## COMSW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

Lecture 1
Spring 2011
Instructor: Michele Merler

## Course Information - Goals

"A general introduction to computer science concepts, algorithmic problemsolving capabilities, and programming skills in C"

University bulletin

- Learn how to program, in C
- Understand basic Computer Science problems
- Learn about basic data structures
- Start to think as a computer scientist
- Use all of the above to solve real world problems


## Course Information - Instructor

- Michele Merler
- Email: mmerler@cs.columbia.edu or mm3233@columbia.edu
- Office : 624 CEPSR
- Office Hours: Friday 12pm-2pm
- $4^{\text {th }}$ year PhD Student in CS Department
- Research Interests:
- Image \& Video Processing
- Multimedia
- Computer Vision


## Course Information- TA

- TDB
- Email:TDB@columbia.edu
- Office : TA room
- Office Hours: TDB


## Course Information- Courseworks

We will be using Courseworks (https://courseworks.columbia.edu/) for:

- Message board for discussions
- Submit Homeworks
- Grades

Check out the board before you send an email to the instructor or the TA, the answer you are looking for could already be there!

## Course Information Requirements and Books

Requirements

- Basic computer skills
- CUNIX account

Textbooks
 by Brian Kernighan and Dennis Ritchie
http://www1.cs.columbia.edu/~mmerler/coms1003-1/C Programming Language.rar

- Practical C Programming (3rd Edition) by Steve Oualline


## Course Information - Grading

- 5 Homeworks ( $10 \%, 10 \%, 10 \%, 10 \%, 10 \%$ )
- Midterm Exam (20\%)
- Final Exam (30\%)



## Course Information Academic Honesty

It's quite simple:

- Do not copy from others
- Do not let others copy from you

Do your homework individually
Please read through the department's policies on academic honesty http://www.cs.columbia.edu/education/honesty/

## Course Information - Syllabus

Go to class webpage
http://www1.cs.columbia.edu/~mmerler/coms1003-1 files/Syllabus.html

## What is Computer Science?

Computer science (sometimes abbreviated CS) is the study of the theoretical foundations of information and computation, and of practical techniques for their implementation and application in computer systems

Wikipedia
"Computer science and engineering is the systematic study of algorithmic processes-their theory, analysis, design, efficiency, implementation, and application-that describe and transform information"
Comer, D. E.; Gries, D., Mulder, M. C., Tucker, A., Turner, A. J., and Young, P. R. (Jan. 1989). "Computing as a discipline". Communications of the ACM 32 (1): 9.
"Computer science is the study of information structures"
Wegner, P. (October 13-15, 1976). "Research paradigms in computer science". Proceedings of the 2nd international Conference on Software Engineering. San Francisco, California, United States
"Computer Science is the study of all aspects of computer systems, from the theoretical foundations to the very practical aspects of managing large software projects."
Massey University


## What is Computer Science?

Computer Science is the discipline that studies how to make computers perform tasks that are too complex or boring for humans


## Computer Science Areas



Algorithms and data structures


## Why programming?

- We need a way to tell computers what to do
- It would be nice to communicate with computers in English, but...
- English can be ambiguous!
- Computers only understand binary!
- Solution: programming languages

```
    There are lo
kinds of people
    in the world:
        those who
        understand
binary coder and
those who don't.
```


## What is a Program?

- A Program is a sequence of instructions and computations
- We'll be designing programs in this course.
- These programs will be based on algorithms
- An Algorithm is a step-by-step problemsolving procedure


## Example

- Add 3 large numbers
- $453+782+17,892$
- Hard to do all at once
- Solution: "divide and impera"!
- $(453+782)+17,892=$
- $1,235+17,892=19,127$

- Algorithms help us divide and organize complex problems into sub-problems which are easier to solve (bottom-up approach)


## Programming

- Back in the day, programmers wrote in Assembly, a language where each word stands for a single instruction

| add | eax, ēdx |
| :--- | :--- |
| shl | eax, 2 |
| add | eax, edx |
| shr | eax, 8 |
| sub | cl, al |

- But then they had to hand translate each instruction into binary!!!
- Solution: the assembler, a computer program to do the translation
- From then, programmers could worry only about writing assembly code
- Then they started to devise higher level languages (FORTRAN, COBOL, PASCAL, C, C++, JAVA, Perl, Python, etc.), which get translated into Assembly by compilers (we will use GCC, a C compiler for Unix)

> High level language

Assembly
Compiler

Machine Instructions

## What is $\boldsymbol{C}$ ?

- Programming language developed by Dennis Ritchie in 1972 at AT\&T Bell labs
- Why is it named "C"?

Well... the B programming language already existed !

- $C$ is still the most used programming language for Operating Systems
- Popular because:
- Flexible
- C compiler was widely available
- Basis for other popular programming languages: $\mathrm{C}++$, $\mathrm{C} \#$


## What is C?

- Among the "high level" programming languages, C is one with the lowest level of abstraction
- Close to English, but more precise!
- Easy to compile into Assembly => Fast
- Rich set of standard function = we don't have to implement everything from scratch!



## Why H? Interesting Facts ...

Approximation of popularity of language using Yahoo API http://www.langpop.com/


Slide credit: Priyank Singh

## Why H? Interesting Facts ...

Available language code available using Google code search http://www.langpop.com/


Slide credit: Priyank Singh

## Why H? Interesting Facts

Jobs posting on craiglist.org, from website http://www.langpop.com/


Slide credit: Priyank Singh

## C/C++ Industry



## Example of C program

## Hello world!

## Announcements

- Homework 0 is out! Due at the beginning of next class
- Bring your laptop to class


## COM SW 1003-1

## Introduction to Computer Programming in $\boldsymbol{C}$

## Lecture 2

Spring 2011
Instructor: M ichele M erler

## Announcements

- Exercise1 is out
- We have a TA!

