt = check_find_many_arguments env e1 in
 List(lt)
else raise(Error("Could not find List or Graph ID: " ^
 id ^ " to run find_many on"))
| Neighbors_Func(id, nid) ->

```

```

not
match
expected
argument
type
")
)
)
el
func
.
formals
with

```

```

if List.exists (fun (gid, _) => gid = id) env.locals.
 graphs then
 if check_if_id_is_node env nid then List(Node)
 else raise(Error("Locals: Argument to
 neighbors must be a node id"))
else if List.exists (fun (gid, _) => gid = id) env.
 globals.graphs then
 if check_if_id_is_node env nid then List(Node)
 else raise(Error("Globals: Argument to
 neighbors must be a node id"))
else raise(Error("Could not find List or Graph ID: " ^
 id ^ " to run neighbors on"))
| Map(id, e2) =>
prerr_string("Checking collection id: " ^ id ^ " for
 map\n");
let ty = try get_type_from_id env.locals id with
 Not_found => try get_type_from_id env.globals
 id with
 Not_found => raise(Error("Could not find List
 ID: " ^ id ^ " to run map on")) in
prerr_string(id ^ " is a " ^ print_type ty ^
 "\n");
match ty with
List(t) => List(t)
| Graph => Graph
| _ => raise(Error("Trying to run map on non-
 list or graph. Id: " ^ id ^ "\n"))

and check_map_func env ty map_func =
match map_func with
Map_Func(id, stmt_list) =>
let new_locals = (match ty with
Graph => let new_nodes = (id, id, []) :: env.locals.
 nodes in
{env.locals with nodes = new_nodes}
| List(t) => (match t with
Node => let new_nodes = (id, id, []) :: env.locals.
 nodes in

```

```

{env.locals with nodes = new_nodes}
| Rel -> let new_rels = (id, id, []) :: env.locals.rels
 in
{env.locals with rels = new_rels}
| Graph -> let new_graphs = (id, []) :: env.locals.
 graphs in
{env.locals with graphs = new_graphs}
| List(t2) -> raise(Error("Cannot have List of Lists"))
| Void -> raise(Error("Cannot have List of Voids"))
| - -> let new_prims = (id, t, set_default_val t) ::
 env.locals.prims in
{env.locals with prims = new_prims})
| - -> raise(Error("Cannot have a map function operate
 on a non-collection")) in
let new_env = {env with locals = new_locals} in
let (checked_stmts, up_env) = get_checked_statements
 new_env stmt_list [] in
checked_stmts

and check_find_many_arguments env e =
match e with
Find_Many_Node(complex) -> prerr_string("Find_many is
 the general type\n"); (match complex with
Graph_Element(s, ll) -> let t =
 check_node_or_rel_literal_matching env s ll in
if t = Node then t else raise(Error("Find_many_node
 does not have node as argument"))
| Graph_Literal(gcl) -> raise(Error("Find_many_node
 does not have node as argument!"))
| Find_Many_Gen(gt1, gt2) -> let t1 = check_graph_type
 env gt1 and t2 = check_graph_type env gt2 in
(match (t1, t2) with
(Node, Rel) -> prerr_string("Find_many is returning a
 node list\n"); Node
| (Rel, Node) -> prerr_string("Find_many is returning a
 node list\n"); Node
| (Node, Node) -> prerr_string("Find_many is returning
 a rel list\n"); Rel

```

```

| (Rel, Rel) -> raise(Error("Cannot have two rel
 arguments in Find_Many"))
| (_,_) -> raise(Error("Must have (Node, Node), (Rel,
 Node), or (Node, Rel) as arguments to find_many"))

and is_node env id =
let isNode = List.exists (fun (fid, _) -> fid = id) env
 .node_types in
isNode

and is_rel env id =
let isRel = List.exists (fun (fid, _) -> fid = id) env.
 rel_types in
isRel

and check_node_literal env id lit_list =
let (_, l) = List.find (fun (fid, _) -> fid = id) env.
 node_types in
(try List.iter2 (fun lit f -> let t2 = (match f with
Formal(ty, _) -> ty
) in
if (lit = Any) then raise(Error("Cannot have a complex
 type with any value. Can only use in find_many
 matching.)) else
let t1 = get_literal_type lit in
if t1 <> t2 then raise(Error("Type mismatch between
 arguments and expected type for given node object.")
)) lit_list l with
Invalid_argument s -> raise(Error("Lists have unequal
 sizes. Check number of literals in your assignment.\
n" ^
"Constructor list size: " ^ string_of_int (List.length
l) ^ "\n" ^ "Literal List size: " ^ string_of_int(
List.length lit_list) ^ "\n")); Node

and check_rel_literal env id lit_list =
let (_, l) = List.find (fun (fid, _) -> fid = id) env.
 rel_types in

```

```

(tr try List.iter2 (fun lit f -> let t2 = (match f with
Formal(ty, -) -> ty
) in
if(lit = Any) then raise(Error("Cannot have a complex
type with any value. Can only use in find_many
matching.")) else
let t1 = get_literal_type lit in
if t1 <> t2 then raise(Error("Type mismatch between
arguments and expected type for given rel object."))
) lit_list l with
Invalid_argument s -> raise(Error("Lists have unequal
sizes. Check number of literals in your assignment.\
n" ^
"Constructor list size: " ^ string_of_int(List.length l
) ^ "\n" ^ "Literal List size: " ^ string_of_int(
List.length lit_list) ^ "\n")); Rel

and check_node_or_rel_literal env id lit_list =
if is_node env id then check_node_literal env id
lit_list
else if is_rel env id then check_rel_literal env id
lit_list
else raise(Error("Could not find constructor for your
node or rel"))

and check_node_literal_matching env id lit_list =
let (_, l) = List.find (fun (fid, _) -> fid = id) env.
node_types in
(tr try List.iter2 (fun lit f -> let t2 = (match f with
Formal(ty, -) -> ty
) in
let t1 = if(lit = Any) then t2 else get_literal_type
lit in
if t1 <> t2 then raise(Error("Type mismatch between
arguments and expected type for given node object."))
)) lit_list l with
Invalid_argument s -> raise(Error("Lists have unequal
sizes. Check number of literals in your assignment.\

```

```

n" ^
"Constructor list size: " ^ string_of_int (List.length
 l) ^ "\n" ^ "Literal List size: " ^ string_of_int(
 List.length lit_list) ^ "\n")); Node

and check_rel_literal_matching env id lit_list =
let (_, l) = List.find (fun (fid, _) -> fid = id) env.
 rel_types in
(trial List.iter2 (fun lit f -> let t2 = (match f with
Formal(ty, -) -> ty
) in
let t1 = if (lit = Any) then t2 else get_literal_type
 lit in
if t1 <> t2 then raise(Error("Type mismatch between
 arguments and expected type for given rel object.))
) lit_list l with
Invalid_argument s -> raise(Error("Lists have unequal
 sizes. Check number of literals in your assignment.\
n" ^
"Constructor list size: " ^ string_of_int(List.length l
) ^ "\n" ^ "Literal List size: " ^ string_of_int(
 List.length lit_list) ^ "\n")); Rel

and check_node_or_rel_literal_matching env id lit_list
=
if is_node env id then check_node_literal_matching env
 id lit_list
else if is_rel env id then check_rel_literal_matching
 env id lit_list
else raise(Error("Could not find constructor for your
 node or rel"))

and check_graph_ID env id =
if List.exists (fun (nid, -, -) -> nid=id) env.locals.
 nodes then Node
else if List.exists (fun (rid, -, -) -> rid=id) env.
 locals.rels then Rel

```

```

else if List.exists (fun (rid , _ , _) -> rid = id) env .
 globals.rels then Rel
else if List.exists (fun (nid , _ , _) -> nid=id) env .
 globals.nodes then Node
else raise(Error("Id (" ^ id ^") doesn't exist"))

and check_graph_type env gt =
let t = (match gt with
Graph_Id (gid) -> check_graph_ID env gid
| Graph_Type (complex) -> (match complex with
| Graph_Element (id , lit_list) ->
 check_node_or_rel_literal env id lit_list
| _ -> raise (Error("There is no spoon")))) in t

and check_nrn_expr env n1 r n2 =
let t1 = check_graph_type env n1 and t2 =
 check_graph_type env n2 and tr = check_graph_type env r in
match (t1 , tr , t2) with
(Node , Rel , Node) -> true
| _ -> false

and check_nrn_list env nrn_list =
List.iter (fun nrn_expr ->
match nrn_expr with
Node_Rel_Node_Tup (n1 , r , n2) -> if check_nrn_expr env n1 r n2 != true then
 raise (Error("Combination is not a Node-Rel-Node combination"))) nrn_list ; Graph

and get_type_from_constructor env id =
if List.exists (fun (fid , _) -> fid=id) env . node_types
 then Node
else Rel

and get_field_list env id ll =

```



```

let (l, l) ← try List.find (fun (fid, _) → fid=id) env .
 node_types with
Not_found → List.find (fun (fid, _) → fid=id) env .
 rel_types in
get_field_lists l ll []

and get_field_lists l ll fl ←
match l, ll with
| [], _ [] → List.rev fl
| _headl :: tail, _headll :: tailll →
(match _headl with
Formal (ty, _head) → let sast_literal ← (match _headll
 with
Int_Literal (i) → SInt_Literal (i)
| _Double_Literal (d) → SDouble_Literal (d)
| _Bool_Literal (b) → SBool_Literal (b)
| _String_Literal (s) → SString_Literal (s)
| _Any → SAny) in
let new_fl ← (head, ty, sast_literal) :: fl in
get_field_lists tail tailll new_fl)
| _ → raise (Error ("Lit list does not match constructor
 list"))

and get_scomplex env complex ←
(match complex with
Graph_Literal (gcl) → SGraph_Literal (get_sgc_list env
 gcl [])
| _Graph_Element (str, ll) → SGraph_Element ((
 get_type_from_constructor env str, str),
 get_field_list env str ll))

and get_sgc_list env gcl sgcl ←
match gcl with
[] → List.rev sgcl
| _head :: tail → let new_gcl ← get_sgc env head ::
 sgcl in
get_sgc_list env tail new_gcl

```

```

and_get_sgc_env_gc =
 (match_gc_with
 Node_Rel_Node_Tup(n1, r, n2) -> SNode_Rel_Node_tup(
 get_sgraph_type_env_n1, get_sgraph_type_env_r,
 get_sgraph_type_env_n2))

and_get_sgraph_type_env_gt =
 (match_gt_with
 Graph_Id(s) -> SGraph_Id(s)
 | Graph_Type(complex) -> SGraph_type(get_scomplex_env_complex))

and_get_sformal_form =
 (match_form_with
 Formal(ty, s) -> SFormal(ty, s))

and_get_sformal_list_fl_sfl =
 match_fl_with
 [] -> List.rev_sfl
 | head::tail -> let new_sfl = (get_sformal_head) ::
 sfl in
 get_sformal_list_tail_new_sfl

and_get_sfm_env_fm =
 match_fm_with
 Find_Many_Node(complex) -> SFind_Many_Node(get_scomplex_env_complex)
 | Find_Many_Gen(gt1, gt2) -> SFind_Many_Gen(
 get_sgraph_type_env_gt1, get_sgraph_type_env_gt2)

and_get_smap_env_ty_mf =
 match_mf_with
 Map_Func(s, sl) -> SMap_Func(s, check_map_func_env_ty_mf)

and_get_sbuilt_in_function_call_env_f =
 match_f_with
 Find_Many(s, fm) -> SFindMany(s, get_sfm_env_fm)

```

```

| _Map(s, mf) → SMap(s, (check_expr_env (Id(s)),
 get_smap_env (check_expr_env (Id(s))) mf)
| _Neighbors_Func(s1, s2) → SNeighbors_Func(s1, s2)

and_get_sexpr_env_ex := match_ex_with
Literal(l) → (match_l_with
Int_Literal(i) → SLiteral(SInt_Literal(i), Int)
| _Double_Literal(d) → SLiteral(SDouble_Literal(d),
 Double)
| _String_Literal(s) → SLiteral(SString_Literal(s),
 String)
| _Bool_Literal(b) → SLiteral(SBool_Literal(b), Bool)
| _Any → SLiteral(SAny, String))
| _Id(v) → SId(v, check_expr_env_ex)
| _Unop(u, e) → SUNop(u, get_sexpr_env_e, check_expr_env_ex)
| _Binop(e1, op, e2) → SBinop(get_sexpr_env_e1, op,
 get_sexpr_env_e2, check_expr_env_ex)
| _Grop(e, grop, gc) → SGrop(get_sexpr_env_e, grop,
 get_sgc_env_gc, check_expr_env_ex)
| _Geop(e, geop, form) → SGeop(get_sexpr_env_e, geop,
 get_sformal_form, check_expr_env_ex)
| _Access(str, str2) → SAccess(str, str2, check_expr_env_ex)
| _Call(str, el) → SCall(str, List.map(fun e →
 get_sexpr_env_e) el, check_expr_env_ex)
| _Func(f) → SFunc(get_sbuilt_in_function_call_env_f,
 check_expr_env_ex)
| _Complex(comp) → SComplex(get_scomplex_env_comp,
 check_expr_env_ex)

and_resolve_envs old_env new_env :=
let new_prims := List.map(fun (id, ty, e) → let v := (
 try List.find(fun (vid, _, _) → vid=id) new_env .
 locals . prims with
Not_found → (id, ty, e)) in v) old_env . locals . prims
and new_nodes := List.map(fun (id, idt, l) → let v := (
 try List.find(fun (vid, _, _) → vid=id) new_env .

```

```

 locals . nodes . with
Not_found -> (id , _id , _l) . in _v) . old_env . locals . nodes
and _new_rels := List . map (fun (id , _idt , _l) -> let _v := (
 try List . find (fun (vid , _ , _) -> vid = id) . new_env .
 locals . rels . with
Not_found -> (id , _idt , _l) . in _v) . old_env . locals . rels
and _new_graphs := List . map (fun (id , _gcl) -> let _v := (
 try List . find (fun (gid , _ , _) -> gid = id) . new_env .
 locals . graphs . with
Not_found -> (id , _gcl) . in _v) . old_env . locals . graphs
and _new_lists := List . map (fun (id , _ty , _el) -> let _v := (
 try List . find (fun (lid , _ , _) -> lid = id) . new_env .
 locals . lists . with
Not_found -> (id , _ty , _el) . in _v) . old_env . locals . lists .
in
{new_env . with . locals := {prims := new_prims ; nodes :=
 new_nodes ; rels := new_rels ; graphs := new_graphs ;
 lists := new_lists }}

and _check_stmt _env _stmt :=
prerr_string (" Calling _check_stmt \n") ; _match_stmt . with
| _Block (stmt_list) ->
prerr_string (" Calling _Block . from _check_stmt \n") ;
let _new_env := _env . in
let (checked_stmts , _up_env) := List . fold_left (fun (l , _e
) _s -> let (checked_statment , _up_e) := _check_stmt _e _s
 in
.....
checked_statment _:: _l , _up_e) ([] , _env) . stmt_list . in
let _resolved_env := _resolve_envs _new_env _up_env . in
(SBlock (List . rev . checked_stmts) , _resolved_env)

| _Expr (e) ->
prerr_string (" Calling _expression . from _check_stmt ") ;
let _new_env := (match _e . with
Geop (e1 , _geop , _formal) ->
let _norid := _get_id_from_expr _e1 . in

```