Gaurav Agarwal

- M S student in CS department
- Email: ga2310@columbia.edu
- Office Hours: Tuesday 11am-12pm in M udd 122A (TA room)


## What is a Program?

- A Program is a sequence of instructions and computations
- We'll be designing programs in this course.
- These programs will be based on algorithms
- An Algorithm is a step-by-step problemsolving procedure


## Example

- Add 3 large numbers
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- $1,235+17,892=19,127$

- Algorithms help us divide and organize complex problems into sub-problems which are easier to solve (bottom-up approach)


## What is $\mathbf{C}$ ?

- Programming language developed by Dennis Ritchie in 1972 at AT\&T Bell labs
- Why is it named "C"? Well... the B programming language already existed !

- $C$ is one of the high level programming language with the lowest level of abstraction
- Low to be close to assembly and machine language $\longrightarrow$ fast!
- High to be programmable by humans without (too many) headaches



## CUNIX

- CUNIX refers to the Columbia Unix environment
- For you: place where you develop your programs!


## Accessing CUNIX remotely

- Secure Shell or SSH is a network protocol that allows data to be exchanged using a secure channel between two networked devices
- The SCP protocol is a network protocol that supports file transfers



## Code Developing Tools - Linux and Mac

- Open terminal
- SSH to cunix.cc.columbia.edu
ssh yourUNI@cunix.cc.columbia.edu
- Data transfer: scp or get/put
- Copying file to host: scp SourceFile user@host:directory/TargetFile
- Copying file from host:
scp user@host:/directory/SourceFile TargetFile
For MAC: use FUGU (graphical data transfer tool) http://www.columbia.edu/acis/software/fugu/ http://download.cnet.com/Fugu/3000-2155_4-26526.html


## Code Developing Tools - Linux and Mac

To use windowing environment:

M ac users need only start X11 (found in the Utilities folder) and log in to the X11 terminal like this:
ssh -X username@cunix.cc.columbia.edu

- Linux users: see X-Windows section in CUNIX tutorial


## Code Developing Tools - Windows

- Xming and Putty to SSH and visualization
- http://sourceforge.net/projects/xming/
- http://www.chiark.greenend.org.uk/ -sgtatham/putty/download
- WinSCP for data transfer
- http://winscp.net/eng/download.php\#download2
- Notepad++for editing (can be used in combination with WinSCP)
- http://notepad-plus-plus.org/


## Code Developing Tools - Windows

- Launch Xming
- Open a session in putty with Host Name
- cunix.cc.columbia.edu



## Code Developing Tools - Windows

- M ake sure the X11 option of the SSH category is enabled



# Code Developing Tools - Windows <br> - Use WinScp to transfer files 



## Code Developing Tools－Windows

－Use WinScp to transfer files

| §y mm3233－mm3233＠cunix．cc．columbia．edu－WinSCP |  |  |  |  |  |  |  |  | $\square \times$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Local Mark Files Commands Session Options Remote Help |  |  |  |  |  |  |  |  |  |
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| （t）． |  | Parent directory | 6／2／2010 | t． |  |  |  |  |  |
| －Ask and Record Toolbar |  | File Folder | 12／7／2009 | （1）public＿html |  |  |  |  |  |
| $\square$ Downloads |  | File Folder | 9／10／2010 | Comsiou3 |  |  |  |  |  |
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| Didm Music |  | File Folder | 6／20／2009 | $0 . e m a c s, d$ |  |  |  |  |  |
| 國My Pictures |  | File Folder | 917／2010 | 句mbox |  |  |  |  | 1，485 |
| My Recordings |  | File Folder | 12／8／2009 | 句 Xauthority |  |  |  |  | 361 |
| ${ }^{\text {m M M Videos }}$ |  | File Folder | 1／22／2008 |  |  |  |  |  | 172 |
| OUDC Output Files |  | File Folder | 10／22／200 | 畇 forward |  |  |  |  | 31 |
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| 7 Visual Studio 2008 |  | File Folder | 7／30／2010 | 國 bash＿history |  |  |  | 1，178 |  |
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| 0 B of $12,807 \mathrm{KiB}$ in 0 of 15 |  |  |  | 0 B of $3,272 \mathrm{~B}$ in 0 of 10 |  |  |  |  |  |
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## Code Developing Environment

CUNIX Tutorial

## Compiling your C code

- GCC : GNU Compiler Collection
- When you invoke GCC, it normally does preprocessing, compilation, assembly and linking
- Basic Command
- gcc myProgram.c
- ./a.out

Run compiled program (executable)

- M ore advanced options
- gcc -Wall -o myProgram myProgram.c
- ./myProgram


## Compiling your C code

- GCC : GNU Compiler Collection
- When you invoke GCC, it normally does preprocessing, compilation, assembly and linking
- Basic Command
- gcc myProgram.c
- ./a.out

Run compiled program (executable)
Display all types of warnings, not only errors

Specify name of the executable

- gcc -Wall -o myProgram myProgram.c
- ./myProgram Run compiled program (executable)


## Assignment

- Read PCP Ch 1
- Read PCP Ch 2, pages 11 to 15,33


## COM SW 1003-1

## Introduction to Computer Programming in $\boldsymbol{C}$

Lecture 3
Spring 2011
Instructor: M ichele M erler

## Today

- Computer Architecture (Brief Overview)
- "Hello World" in detail
- C Syntax
- Variables and Types
- Operators
- printf (if there is time)


## Von Neumann Architecture



## Computer M emory Architecture



## Von Neumann Architecture

The Program Counter (PC) points (=tells the CPU) to the address in memory where the next instruction to be executed resides

Address 1
Address 2
. .