```

let (id, t, fl) ← try List.find (fun (lid, _, _) → lid =
 norid) env.locals.nodes with
Not_found → try List.find (fun (lid, _, _) → lid = norid)
 env.locals.rels with
Not_found → try List.find (fun (lid, _, _) → lid = norid)
 env.globals.nodes with
Not_found → try List.find (fun (lid, _, _) → lid = norid)
 env.globals.rels with
Not_found → raise (Error(" Id: " ^ norid ^ " not found as
 a node or rel.\n")) in
let (id, ut, new_field_list) ← (match formal with
 Formal(ty, vid) →
 (match geop with
 Data_Insert → if List.exists (fun (
 vcid, _, _) → vcid = vid) fl then raise (Error(" Your
 field: " ^ vid ^ " already exists in node: " ^ id ^ ".\n
 "))
 else (id, ut, (vid, vty, set_default_val
 ty) :: fl)
 | Data_Remove → if List.exists (fun (
 vcid, t, _) → (vcid = vid && t = ty)) fl then
 let new_fl ← List.fold_left (
 fun l (vcid, vcty, vcfl) → if (vcid = vid) then
 else (vcid, vcty, vcfl) :: l) [] fl in
 (id, ut, new_fl) else raise (
 Error(" Field to remove does not exist"))
)
) in add_tuple_to_list env (id, id, new_field_list)
| _ → env) in (SExp(get_sexpr env e), new_env)

| Return(e) →
prerr_string(" Return from check_stmt");
let t1 ← check_expr env e in
(if not ((t1 = env.return_type)) then
raise (Error(" Incompatible Return Type")));

```

```

let new_env = {env with has_return = true; return_type =
 t1; return_val = e} in
(SReturn(get_sexpr env e), new_env)

| If(e, s1, s2) ->
prerr_string(" Calling If from check_stmt\n");
let t1 = check_expr env e in
(if not(t1=Bool) then
raise(Error(" If statement must be a boolean")));
let (st1, new_env) = check_stmt env s1 in
let (st2, new_env2) = check_stmt new_env s2 in
(SIf((get_sexpr env e), st1, st2), new_env2)

| Var_Decl(decl) ->
prerr_string(" Calling Var_decl from check_stmt\n");
let (checked_stmt, up_env) =
(match decl with
Var(ty, id) -> prerr_string(" Local_Var: Checking " ^ id
 ^ "\n");
let new_table = (match ty with
Node | Rel -> update_node_or_rel_table env env.locals id
 id ty (set_default_val ty)
| Graph -> update_graph_table env.locals id (
 set_default_val ty)
| List(t) -> prerr_string(" Var: Id: " ^ id ^ " is a
 List\n");
| _ -> raise(Error(" Can't have List of
 Lists! THAT'S INSANE")))
| _ -> update_list_table env.locals id t [])
| Void -> raise(Error(" Can't declare void"))
| _ -> update_prim_table env.locals id ty (
 set_default_val ty)) in
let new_env = {env with locals = new_table} in
(SVar(ty, id), new_env)
| Var_Decl_Assign(id, ty, e) -> prerr_string("
 Var_Decl_Assign: Checking " ^ id ^ "\n");
let t_ex = check_expr env e in

```

```

if (t_ex = ty) then
 let new_table = (match ty with
 Node | Rel-> let idt = (match e with
 Complex(Graph_Element(s, _)) -> s
 | _ -> raise(Error("Trying to assign
 non Node or Rel to non Node or Rel"))
)) in update_node_or_rel_table env
 env.locals id idt ty e
 | Graph -> update_graph_table env.locals id e
 | List(t) -> update_list_table env.locals id t
 []
 | Void -> raise(Error("Can't declare void"))
 | _ -> update_prim_table env.locals id ty e) in
 let new_env = {env with locals = new_table} in
 (SVar_Decl_Assign(id, ty, get_sexpr env e),
 new_env)
else
 raise(Error("Type mismatch in local variable
 assignment"))
| Access_Assign(e1, e2) ->
prerr_string(" Access_Assign being called from check_stmt
 \n");
let tl = check_expr env e1 and tr = check_expr env e1
 in
prerr_string(" tl = " ^ print_type tl ^ " tr = " ^
 print_type tr ^ ".\n");
if (tl = tr) then
 (SAccess_Assign(get_sexpr env e1, get_sexpr env
 e2), env)
else
 raise(Error("Type mismatch in assignment!"))
| Constructor(_,_,_) -> raise(Error("Can't declare
 constructor locally")) in (SVar_Decl(check_stmt),
 up_env)

and get_checked_statements env stmts checked_statments
 =
match stmts with

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| stmt :: tail ->
let (checked_statement, new_env) = check_stmt env stmt
 in
get_checked_statements new_env tail (checked_statement
 :: checked_statments)
| [] -> (List.rev checked_statments, env)

let check_function env func =
prerr_string(" Starting to check function: " ^ func.
 fname ^ ".\n");
let (sfstatements, up_env) = get_checked_statements (
 get_new_env env func) func.body [] in
({sfname = func.fname; sformals = get_sformal_list func
 .formals []; sbody = sfstatements; sreturn_type =
 func.return_type}, {up_env with locals = env.locals
 })

let rec check_functions env funcs checked_funcs =
let checked_functions =
(match funcs with
func :: tail ->
let (checked_func, up_env) = check_function env func in
check_functions up_env tail (checked_func ::
 checked_funcs)
| [] -> checked_funcs) in
checked_functions

let check_global env var =
let (checked_global, up_env) =
(match var with
Var(ty, id) -> prerr_string(" Global_Var: Checking " ^
 id ^ ".\n");
let new_table = (match ty with
Node | Rel-> update_node_or_rel_table env env.globals
 id id ty (set_default_val ty)
| Graph -> update_graph_table env.globals id (
 set_default_val ty)
| List(t) -> update_list_table env.globals id t [])

```