## Von Neumann Architecture

Hello World
\#include <stdio.h>
int main()\{ printf("Hello World ${ }^{n}$ ");
return(0);
\}


## Von Neumann Architecture

Hello World
\#include <stdio.h>
int main()\{ printf("Hello World\n"); PC $\longrightarrow$-Address n+1 return(0); return(0);
\}

| Address 1 |  |
| :---: | :---: |
| - Address 2 | M ain |
| $\cdot$ | Memory |
| $\cdot$ |  |
| Address n | printf("Hello Worldln"); |
| Address n+1 | return(0); |
|  |  |
| $\cdot$ |  |
| $\cdot$ |  |
| - Address N |  |

## The Operating System

- M anages the hardware
- Allocates resources to programs
- Accommodates user requests
- First program to be executed when computer starts (loaded from ROM)

- Windows
- Unix
- Mac OS
- Android
- Linux
- Solaris
- Chrome OS

Hardware

Program

## Hello World



## C Syntax

- Statements
- one line commands
- always end with ;
- can be grouped between \{ \}
- spaces are not considered
- Comments
// single line comment
/* multiple lines comments
*/


## Hello World +Comments

```
/*
    * My first C program
    */
#include <stdio.h>
int main(){
printf("Hello World\n");
\[
\text { return(0); } / / \text { return } 0 \text { to the } O S=O K
\]

\section*{Variables and types}
- Variables are placeholders for values
\[
\begin{aligned}
& \text { int } x=2 ; \\
& x=x+3 ; / / x \text { value is } 5 \text { now }
\end{aligned}
\]
- In C, variables are divided into types, according to how they are represented in memory (always represented in binary)
- int
- float
- double
- char

\section*{Variables Declaration}
- Before we can use a variable, we must declare (= create) it
- When we declare a variable, we specify its type and its name
```

int x;
float y = 3.2;

```
- M ost of the time, the compiler also allocates memory for the variable when it's declared. In that case declaration \(=\) definition
- There exist special cases in which a variable is declared but not defined, and the computer allocates memory for it only at run time (will see with functions and external variables)

\section*{int}
- No fractional part or decimal point (ex. +3, -100)
- Represented with 4 bytes ( 32 bits) in UNIX
- Sign
- unsigned : represents only positive values, all bites for value Range: from 0 to 2^32
- signed (default) : 1 bit for sign +31 for actual value Range: from -2^31 to 2^31
- Size
- short int : at least 16 bits
- long int: at least 32 bits
```