```

| Void -> raise (Error ("Can't declare void"))
| _ -> update_prim_table env.globals id ty (
 set_default_val ty) in
let new_env = {env with globals = new_table} in
(SVar (ty, id), new_env)
| SVar_Decl_Assign (id, ty, e) -> prerr_string ("
 Var_Decl_Assign: Checking " ^ id ^ "\n");
let t_ex = check_expr env e in
if (t_ex = ty) then
 let new_table = (match ty with
 Node | Rel -> let idt = (match e with
 Complex (Graph_Element (s, _)) -> s
 | _ -> raise (Error ("Trying to assign
 non_Node_or_Rel to non_Node_or_Rel"))) in
 update_node_or_rel_table env env.globals id idt ty e
 | Graph -> update_graph_table env.globals id e
 | List (t) -> update_list_table env.globals id t
 | [])
 | Void -> raise (Error ("Can't declare void"))
 | _ -> update_prim_table env.globals id ty e)
 in
 let new_env = {env with globals = new_table} in
 (SVar_Decl_Assign (id, ty, get_sexpr env e),
 new_env)
else
 raise (Error ("Type mismatch in global variable
 assignment"))
| Access_Assign (e1, e2) -> let t1 = check_expr env e1
 and tr = check_expr env e1 in
if (t1 = tr) then
 (SAccess_Assign (get_sexpr env e1, get_sexpr env
 e2), env)
else
 raise (Error ("Type mismatch in assignment!"))
| Constructor (ty, id, l) ->
prerr_string ("Constructor " ^ id ^ " being created\n")
;
let list_to_check = (match ty with

```

```

 Node -> env.node_types
 | Rel -> env.rel_types
 | _ -> raise (Error("Can't declare constructors
 of non Node or Rel types")) in
let does_type_exist = List.exists (fun (cid, _) -> cid =
 id) list_to_check in
if does_type_exist then
 raise (Error("Already have a constructor of this
 name for this type"))
else
 let new_constructors = (id, l) :: list_to_check
 in
 let new_env = (match ty with
 Node -> if List.exists (fun (vid, _) -> vid = id)
 env.rel_types then
 raise (Error("Can't have a constructor
 for both Node and Rel types")) else
 {env with node_types = new_constructors
 }
 | Rel -> if List.exists (fun (vid, _) -> vid = id)
 env.node_types then
 raise (Error("Can't have a constructor
 for both Rel and Node types"))
 else {env with rel_types =
 new_constructors}
 | _ -> raise (Error("Can only have constructors
 for Node and Rel types")) in
 (SConstructor(ty, id, get_sformal_list l []),
 new_env) in
 (checked_global, up_env)

let rec check_globals_and_update_env env vars =
 checked_vars =
let (checked_globals, new_env) =
 (match vars with
 | var :: tail ->
let (checked_global, up_env) = check_global_env var in

```

```

check_globals_and_update_env up_env tail (
 checked_global :: checked_vars)
| [] -> (checked_vars, env)
in (checked_globals, new_env)

let run_program program =
let (vars, funcs) = program in
let env = beginning_environment in
let (checked_globals, new_env) =
 check_globals_and_update_env env (List.rev vars) [] in
let checked_functions = check_functions new_env (List
 rev funcs) [] in
SProg(checked_globals, checked_functions)

```

## 8.5 SAST

```

open Ast

type sformal =
 SFormal of Ast.n2n_type * string

type svar_decl =
 SVar of Ast.n2n_type * string
 | SConstructor of Ast.n2n_type * string * sformal
 list
 | SVar_Decl_Assign of string * Ast.n2n_type * sexpr
 | SAccess_Assign of sexpr * sexpr

and sexpr =
 SLiteral of sliteral * Ast.n2n_type
 | SId of string * Ast.n2n_type
 | SBinop of sexpr * Ast.op * sexpr * Ast.n2n_type
 | SGrop of sexpr * Ast.grop * sgraph_component * Ast.
 n2n_type
 | SGeop of sexpr * Ast.geop * sformal * Ast.n2n_type

```

```

| SUNop of Ast.uop * sexpr * Ast.n2n_type
| SAccess of string * string * Ast.n2n_type
| SCall of string * sexpr list * Ast.n2n_type
| SFunc of sbuilt_in_function_call * Ast.n2n_type
| SComplex of scomplex_literal * Ast.n2n_type

and sliteral =
 SInt_Literal of int
 | SDouble_Literal of float
 | SString_Literal of string
 | SBool_Literal of bool
 | SAny

and sbuilt_in_function_call =
 SFindMany of string * sfind_many
 | SMap of string * Ast.n2n_type * smap_function
 | SNeighbors_Func of string * string

and sfind_many =
 | SFind_Many_Node of scomplex_literal
 | SFind_Many_Gen of sgraph_type * sgraph_type

and smap_function =
 | SMap_Func of string * sstatement list

and scomplex_literal =
 | SGraph_Literal of sgraph_component list
 | SGraph_Element of (Ast.n2n_type * string) * (string
 * Ast.n2n_type * sliteral) list

and sgraph_component =
 SNode_Rel_Node_tup of sgraph_type * sgraph_type *
 sgraph_type

and sgraph_type =
 | SGraph_Id of string
 | SGraph_type of scomplex_literal

```

```

and sstatement =
 SBlock of sstatement list
 | SExpr of sexpr
 | SReturn of sexpr
 | SIf of sexpr * sstatement * sstatement
 | SVar_Decl of svar_decl

type sfunc_decl = {
 sfname : string;
 sformals : sformal list;
 sbody : sstatement list;
 sreturn_type : Ast.n2n_type;
}

type sprogram =
 SProg of svar_decl list * sfunc_decl list

```

## 8.6 Code Generation

```

open Ast
open Sast
open Printf

let imports =
 "package com.n2n;\n\n" ^
 "import java.util.*;\n\n"

let set_default_val ty = match ty with
| Int -> "0"
| String -> "\"\""
| Double -> "0.0"
| Bool -> "false"
| - -> "\"\""

let rec gen_var_type = function

```

```

Int -> "int"
| Double -> "double"
| Bool -> "boolean"
| String -> "String"
| Void -> "void"
| Rel -> "Relationship"
| Node -> "Node"
| Graph -> "Graph"
| List (ty) -> "Set<" ^ gen_var_type ty ^ ">"

let gen_binop = function
| Add -> "+"
| Sub -> "-"
| Mult -> "*"
| Div -> "/"
| Mod -> "%"
| Equal -> "=="
| Neq -> "!="
| Less -> "<"
| Leq -> "<="
| Greater -> ">"
| Geq -> ">="
| And -> "&&"
| Or -> "||"
| Concat -> "+="

let gen_unop = function
| Not -> "!"
| Neg -> "-"

let gen_literal lit = match lit with
| SInt_Literal(i) -> string_of_int i
| SDouble_Literal(d) -> string_of_float d
| SBool_Literal(b) -> string_of_bool b
| SString_Literal(str) -> "\"" ^ str ^ "\""
| SAny -> "Any"

let rec gen_literal_list ll = match ll with

```

```

| [] -> ""
| head::[] -> gen_literal head
| head::tail -> gen_literal head ^ ", " ^
 gen_literal_list tail

let rec gen_expr expr = match expr with
| SLiteral(l,t) -> gen_literal l
| SId(v,t) -> v
| SComplex(c,t) -> gen_scomplex c
| SUNop(u, e, t) -> gen_unop u ^ "(" ^
 gen_expr e ^ ")"
| SBinop(e1, op, e2, t) -> (match e1, op with
 | SAccess(-,-,-), Equal | SAccess(-,-,-),
 Neq -> gen_expr e1 ^ ".equals(" ^
 gen_expr e2 ^ ")"
 | _ -> gen_expr e1 ^ gen_binop op ^ gen_expr
 e2)
| SAccess(e1,e2, t) -> e1 ^ ".getValueFor(\"" ^
 e2 ^ "\")"
| SCall(id, e1, t) -> if(id="print") then
 gen_print e1 else id ^ "(" ^ gen_expr_list e1 ^ ")"
 "
| SFunc(fname, t) -> gen_sfunc fname
| SGrop(e1, grop, nrn, t) -> gen_expr e1 ^
 gen_graph_op grop ^ gen_nrn_tup nrn ^ ")"
| SGeop(e1, geop, f1, t) -> gen_expr e1 ^
 gen_graph_elem_op geop ^
 (match f1 with
 SFormal(_, id) -> id) ^
 (match geop with
 Data_Insert -> let t =
 (match f1 with
 SFormal(ty, _) -> ty) in "\"", " ^ (
 set_default_val t) ^ ")"
 | Data_Remove -> "\"")

and gen_expr_list expr_list = match expr_list with
| [] -> ""

```

```

| head::[] -> gen_expr head
| head::tail -> gen_expr head ^ ", " ^ gen_expr_list
 tail

and gen_formal h = match h with
 SFormal(type_spec, id) -> gen_var_type type_spec ^ "
 " ^ id

and gen_formal_list fl = match fl with
| [] -> ""
| head::[] -> gen_formal head
| head::tail -> gen_formal head ^ ", " ^
 gen_formal_list tail

and gen_print p = match p with
| [] -> ""
| head::[] -> "System.out.print(" ^ gen_expr head ^
 ")"
| head::tail -> "System.out.print(" ^ gen_expr head ^
 ")" ^ gen_print tail

and gen_sgraph_type gt = match gt with
| SGraph_Id(id) -> id
| SGraph_type(s1) -> (match s1 with
 SGraph_Element(element_type, field_info) ->
 gen_gt_instantiation element_type field_info
 | _ -> "")

and gen_gt_instantiation element_type field_info =
 match element_type with
| (graph_element_type, id) -> "new " ^ gen_var_type
 graph_element_type ^ "(\\" ^ id ^ "\\" ^
 ", new HashMap<String, Object>() {{\n\t" ^ (
 gen_graph_elem element_type field_info ^ "\n
 }})\n"

and gen_scomplex c = match c with

```



```

| SGraph_Literal(nrn_list) ->
 gen_node_rel_node_tup_list nrn_list (* Should just
 be a list of graph elements *)
| SGraph_Element(element_type, field_info) ->
 gen_element_instantiation element_type field_info

and gen_element_instantiation element_type field_info =
 match element_type with
 | (graph_element_type, id) -> "\"" ^ id ^ "\"" ^
 ", new HashMap<String, Object>() {{\n\t" ^ (
 gen_graph_elem element_type field_info ") ^ "\n
 }}\n"

and gen_graph_elem element_type field_info out_str =
 match element_type with
 | (n2n_type, id) -> (match field_info with
 | (field_name, field_type, field_value)::tail ->
 let put_string = (if (field_type = String) then
 sprintf "put(%s, %s);\n" ("\" ^ field_name ^
 "\"") (if field_value = SAny then "\"Any\""
 else (gen_literal field_value)))
 else
 sprintf "put(%s, %s);\n" ("\" ^ field_name ^
 "\"") (if field_value = SAny then "\"Any\""
 else (gen_literal field_value))) in
 let new_str = out_str ^ put_string in (
 gen_graph_elem element_type tail new_str)
 | [] -> out_str)

and gen_node_rel_node_tup_list nrn_tup = match nrn_tup
 with
 | [] -> ""
 | head::[] -> gen_nrn_tup head
 | head::tail -> gen_nrn_tup head ^ ", " ^
 gen_node_rel_node_tup_list tail

and gen_nrn_tup nrn_tup = match nrn_tup with

```

```

SNode_Rel_Node_tup(sg1 , sg2 , sg3) -> "new Graph.
 Member<>(" ^ gen_sgraph_type sg1 ^ ", " ^
 gen_sgraph_type sg2 ^ ", " ^ gen_sgraph_type sg3 ^
 ")”

and gen_sfunc fname = match fname with
 (* TOASK How call Map and Neighbors function? And
 Graph/Data inserts *)
 | SFindMany(id , sfm) -> id ^ ".findMany(" ^
 gen_find_many sfm ^ ")”
 | SMap(id , ty , smf) ->
 (match ty with
 Graph -> "for(" ^ "Node " ^ gen_map id ty
 smf
 | List(t) -> "for(" ^ gen_var_type t ^ " " ^
 gen_map id ty smf
 | _ -> raise Not_found)
 | SNeighbors_Func(id1 , id2) -> id1 ^ ".neighbors("
 ^ id2 ^ ")”

and gen_find_many sfind = match sfind with
 | SFind_Many_Node(scomp) -> gen_scomplex scomp
 | SFind_Many_Gen(gt1 , gt2) -> gen_sgraph_type gt1 ^
 ", " ^ gen_sgraph_type gt2

(* TODO: Figure out what to pass into the Map function
when created in Javac Backedn *)
(* Cuz this is not right! *)
and gen_map id ty smap = match smap with
 | SMap_Func(nid , sl) -> (match ty with
 Graph -> nid ^ " : " ^ id ^ ".getMapSet()){\n" ^
 gen_sstmt_list sl ^ "}\n”
 | List(t) -> nid ^ " : " ^ id ^ "){\n" ^
 gen_sstmt_list sl ^ "}\n”
 | _ -> raise Not_found)

and gen_sstmt stmt = match stmt with

```

```

| SBlock(stmt_list) -> gen_sstmt_list stmt_list
| SExpr(expr) -> (match expr with
 SFunc(SMap(-,-,-), -) -> gen_expr expr ^ "\n\
 t"
 | - -> gen_expr expr ^ ";\n\t")
| SReturn(expr) -> "return " ^ gen_expr expr ^ ";\n\t"
"
| SIf(expr, s1, s2) -> "if(" ^ gen_expr expr ^ ") {\n\t
\t" ^ gen_sstmt s1 ^ "}\n\telse {\n\t" ^ gen_sstmt
s2 ^ "}\n\n"
| SVar_Decl(vdec) -> gen_var_dec vdec ^ ";\n\t"

and gen_sstmt_list stmt_list = match stmt_list with
| [] -> ""
| head::[] -> gen_sstmt head
| head::tail -> gen_sstmt head ^ gen_sstmt_list tail

and gen_var_dec dec = match dec with
| SVar(ty, id) -> gen_var_type ty ^ " " ^ id
| SConstructor(ty, id, formals) -> "String " ^ id ^ " = "
^ "\"" ^ id ^ "\"";\n"
| SAccess_Assign(e1, e2) -> (match e1 with
 SAccess(el, er, t) -> el ^ ".getData().put(" ^
 "\"" ^ er ^ "\", " ^ gen_expr e2 ^ ")")
 | - -> gen_expr e1 ^ " = " ^ gen_expr e2)
| SVar_Decl_Assign(id, ty, e) -> (match ty with
| Int | Double | Bool | String -> gen_var_type ty ^
" " ^ id ^ " = " ^ gen_expr e ^ ";\n"
| Rel | Node -> gen_var_type ty ^ " " ^ id ^ " =
new " ^ gen_var_type ty ^ "(" ^ gen_expr e ^ ")
;"
| List(-) -> (match e with
 SFunc(fname, t) ->
 (match fname with
 - -> gen_var_type ty ^ " " ^ id ^ " = " ^
 gen_sfunc fname)
 | - -> gen_var_type ty ^ " " ^ id ^ " = new " ^
 gen_var_type ty ^ "(" ^ gen_expr e ^ ");")

```