int x = -12;
unsigned int x = 5;
short (int) x = 2;

```
- long long int: at least 64 bits
- size(short) \(\leq\) size(int) \(\leq\) size(long)

\section*{float}
- Single precision floating point value
- Fractional numbers with decimal point
- Represented with 4 bytes (32 bits)
- Range: -10^(38) to \(10^{\wedge}(38)\)
\[
\text { float } x=11.5
\]
- Exponential notation :-0.278* \(10 \wedge \underset{m}{3}\)

\[
n_{10}=(-1)^{s} \cdot\left(f \cdot 2^{-23}\right) \cdot 2^{m-127}
\]

\section*{double}
- Double precision floating point
- Represented with 8 bytes (64 bits)

\[
\text { double } x=121.45
\]

\section*{char}
- Character
- Single byte representation
- 0 to 255 values expressed in the ASCII table
\[
\operatorname{char} \mathrm{C}={ }^{6} \mathrm{~W}
\]

\section*{ASCII Table}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline Dec & HxO & Oct Char & & Dec & Hx & & Html & Chr & & Dec & & & Html & Chr & Dec & c Hx & Oct & , & \\
\hline & & 000 NUL & (null) & 32 & 20 & 040 & \&\#32; & Space & & 64 & & 100 & \&\#64; & 0 & 96 & 60 & 140 & \&\#96; & \\
\hline & & 001 50H & (start of heading) & 33 & 21 & 041 & \&\#33; & ! & & 65 & & 101 & \&\#65; & ; A & 97 & 61 & 141 & \&\#97: & a \\
\hline 2 & 20 & 002 STX & (start of text) & 34 & 22 & 042 & \&\#34; & & & 66 & & 102 & \&\#66; & ; B & 98 & 62 & 142 & \&\#98; & b \\
\hline 3 & 30 & 003 ETX & (end of text) & 35 & 23 & 043 & \&\#35; & \# & & 67 & & 103 & \&\#67; & - C & 99 & 63 & 143 & \&\#99; & \\
\hline 4 & 40 & 004 E0T & (end of transmission) & 36 & 24 & 044 & \&\#36; & - & & 68 & & 104 & \&\#68; & ; & 100 & 64 & 144 & \&\#100; & d \\
\hline 5 & 50 & 005 ENQ & (enquiry) & 37 & 25 & 045 & \&\#37; & \% & & 69 & 45 & 105 & \&\#69; & ; E & 101 & 65 & 145 & \&\#101; & \\
\hline 6 & 60 & 006 ACK & (acknowledge) & 38 & 26 & 046 & \&\#38; & \& & & 70 & & 106 & \& \({ }_{\text {\# }} 70\) & ; F & 102 & 66 & 146 & \&\#102; & f \\
\hline 7 & 70 & 007 BEL & (bell) & 39 & 27 & 047 & \&\#39; & & & 71 & & 107 & \&\#71; & ; G & 103 & 67 & 147 & \&\#103 & \\
\hline 8 & 80 & 010 BS & (backspace) & 40 & 28 & 050 & \& \# 40; & ( & & 72 & 48 & 110 & \&\#72; & ; & 104 & 68 & 150 & \&\#104; & \\
\hline 9 & 90 & 011 TAB & (horizontal tab) & 41 & 29 & 051 & \&\#41; & ) & & 73 & & 111 & \&\#73; & ; I & 105 & 69 & 151 & \&\#105; & i \\
\hline 10 & A 0 & 012 LF & (NL line feed, new line) & 42 & 2A & 052 & \&\#42; & * & & 74 & & 112 & \&\#74; & ; & 106 & 6A & 152 & \&\#106; & \\
\hline 11 & B 0 & 013 VT & (vertical tab) & 43 & 2B & 053 & \&\#43; & + & & 75 & 4 B & 113 & \&\#75; & ; & 107 & 6B & 153 & \&\#107; & \\
\hline 12 & C 0 & 014 FF & (NP form feed, new page) & 44 & 2 C & 054 & \&\#44; & & & 76 & 4 C & 114 & \& \({ }^{\prime \prime} 76\); & - & 108 & 6C & 154 & \&\#108; & \\
\hline 13 & D 0 & 015 CR & (carriage return) & 45 & 2D & 055 & \&\#45; & - & & 77 & 4 D & 115 & \&\#77; & ; M & 109 & 6D & 155 & \&\#109; & \\
\hline 14 & E 0 & 016 S0 & (shift out) & 46 & 2 E & 056 & \&\#46; & & & 78 & & 116 & \&\#78; & ; N & 110 & 6 E & 156 & \&\#110; & \\
\hline 15 & F 0 & 017 SI & (shift in) & 47 & 2 F & 057 & \&\#47; & / & & 79 & & 117 & \&\#79; & ; 0 & 111 & 6 F & 157 & \&\#111; & \\
\hline 16 & 100 & 020 DLE & (data link escape) & 48 & 30 & 060 & \&\#48; & 0 & & 80 & 50 & 120 & \&\#80; & ; P & 112 & 70 & 160 & \&\#112; & \\
\hline 17 & 110 & 021 DCL & (device control 1) & 49 & 31 & 061 & \& \({ }^{\text {\# }} 49\); & 1 & & 81 & & 121 & \&\#81; & ; Q & 113 & 71 & 161 & \&\#113 & \\
\hline 18 & 120 & 022 DC2 & (device control 2) & 50 & 32 & 062 & \&\#50; & 2 & & 82 & & 122 & \&\#82; & ; R & 114 & 72 & 162 & \&\#114; & \\
\hline 19 & 130 & 023 DC3 & (device control 3) & 51 & 33 & 063 & \&\#51; & 3 & & 83 & 53 & 123 & ¢\#83; & ; & 115 & 73 & 163 & \&\#115; & \\
\hline 20 & 140 & 024 DC4 & (device control 4) & 52 & 34 & 064 & \&\#52; & 4 & & 84 & & 124 & \&\#84; & ; & 116 & 74 & 164 & \&\#116; & \\
\hline 21 & 150 & 025 NAK & (negative acknowledge) & 53 & 35 & 065 & \&\#53; & 5 & & 85 & 55 & 125 & \&\#85; & ; U & 117 & 75 & 165 & \&\#117 & \\
\hline 22 & 160 & 026 SYN & (synchronous idle) & 54 & 36 & 066 & \&\#54; & - & & 86 & 56 & 126 & Q\#86; & ; V & 118 & 76 & 166 & \&\#118; & \\
\hline 23 & 170 & 027 ETB & (end of trans. block) & 55 & 37 & 067 & \&\#55; & 7 & & 87 & 57 & 127 & \&\#87; & ; & 119 & 77 & 167 & \&\#119; & \\
\hline 24 & 18 & 030 CAN & (cancel) & 56 & 38 & 070 & \&\#56; & 8 & & 88 & 58 & 130 & \&\#88; & ; X & 120 & 78 & 170 & \&\#120; & \\
\hline 25 & 19 & 031 EM & (end of medium) & 57 & 39 & 071 & \&\#57; & 9 & & 89 & 59 & 131 & \&\#89; & ; Y & 121 & 79 & 171 & \&\#121 & \\