```

| Graph -> (match e with
 SFunc(fname, t) ->
 (match fname with
 SMap(-,-,-) -> gen_sfunc fname
 | _ -> raise Not_found)
 | SCall(s, e1, t) -> gen_var_type ty ^ " " ^ id
 ^ " = " ^ s ^ "(" ^ gen_expr_list e1 ^ ");"
 | _ -> gen_var_type ty ^ " " ^ id ^ " = new Graph
 (Arrays.asList(" ^ gen_expr e ^ ");")
| Void -> "void")(* impossible case *)

and gen_var_dec_list var_dec_list = match var_dec_list
with
| [] -> ""
| head::[] -> gen_var_dec head
| head::tail -> gen_var_dec head ^ gen_var_dec_list
tail

and gen_global_var_dec_list var_dec_list = match
var_dec_list with
| [] -> ""
| head::[] -> "static " ^ gen_var_dec head ^ ";"
| head::tail -> gen_var_dec head ^ gen_var_dec_list
tail

(* TODO: These dont currently exist in java backend *)
and gen_graph_op grop = match grop with
| Graph_Insert -> ".insert("
| Graph_Remove -> ".remove("

and gen_graph_elem_op geop = match geop with
| Data_Insert -> ".getData().put(\""
| Data_Remove -> ".getData().remove(\""

and gen_func_dec func =
 if(func.sfname = "main") then "public static void
 main(String [] args) {\n\n\t\t" ^ gen_sstmt_list
 func.sbody ^ "}\n"

```

```

else "public static " ^ gen_var_type func.
 sreturn_type ^ " " ^ func.sfname ^
 "(" ^ gen_formal_list func.sformals ^ ") {\n" ^
 gen_sstmt_list func.sbody ^ "}\n"

and gen_func_dec_list fl = match fl with
| [] -> ""
| head::[] -> gen_func_dec head
| head::tail -> gen_func_dec head ^ gen_func_dec_list
 tail

let prog_gen = function
 SProg(checked_globals, checked_functions) ->
 imports ^
 "class Main {\n\n\t" ^
 gen_global_var_dec_list checked_globals ^ ";\n" ^
 gen_func_dec_list checked_functions ^
 "\n}\n"

```

## 8.7 Java Backend

### 8.7.1 Graph.java

```

package com.n2n;

import java.util.*;
import java.util.function.Function;
import java.util.function.Predicate;
import java.util.stream.Collectors;
import java.util.stream.Stream;

public class Graph {

 private Set<Relationship> relationships = new
 HashSet<>();

```

```

/**
 * A class that encapsulates {Node, Relationship,
 * Node} triplets. Used
 * when constructing graphs to enforce this type
 * union.
 *
 * @param <N> A node.
 * @param <R> A relationship.
 */
public static class Member<N, R> {
 private N from;
 private R rel;
 private N to;

 public Member(N from, R rel, N to) {
 this.from = from;
 this.rel = rel;
 this.to = to;
 }

 public N getFrom() { return from; }
 public R getRel() { return rel; }
 public N getTo() { return to; }
}

public Set<Relationship> getRelationships(){
 return this.relationships;
}

public Graph(List<Member<Node, Relationship>>
relatedMemberList) {
 addToGraph(relatedMemberList);
}

private void addToGraph(List<Member<Node,
Relationship>> relatedMemberList) {
 relatedMemberList.stream().forEach((members) ->
 {

```

```

 members.getRel().addNodes(members.getFrom()
 , members.getTo());
 relationships.add(members.getRel());
 });
}

/**
 * An operation for finding nodes based on loose
 * relationship equality.
 *
 * Loose equality is defined as a match on the '
 * type' of the relationship.
 *
 * Example:
 * + Find movies in which Keanu acted_in
 * keanu_movies: List<Node> = find_many(keanu
 * acted_in)
 *
 * Will search the graph for matches on a 'keanu'
 * node that has a relationship of type 'acted_in'
 *
 *
 * @param node The source node from which
 * relationships start.
 * @param relationshipType The relationship type (
 * its name) that joins source node and potential
 * target nodes.
 * @return The set of target nodes, or an empty set
 * if no nodes are found.
 */
public Set<Node> findMany(Node node, String
 relationshipType) {
 return findManyHelper(r -> r.looselyEquals(
 relationshipType), r -> r.getNodesFrom(node)
 .stream());
}

```

```

public Set<Node> findMany(String type, Map<String,
 Object> data) {
 return findMany(new Node(type, data));
}

/**
 * An operation for finding nodes based on strict
 * relationship equality.
 *
 * Strict equality is defined as a match on the '
 * type' of the relationship and all the fields
 * provided in the
 * relationship.
 *
 * Example:
 * + Find movies in which Keanu acted_in as Neo
 * keanu_movies: List<Node> = find_many(keanu
 * acted_in[Neo])
 *
 * Will search the graph for matches on a 'keanu'
 * node that has a relationship of type 'acted_in'
 * with
 * first field of the relationship as "Neo"
 *
 * @param node The source node.
 * @param relationship A relationship that joins
 * source node and target nodes.
 * @return The set of target nodes, or an empty set
 * if no nodes are found.
 */
public Set<Node> findMany(Node node, Relationship
 relationship) {
 return findManyHelper(r -> r.strictlyEquals(
 relationship), r -> r.getNodesFrom(node).
 stream());
}

/**

```



```

* An operation for finding nodes based on an
 inverse, loose relationship equality.
*
* Loose equality is defined as a match on the '
 type' of the relationship.
*
* Example:
* + Find movies in which Keanu acted_in
* matrix_actors: List<Node> = find_many(
 acted_in matrix)
*
* Will search the graph for actor nodes that
 point to node 'matrix' through an 'acted_in'
 relationship
*
* @param node The source node from which
 relationships start.
* @param relationshipType The relationship type (
 its name) that joins source node and potential
 target nodes.
* @return The set of target nodes, or an empty set
 if no nodes are found.
*/
public Set<Node> findMany(String relationshipType,
 Node node) {
 return findManyHelper(r -> r.looselyEquals(
 relationshipType), r -> r.getNodesTo(node).
 stream());
}

/**
* An operation for finding nodes based on an
 inverse, strict relationship equality.
*
* Strict equality is defined as a match on the '
 type' of the relationship and all the fields
 provided in the
* relationship.

```

```

*
* Example:
* + Find actors that acted_in as 'Neo' in 'matrix'
*
* neo_actors: List<Node> = find_many(acted_in
* (Neo) matrix)
*
* Will search the graph for actor nodes that
* point to node 'matrix' through an 'acted_in'
* relationship with
* first field as 'Neo'
*
* @param node The destination node at which the
* relationship ends.
* @param relationship A relationship that joins
* source node and target nodes.
* @return The set of target nodes, or an empty set
* if no nodes are found.
*/
public Set<Node> findMany(Relationship relationship
 , Node node) {
 return findManyHelper(r -> r.strictlyEquals(
 relationship), r -> r.getNodesTo(node).
 stream());
}

/**
* Finds all relationships that join leftNode and
* RightNode.
* TODO: Ugly and inefficient. Fix me.
*
* @param leftNode Left side of the relationship.
* @param rightNode Right side of the relationship.
* @return A set that contains relationships
* between the two nodes.
*/
public Set<Relationship> findMany(Node leftNode ,
 Node rightNode) {

```