\hline 26 & 1A & 032 SUB & (substitute) & 58 & 3A & 072 & \&\#58; & : & & 90 & 5 A & 132 & \&\#90; & ; 2 & 122 & 7A & 172 & \&\#122; & \\
\hline 27 & 1B 0 & 033 ESC & (escape) & 59 & 3B & 073 & \&\#59; & , & & 91 & 5 B & 133 & \&\#91; & ; & 123 & 7 B & 173 & \&\#123; & \\
\hline 28 & 1C 0 & 034 FS & (file separator) & 60 & 3C & 074 & \&\#60; & < & & 92 & & 134 & \&\#92; & & 124 & 7 C & 174 & \&\#124; & \\
\hline 29 & 1D 0 & 035 GS & (group separator) & 61 & 3D & 075 & \&\#61; & = & & 93 & 5 D & 135 & \&\#93; & ; ] & 125 & 7D & 175 & \&\#125 & \\
\hline 30 & 1E & 036 RS & (record separator) & 62 & 3E & 076 & \&\#62; & \(>\) & & 94 & & 136 & \&\#94; & & 126 & 7 E & 176 & \&\#126; & \\
\hline 31 & \(1 F\) & 037 US & (unit separator) & 63 & 3 F & 077 & \&\#63; & ? & & 95 & 5 F & 137 & \&\#95; & & 127 & 7 F & 177 & \&\#127; & DEL \\
\hline
\end{tabular}

\section*{Extended ASCII Table}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline 128 & Ç & 144 & E & 160 & á & 176 & ＋ & 192 & L & 208 & \(\Perp\) & 224 & \(\alpha\) & 240 & 三 \\
\hline 129 & ü & 145 & ＊ & 161 & i & 177 & & 193 & \(\perp\) & 209 & ¢ & 225 & \(\beta\) & 241 & \(\pm\) \\
\hline 130 & é & 146 & 不 & 162 & о́ & 178 & & 194 & T & 210 & \(\pi\) & 226 & \(\Gamma\) & 242 & \(\geq\) \\
\hline 131 & à & 147 & 6 & 163 & ú & 179 & ｜ & 195 & F & 211 & 4 & 227 & \(\pi\) & 243 & \(\leq\) \\
\hline 132 & a & 148 & 0 & 164 & nit & 180 & － & 196 & － & 212 & t & 228 & \(\Sigma\) & 244 & 1 \\
\hline 133 & à & 149 & ò & 165 & Ñ & 181 & ＝ & 197 & \(t\) & 213 & F & 229 & \(\square\) & 245 & J \\
\hline 134 & a & 150 & ט̂ & 166 & a & 182 & － & 198 & ＝ & 214 & \(\pi\) & 230 & \(\mu\) & 246 & \(\div\) \\
\hline 135 & Ç & 151 & ù & 167 & － & 183 & \(\pi\) & 199 & 1 & 215 & \＃ & 231 & \(\tau\) & 247 & \(\approx\) \\
\hline 136 & ê & 152 & y & 168 & ¿ & 184 & 7 & 200 & L & 216 & キ & 232 & \(\Phi\) & 248 & 。 \\
\hline 137 & ë & 153 & O & 169 & － & 185 & \(\downarrow\) & 201 & 『 & 217 & 」 & 233 & （1） & 249 & \\
\hline 138 & è & 154 & Ü & 170 & ᄀ & 186 & \｜ & 202 & \(\xrightarrow{\Perp}\) & 218 & \(\Gamma\) & 234 & \(\Omega\) & 250 & \\
\hline 139 & 1 & 155 & \(\phi\) & 171 & 1／2 & 187 & ง & 203 & \(\bar{\pi}\) & 219 & \(\square\) & 235 & \(\delta\) & 251 & \(\checkmark\) \\
\hline 140 & i & 156 & E & 172 & 1／4 & 188 & 』 & 204 & 15 & 220 & ■ & 236 & \(\infty\) & 252 & \({ }^{\text {n }}\) \\
\hline 141 & 1 & 157 & ¥ & 173 & i & 189 & \(\Perp\) & 205 & ＝ & 221 & I & 237 & 中 & 253 & 2 \\
\hline 142 & A & 158 & P & 174 & ＜ & 190 & \(\pm\) & 206 & \＃ & 222 & I & 238 & \(\varepsilon\) & 254 & － \\
\hline 143 & \＆ & 159 & \(f\) & 175 & 》 & 191 & 7 & 207 & \(\stackrel{ }{1}\) & 223 & \(\square\) & 239 & \(\bigcirc\) & 255 & \\
\hline & & & & & & & & & & & our & www & Look & Tabl & \\
\hline
\end{tabular}

\section*{Casting}
- Casting is a method to correctly use variables of different types together
- It allows to treat a variable of one type as if it were of another type in a specific context
- When it makes sense, the compiler does it for us automatically
- Implicit (automatic)
\[
\begin{aligned}
& \text { int } x=1 ; \\
& \text { float } y=2.3 ; \\
& x=x+y ;
\end{aligned}
\]
x=3 compiler automatically casted (=converted) y to be an integer just for this instruction
- Explicit (non-automatic)
\begin{tabular}{ll} 
char \(\mathrm{c}=\mathrm{A} \mathrm{A}\) ' ; & \begin{tabular}{l} 
Explicit casting from char to int. The value of x \\
int \(\mathrm{x}=(\) int \() \mathrm{c} ;\) \\
here is 65
\end{tabular}
\end{tabular}

\section*{Operators}
- Assignment
- Arithmetic
* / \% + -
- Increment
++ -- += -=
- Relational
\(\ll=\gg=\) = !=
- Logical
- Bitwise
\& \| \| !
\& | ~ ^ << >>
- Comma
,

\title{
Operators - Assignment \\ int \(x=3\); \\ \[
\mathrm{x}=7
\]
}
int \(x_{0} \mathrm{y}=5 ; \longrightarrow\) The comma operator allows
\[
x=y=7 ;
\] us to perform multiple assignments/declarations
int \(\mathrm{i}, \mathrm{j}, \mathrm{k}\);
\(\mathrm{k}=(\mathrm{i}=2, \mathrm{j}=3)\);
printf( "i \(\left.=\% \mathrm{~d}, \mathrm{j}=\% \mathrm{~d}, \mathrm{k}=\%_{\mathrm{d}} \backslash \mathrm{n} ", \mathrm{i}, \mathrm{j}, \mathrm{k}\right)\);

\section*{Operators - Arithmetic}
- Arithmetic operators have a precedence
int \(x\);
\(x=3+5 * 2-4 / 2 ;\)
- We can use parentheses () to impose our precedence order int \(x\);
\(\mathrm{x}=(3+5) *(2-4) / 2\);
- \% returns the module (or the remainder of the division)
int \(x\);
\(x=5 \% 3 ; \quad / / x=2\)
- We have to be careful with integer vs. float division : remember automatic casting!
```