```

 Set<Relationship> result = new HashSet<>();
 result.addAll(relationshipFinder(leftNode,
 rightNode));
 result.addAll(relationshipFinder(rightNode,
 leftNode));
 return result;
}

public Set<Node> findMany(Node target) {
 Set<Node> nodes = new HashSet<>();
 for (Relationship relationship : relationships)
 {
 nodes.addAll(relationship
 .getAll()
 .stream()
 .filter(node -> node.looselyEquals(
 target)).collect(Collectors.
 toList()));
 }
 return nodes;
}

public Set<Node> neighbors(Node target) {
 return getNodesFromRelationships(r -> r.
 getNodesFrom(target).stream());
}

public Set<Node> getMapSet() {
 return getNodesFromRelationships(r -> r.getAll
 ().stream());
}

private Set<Node> getNodesFromRelationships(
 Function<Relationship, Stream<? extends Node>>
 mapper) {
 return relationships.stream().flatMap(mapper).
 collect(Collectors.toSet());
}

```

```

private Set<Relationship> relationshipFinder (Node
left , Node right) {
 Set<Relationship> result = new HashSet<>();
 for (Relationship relationship : relationships)
 {
 Set<Node> nodesFromLeft = relationship.
 getNodesFrom (left);
 Set<Node> nodesToRight = relationship.
 getNodesTo (right);
 if (!nodesToRight.isEmpty()) {
 boolean modified = nodesFromLeft.
 retainAll (nodesToRight);
 if (modified) {
 result.add (relationship);
 }
 }
 }
 return result;
}

private Set<Node> findManyHelper (Predicate<
Relationship> predicate , Function<Relationship ,
Stream<? extends Node>> mapper) {
 return relationships.stream ()
 .filter (predicate)
 .flatMap (mapper)
 .collect (Collectors.toSet ());
}

public void insert (List<Member<Node , Relationship>>
relatedMemberList) {
 addToGraph (relatedMemberList);
}

@Override
public String toString () {
 return "Graph{" +

```

```

 "relationships=" + relationships +
 '>';
 }
}

```

## 8.7.2 Node.java

```

package com.n2n;

import java.util.Map;

public class Node {

 private String type;
 private Map<String, Object> data;

 public Node(String type, Map<String, Object> data)
 {
 this.type = type;
 this.data = data;
 }

 public Object getValueFor(String field) {
 return this.data.get(field);
 }

 public Map<String, Object> getData() {
 return this.data;
 }

 public boolean looselyEquals(Object other) {
 if (this == other) return true;
 if (other == null || getClass() != other.
 getClass()) return false;

 Node otherNode = (Node) other;
 }
}

```

```

 return type.equals(otherNode.type) &&
 dataLooselyEquals(otherNode.getData());
 }

 private boolean dataLooselyEquals(Map<String,
 Object> other) {
 /*if (!this.data.keySet().containsAll(other.
 keySet())) {
 return false;
 }*/

 for (Map.Entry<String, Object> entry : other.
 entrySet()) {
 if (!entry.getValue().equals("Any") && !
 this.data.get(entry.getKey()).equals(
 entry.getValue())) {
 return false;
 }
 }
 return true;
 }

 @Override
 public boolean equals(Object o) {
 if (this == o) return true;
 if (o == null || getClass() != o.getClass())
 return false;

 Node node = (Node) o;

 return data.equals(node.data) && type.equals(
 node.type);
 }

 @Override
 public int hashCode() {
 int result = type.hashCode();
 result = 31 * result + data.hashCode();
 }

```

```

 return result;
 }

 @Override
 public String toString() {
 StringBuilder sb = new StringBuilder();
 int count = this.data.size();
 int i=0;
 sb.append(type).append("{ ");
 for (Map.Entry<String, Object> data : this.data
 .entrySet()) {
 sb.append(data.getKey()).append(" = ").
 append(data.getValue());
 if (i < count-1)
 sb.append(", ");
 i++;
 }
 sb.append("}");
 return sb.toString();
 }
}

```

### 8.7.3 Relationship.java

```

package com.n2n;

import java.util.*;

public class Relationship {

 private String type;
 private Map<String, Object> data;
 private Map<Node, Set<Node>> fromTo = new HashMap
 <>();
 private Map<Node, Set<Node>> toFrom = new HashMap
 <>();
}

```

```

public Relationship(String type, Map<String, Object
 > data) {
 this.type = type;
 this.data = data;
}

public Relationship(String type) {
 this(type, Collections.emptyMap());
}

public Object getValueFor(String field) {
 return this.data.get(field);
}

public Set<Node> getAll() {
 Set<Node> nodes = new HashSet<>();
 nodes.addAll(fromTo.keySet());
 nodes.addAll(toFrom.keySet());
 return nodes;
}

public void addNodes(Node from, Node to) {
 if (fromTo.containsKey(from)) {
 fromTo.get(from).add(to);
 } else {
 fromTo.put(from, new HashSet<>(Arrays.
 asList(to)));
 }
 if (toFrom.containsKey(to)) {
 toFrom.get(to).add(from);
 } else {
 toFrom.put(to, new HashSet<>(Arrays.asList(
 from)));
 }
}

public Set<Node> getNodesFrom(Node from) {

```



```

 return fromTo.containsKey(from) ? new HashSet
 <>(fromTo.get(from)) : Collections.emptySet
 ();
 }

 public Set<Node> getNodesTo(Node to) {
 return toFrom.containsKey(to) ? new HashSet<>(
 toFrom.get(to)) : Collections.emptySet();
 }

 public boolean looselyEquals(String type) {
 return this.type.equals(type);
 }

 public boolean strictlyEquals(Relationship
 relationship) {
 return this.equals(relationship);
 }

 @Override
 public boolean equals(Object o) {
 if (this == o) return true;
 if (o == null || getClass() != o.getClass())
 return false;

 Relationship that = (Relationship) o;

 return type.equals(that.type) && data.equals(
 that.data);
 }

 @Override
 public int hashCode() {
 int result = type.hashCode();
 result = 31 * result + data.hashCode();
 return result;
 }

```

```

@Override
public String toString() {
 StringBuilder sb = new StringBuilder();
 int count = this.data.size();
 int i = 0;
 sb.append(type).append(" {");
 for (Map.Entry<String, Object> data : this.data
 .entrySet()) {
 sb.append(data.getKey()).append(" =").
 append(data.getValue());
 if (i < count-1)
 sb.append(", ");
 i++;
 }
 sb.append("}");
 return sb.toString();
}
}

```

## 8.8 Tests

### 8.8.1 arithmetic1.n2n

```

fn main () -> Void {
 print(2+4);
}

```

### 8.8.2 arithmetic2.n2n

```

fn main () -> Void {
 print(3-6);
 print("\n");
 print(6*3);
 print("\n");
 print(6/3);
}

```

```
 print("\n");
 print(15%4);
 print("\n");
}
```

### 8.8.3 comparison1.n2n

```
fn main () -> Void {
 print(2>3);
 print("\n");
 print(2<3);
 print("\n");
}
```

### 8.8.4 comparison2.n2n

```
fn main () -> Void {
 print(2==4);
 print("\n");
 print(2!=4);
 print("\n");
 print(2>=4);
 print("\n");
 print(2<=4);
 print("\n");
}
```

### 8.8.5 declaration.n2n

```
;;Test the multiple declaration of a variable local
 scoping;;

a: Int = 100;

fn main () -> Void {
```

```

 print(a);
 print("\n");
 a = 300;
 print(a);
 print("\n");
 a: Int = 1; ;;Not sure what happens when
 declared again;;
 print(a);
 print("\n");

 b: Int = 200;
 print(b);
 print("\n");
}

```

### 8.8.6 escapestr1.n2n

```

fn main() -> Void {
 a: String = "a\n";
 print(a);
}

```

### 8.8.7 escapestr2.n2n

```

fn main() -> Void {
 a: String = "a\t";
 print(a);
 print("\n");
}

```

### 8.8.8 escapestr3.n2n

```

fn escapestr3() -> Void {
 a: String = "a\\";
 print(a);
}

```

```
}

fn main() -> Void {
 escapestr3();
}
```

### 8.8.9 escapestr5.n2n

```
fn escapestr3() -> Void {
 a: String = "a\'";
 print(a);
}

fn main() -> Void {
 escapestr3();
}
```

### 8.8.10 float.n2n

```
fn main () -> Void {
 print(8/6);
}
```

### 8.8.11 float2.n2n

```
fn main () -> Void {
 print(8.0/6.0);
}
```

### 8.8.12 helloworld.n2n

```
fn main () -> Void {
 print("hello \uworld");
}
```

### 8.8.13 `init1.n2n`

```
fn main() -> Void {
 a: Int = 66;
 b: Int = 123456;
 print(a);
 print("\n");
 print(b);
 print("\n");
}
```

### 8.8.14 `functions1.n2n`

```
fn foo () -> Void {
 print(5);
}

fn main () -> Void {
 foo();
 print("\n");
}
```

### 8.8.15 `functions2.n2n`

```
fn foo () -> Int {
 ;;Do something;;
 return 5;
}

fn main () -> Void {
 a: Int = foo();
 print(a);
}
```

```
 print("\n");
}
```

### 8.8.16 functions3.n2n

```
fn foo (c: Int) -> Int {
 ;;Do something;;
 return c+5;
}

fn bar (a: Int) -> Int {
 b: Int = a;
 return foo(b);
}

fn main () -> Void {
 a: Int = bar(61);
 print(a);
}
```

### 8.8.17 find\_many1.n2n

```
actor: Node = { name: String , age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String , year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];
}
```

```

Cast: Graph = <
 Keanu neo matrix ,
 Leo jordan wolf
>;

missing_rel: List<Rel>;
missing_rel = Cast.find_many(Keanu, matrix);

print(missing_rel);
}

```

### 8.8.18 find\_many2.n2n

```

actor: Node = { name: String , age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String , year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];
 nelson: Rel = actedIn["Nelson"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];
 sweet_nov: Node = movie["Sweet November",
 2000];

 Cast: Graph = <
 Keanu neo matrix ,
 Leo jordan wolf ,
 Keanu nelson sweet_nov
 >
}

```



```

>;

point_to: List<Node>;
point_to = Cast.find_many(Keanu, neo);

print(point_to);

}

```

### 8.8.19 find\_many3.n2n

```

actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];
 nelson: Rel = actedIn["Nelson"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];
 sweet_nov: Node = movie["Sweet November",
 2000];

 Cast: Graph = <
 Keanu neo matrix,
 Leo jordan wolf,
 Keanu nelson sweet_nov
 >;

 point_from: List<Node>;
 point_from = Cast.find_many(neo, matrix);
}

```

```
 print(point_from);
 print("\n");
}
```

### 8.8.20 find\_many4.n2n

```
actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];
 nelson: Rel = actedIn["Nelson"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];
 sweet_nov: Node = movie["Sweet November",
 2000];

 Cast: Graph = <
 Keanu neo matrix,
 Leo jordan wolf,
 Keanu nelson sweet_nov
 >;

 node_lit: List<Node>;
 node_lit = Cast.find_many(actor["Keanu", -]);

 print(node_lit);
 print("\n");
}
```

```
}
```

### 8.8.21 graph.n2n

```
actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];

 Cast: Graph = <
 Keanu neo matrix,
 Leo jordan wolf
 >;

 print(Cast);
}
```

### 8.8.22 graph2.n2n

```
actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {
```

```

Cast: Graph = <
 actor["Keanu", 35] actedIn["Neo"] movie
 ["Matrix", 1999],
 actor["Leo", 20] actedIn["Jordan"]
 movie["Wolf", 1994]
>;

print(Cast);
}

```

### 8.8.23 graphInsDel.n2n

```

actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];

 Cast: Graph = <
 Keanu neo matrix,
 Leo jordan wolf
 >;

 print(Cast);
 print("\n");
 print("\n");
}

```

```

Cast ^+ (Keanu actedIn ["Nelson"] movie ["Sweet L
November" , 2000]);
print(Cast);
print("\n");
print("\n");

Cast ^- (Keanu actedIn ["Nelson"] movie ["Sweet L
November" , 2000]);
print(Cast);
print("\n");
print("\n");
}

```

### 8.8.24 if.n2n

```

fn main() -> Void {
 if (true) {
 print("success");
 print("\n");
 }
 else {

 }

 if(4<5){
 print(6);
 print("\n");
 if(0>1){
 print(7);
 }
 else{
 print("else");
 print("\n");
 }
 }
}

```

---

### 8.8.25 map.n2n

```
actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];

 Cast: Graph = <
 Keanu neo matrix,
 Leo jordan wolf
 >;

 visited_node: List<Node>;
 mutated_graph: Graph = Cast.map(node in {node
 [+] visited: Bool;});

 Keanu.visited = true;

 ;;Need to be considered;;

 print(Keanu.visited);
 print(matrix.visited);
 print(Leo.visited);
 print(wolf.visited);
```

```
}
```

### 8.8.26 neighbor.n2n

```
actor: Node = { name: String, age: Int };
actedIn: Rel = { role: String };
movie: Node = { title: String, year: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Leo: Node = actor["Leo", 20];

 neo: Rel = actedIn["Neo"];
 jordan: Rel = actedIn["Jordan"];

 matrix: Node = movie["Matrix", 1999];
 wolf: Node = movie["Wolf", 1994];

 Cast: Graph = <
 Keanu neo matrix,
 Leo jordan wolf
 >;

 Keanu_movies: List<Node>;
 Keanu_movies = Cast.neighbors(Keanu);

 print(Keanu_movies);
 print("\n");
}
```

### 8.8.27 node.n2n

```
actor: Node = { name: String, age: Int };
```

```

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 print(Keanu);
 print("\n");

}

```

### 8.8.28 nodeAccess.n2n

```

actor: Node = { name: String, age: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 Keanu.name = "Reeves";
 print(Keanu.name);
 print("\n");

}

```

### 8.8.29 nodeInsDel.n2n

```

actor: Node = { name: String, age: Int };

fn main() -> Void {

 Keanu: Node = actor["Keanu", 35];
 print(Keanu);
 print("\n");

 Keanu [+] visited: Bool;
 Keanu.visited = true;
 print(Keanu);
 print("\n");

 Keanu [-] visited: Bool;

```



```
 print(Keanu);
 print("\n");
}
```

### 8.8.30 rel.n2n

---

### 8.8.31 relAccess.n2n

```
actedIn: Rel = { role: String };
fn main() -> Void {
 neoRole: Rel = actedIn["Neo"];
 print(neoRole);
 print("\n");
}
```

### 8.8.32 simple.n2n

```
a: Int = 5;
b: String = "test";
c: Node = {name: String, age: Int};
fn t()->Void {}
;; comment;;
fn main()->Void {
 a: Int = 5;
 print(5);
}
```

### 8.8.33 simple2.n2n

```
fn main()->Void {
 a: Int = 5;
}
```