int x = 3;
float y;
y=x / 2; // y = 1.00

```
```

float y;

```
float y;
y=1 / 2; // y = 0.00
```

y=1 / 2; // y = 0.00

```

\section*{Operators - Arithmetic}
- Arithmetic operators have a precedence
int \(x\);
\(x=3+5 * 2-4 / 2 ;\)
- We can use parentheses () to impose our precedence order int \(x\);
\(\mathrm{x}=(3+5) *(2-4) / 2\);
- \% returns the module (or the remainder of the division)
int \(x\);
\(\mathrm{x}=5 \% 3 ; \quad / / \mathrm{x}=2\)
- We have to be careful with integer vs. float division : remember automatic casting!
int \(x=3\);
float \(y\);
\(y=x / 2 ; / / y=1.00\)
Possible fixes:
1) float \(x=3 ;\)
2)y \(=(\) float \() \times / 2 ;\)
Then \(y=1.50\)
```

float y;
y=1 / 2; // y = 0.00
Possible fix: y = 1.0/2;
Then y = 0.50

## Operators - Increment

$$
++\quad--\quad+=\quad-=
$$

int $x=3, y, z ;$
$\mathrm{x}++; \longrightarrow \mathrm{x}$ is incremented at the end of statement
$++\mathrm{x} ; \longrightarrow \mathrm{x}$ is incremented at the beginning of statement

$$
\mathrm{y}=++\mathrm{x}+3 ; / / \mathrm{x}=\mathrm{x}+1 ; \mathrm{y}=\mathrm{x}+3 ;
$$

$$
\mathrm{z}=\mathrm{x}+++3 ; / / \mathrm{z}=\mathrm{x}+3 ; \mathrm{x}=\mathrm{x}+1 ;
$$

$$
\mathrm{x}-=2 ; \quad / / \mathrm{x}=\mathrm{x}-2 ;
$$

## Operators - Relational \ll \ggg = = ! $=$

- Return 0 if statement is false, 1 if statement is true

$$
\begin{array}{ll}
\text { int } x=3, y=2, z, k, t \\
z=x>y ; & / / z=1 \\
\mathrm{k}=\mathrm{z}<=\mathrm{y} ; & / / \mathrm{k}=0 \\
\mathrm{t}=\mathrm{x}!=\mathrm{y} ; & / / \mathrm{t}=1
\end{array}
$$

## Operators - Logical \&\& II!

- A variable with value 0 is false, a variable with value $!=0$ is true

$$
\begin{aligned}
& \text { int } x=3, y=0, z, k, t, q=-3 \text {; } \\
& \mathrm{z}=\mathrm{x} \text { \&\& } \mathrm{y} ; \quad / / \mathrm{z}=0 ; \quad \mathrm{x} \text { is true but } \mathrm{y} \text { is false } \\
& \mathrm{k}=\mathrm{x} \text { || } \mathrm{y} ; \quad \text { // } \mathrm{k}=1 ; \quad \text { x is true } \\
& \mathrm{t}=\text { ! } \mathrm{q} ; \quad \quad / / \mathrm{t}=0 ; \quad \text { is true }
\end{aligned}
$$

## Review: Operators - Bitwise

- Work on the binary representation of data
- Remember: computers store and see data in binary format!

```
int x, y, z , t, q, s, v;
x = 3; 0000000000000000000000000000000011
y = 16; 00000000000000000000000000010000
z = x << 1; equivalent to z=x [ 21 00000000000000000000000000000110
t = y >> 3; equivalent to t=y 2 2-300000000000000000000000000000010
q = x & y;
    00000000000000000000000000000000
s = x | y;
    00000000000000000000000000010011
v = x {
    00000000000000000000000000010011
```


## printf

- printf is a function used to print to standard output (command line)
- Syntax:

```
printf("format1 format2 ...", variable1, variable2,...);
```

- Format characters:
- \%d or \%i integer
- \%f float
- \%lf double
- \%c char
- \%u unsigned
- \%s string


## printf

```
#include <stdio.h>
int main() {
int a,b;
float c,d;
a = 15;
b = a / 2;
printf("%d\n",b);
printf("%3d\n",b);
printf("%03d\n",b);
c = 15.3;
d = c / 3;
printf("%3.2f\n",d);
return(0);
```

Output:
\}

## printf

## Escape sequences

In newline
lt
Iv
\f
lb
Ir
tab
vertical tab
new page
backspace
carriage return

## Assignment

- Read PCP Chapter 3 and 4


## COMSW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

Lecture 4
Spring 2011
Instructor: Michele Merler

## Announcements

- HW 1 is due on Monday, February $14^{\text {th }}$ at the beginning of class, no exceptions
- Read so far: PCP Chapters 1 to 4
- Reading for next Wednesday: PCP Chapter 5


## Review - Access CUNIX

http://www1.cs.columbia.edu/~bert/courses/1003/cunix.html

1) Enable windowing environment

- X11, Xming, X-Server

2) Launch SSH session (login with UNI and password)

- Terminal, Putty

3) Launch Emacs
\$ emacs \&
4) Open/create a file, than save it with .c extension
5) Compile source code into executable with gcc

## Review - Compiling your C code

- GCC : GNU Compiler Collection
- When you invoke GCC, it normally does preprocessing, compilation, assembly and linking
- Basic Command
- gcc myProgram.c
- ./a.out

Run compiled program (executable)

- More advanced options
- gcc -Wall -o myProgram myProgram.c
- ./myProgram


## Review - Compiling your C code

- GCC : GNU Compiler Collection
- When you invoke GCC, it normally does preprocessing, compilation, assembly and linking
- Basic Command
- gcc myProgram.c
- ./a.out Run compiled program (executable)

Display all types of warnings, not only errors

```
Specify name of
the executable
```

- gcc -Wali -o myProgram myProgram.c
- ./myProgram Run compiled program (executable)


## Review: C Syntax

- Statements
- one line commands
- always end with ;
- can be grouped between \{ \}
- Comments
single line comment
/* multiple lines comments
*/


## Review : Variables and types

- Variables are placeholders for values

$$
\begin{aligned}
& \text { int } x=2 ; \\
& x=x+3 ; / / x \text { value is } 5 \text { now }
\end{aligned}
$$

- In C, variables are divided into types, according to how they are represented in memory (always represented in binary)
- int
- float
- double 8 bytes
- char

4 bytes, signed/unsigned
4 bytes, decimal part + exponent

1 byte, ASCII Table

## Review : Casting

- Casting is a method to correctly use variables of different types together
- It allows to treat a variable of one type as if it were of another type in a specific context
- When it makes sense, the compiler does it for us automatically
- Implicit (automatic)

$$
\begin{aligned}
& \text { int } x=1 ; \\
& \text { float y }=2.3 ; \\
& x=x+y ;
\end{aligned}
$$

```
x=3 compiler automatically casted
(=converted) y to be an integer just for
this instruction
```

- Explicit (non-automatic)

| char $c=$ 'A' $;$ | Explicit casting from char to int. The value of $x$ <br> int $x=($ int $)$ <br> here is 65 |
| :--- | :--- |

## Today

- Operators
- printf()
- Binary logic


## Operators

- Assignment
- Arithmetic
* / \% + -
- Increment
++ -- += -=
- Relational
\ll= \gg= == !=
- Logical
- Bitwise
\&\& || !
$\& \mid \sim \wedge \ll>$
- Comma


## Operators - Assignment and Comma

```
int x = 3;
x = 7;
int x
x = y = 7; us to perform multiple
int i,j,k;
k = (i=2, j=3);
printf("i = %d, j = %d, k = %d\n",i,j,k);
```

                                assignments/declarations
    
## Operators - Arithmetic

- Arithmetic operators have a precedence
int $x$;
$\mathrm{x}=3+5 * 2-4 / 2 ;$
- We can use parentheses () to impose our precedence order int $x$;
$\mathrm{x}=(3+5) *(2-4) / 2$;
- \% returns the module (or the remainder of the division)
int $x$;
$\mathrm{x}=5 \% 3 ; \quad / / \mathrm{x}=2$
- We have to be careful with integer vs. float division : remember automatic casting!

```
int x = 3;
float y;
y =x / 2; // y = 1.00
```

```
float y;
```

float y;
y=1 / 2; // y = 0.00

```
y=1 / 2; // y = 0.00
```


## Operators - Arithmetic

```
* / % + -
```

- Arithmetic operators have a precedence
int $x$;
$x=3+5 * 2-4 / 2 ;$
- We can use parentheses () to impose our precedence order int $x$;
$\mathrm{x}=(3+5) *(2-4) / 2$;
- \% returns the module (or the remainder of the division)
int $x$;
$\mathrm{x}=5 \% 3 ; \quad / / \mathrm{x}=2$
- We have to be careful with integer vs. float division : remember automatic casting!

```
int x = 3;
float y;
y =x / 2; // y = 1.00
```

```
Possible fixes:
```

Possible fixes:
1)float x = 3;
1)float x = 3;
2)y = (float) x /2;
2)y = (float) x /2;
Then y = 1.50

```
Then y = 1.50
```

```
float y;
y = 1 / 2; // y = 0.00
Possible fix: y = 1.0/2;
Then y = 0.50

\section*{Operators - Increment/Decrement}
\[
++\quad-\quad+=\quad-=
\]
int \(x=3, y, z\);
\(\mathrm{X}++; \longrightarrow \mathrm{x}\) is incremented at the end of statement
\(++x ; \longrightarrow x\) is incremented at the beginning of statement
\[
\mathrm{y}=++\mathrm{x}+3 ; / / \mathrm{x}=\mathrm{x}+1 ; \mathrm{y}=\mathrm{x}+3
\]
\[
\mathrm{z}=\mathrm{x}+++3 ; / / \mathrm{z}=\mathrm{x}+3 ; \mathrm{x}=\mathrm{x}+1
\]
\[
x-=2 ; \quad / / x=x-2
\]

\section*{Operators - Relational}
\[
\ll=\gg===\text { != }
\]
- Return 0 if statement is false, 1 if statement is true
\[
\begin{array}{ll}
\text { int } \mathrm{x}=3, \mathrm{y}=2, \mathrm{z}, \mathrm{k}, \mathrm{t} \\
\mathrm{z}=\mathrm{x}>\mathrm{y} ; & / / \mathrm{z}=1 \\
\mathrm{k}=\mathrm{x}<=\mathrm{y} ; & / / \mathrm{k}=0 \\
\mathrm{t}=\mathrm{x}!=\mathrm{y} ; & / / \mathrm{t}=1
\end{array}
\]

\section*{Operators - Logical \&\& || !}
- A variable with value 0 is false, a variable with value !=0 is true
\[
\begin{aligned}
& \text { int } x=3, y=0, z, k, t, q=-3 \text {; } \\
& \mathrm{z}=\mathrm{x} \text { \&\& } \mathrm{y} ; \quad / / \mathrm{z}=0 ; \quad \mathrm{x} \text { is true but } \mathrm{y} \text { is false } \\
& \mathrm{k}=\mathrm{x} \text { || } \mathrm{y} ; \quad \text { // } \mathrm{k}=1 ; \quad \text { x is true } \\
& \mathrm{t}=\text { ! } \mathrm{q} ; \quad \quad / / \mathrm{t}=0 ; \quad \text { is true }
\end{aligned}
\]

\section*{Operators - Bitwise}
- Work on the binary representation of data
- Remember: computers store and see data in binary format!
```

int x, y, z , t, q, s, v;
x = 3; 000000000000000000000000000000011
y = 16; 00000000000000000000000000010000
z = x << 1; equivalent to z=x 利 100000000000000000000000000000110
t = y >> 3; equivalent to t=y\cdot 2-3 00000000000000000000000000000010
q = x \& y; 00000000000000000000000000000000
s = x l y; 00000000000000000000000000010011
v = x
00000000000000000000000000010011

```

\section*{printf}
- printf is a function used to print to standard output (command line)
- Syntax:
```

printf("format1 format2 ...", variable1, variable2,...);

```
- Format characters:
- \%d or \%i integer
- \%f float
- \%lf double
- \%c char
- \%u unsigned
- \%s string

\section*{printf}
```

\#include <stdio.h>
int main() {
int a,b;
float c,d;
a = 15;
b = a / 2;
printf("%d\n",b);
printf("%3d\n",b);
printf("%03d\n",b);
c = 15.3;
d = c / 3;
printf("%3.2f\n",d); 5.10
return(0);

```

Output:

\section*{printf}

Escape sequences
In newline
It tab
\(\begin{array}{ll}\text { Iv } & \text { vertical tab } \\ \text { If } & \text { new page }\end{array}\)
lb backspace
Ir
carriage return

\section*{Binary Logic}
- 1 = true, 0 = false
- Decimal to binary conversion
\[
6_{10}=110_{2}
\]

\section*{Binary Logic}
- 1 = true, 0 = false
remainder
- Decimal to binary conversion

\(\longrightarrow\)\begin{tabular}{l|l}
\multicolumn{1}{c}{\begin{tabular}{l} 
remai \\
\(\mathbf{6}\) \\
\(\mathbf{6}\) \\
\hline
\end{tabular}} \\
\hline 3 & 1 \\
\hline 1 & 1 \\
\hline 0 &
\end{tabular}

\section*{Binary Logic}
- 1 = true, 0 = false

- Decimal to binary conversion

- Binary to decimal conversion
\(11001_{2}=1 \times 2^{0}+0 \times 2^{1}+0 \times 2^{2}+1 \times 2^{3}+1 \times 2^{4}=25\)

\section*{Binary Logic}
- 1 = true, 0 = false
- Decimal to binary conversion


Divide by 2

- Binary to decimal conversion
\[
11001_{2}=1 \times 2^{0}+0 \times 2^{1}+0 \times 2^{2}+1 \times 2^{3}+1 \times 2^{4}=25
\]
- AND \(v=x \& y\)
\begin{tabular}{|c|c|c|}
\hline \(\mathbf{x}\) & \(\mathbf{y}\) & \(\mathbf{v}\) \\
\hline 0 & 0 & 0 \\
\hline 0 & 1 & 0 \\
\hline 1 & 0 & 0 \\
\hline 1 & 1 & 1 \\
\hline
\end{tabular}
- OR
\(v=x \mid y\)
\begin{tabular}{|c|c|c|}
\hline\(x\) & \(y\) & \(v\) \\
\hline 0 & 0 & 0 \\
\hline 0 & 1 & 1 \\
\hline 1 & 0 & 1 \\
\hline 1 & 1 & 1 \\
\hline
\end{tabular}
- NOT
\(v=\) ! \(x\)

- EXOR \(v=x^{\wedge} y\)
\begin{tabular}{|c|c|c|}
\hline \(\mathbf{x}\) & \(\mathbf{y}\) & \(\mathbf{v}\) \\
\hline 0 & 0 & 0 \\
\hline 0 & 1 & 1 \\
\hline 1 & 0 & 1 \\
\hline 1 & 1 & 0 \\
\hline
\end{tabular}

\section*{Homework 1 review}

\section*{HOW TO COMPRESS/UNCOMPRESS folders in UNIX}
- Compress folder ~/COMS1003/HW1 to HW1.tar.gz tar -zcvf HW1.tar.gz ~/COMS1003/HW1
- Uncompress HW1.tar.gz to folder ~/COMS1003/HW1new tar -zxvf HW1.tar.gz -C ~/COMS1003/HW1new (note: ~/COMS1003/HW1new must exist already)

\section*{COMSW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

Lecture 5
Spring 2011
Instructor: Michele Merler

\section*{Announcements}
- Exercise 1 solution out
- Exercise 2 out
- Read PCP Ch 6

\section*{Today}
- Review of operators and printf()
- Binary Logic
- Arrays
- Strings

\section*{Review : printf}
- printf is a function used to print to standard output (command line)
- Syntax:
```

printf("format1 format2 ...", variable1, variable2,...);

```
- Format characters:
- \%d or \%i integer
- \%f float
- \%lf double
- \%c char
- \%u unsigned
- \%s string

\section*{Review : printf}
```

\#include <stdio.h>
int main() {
int a,b;
float c,d;
a = 15;
b = a / 2;
printf("%d\n",b) ;
printf("%3d\n",b) ;
printf("%03d\n",b);
c = 15.3;
d = c / 3;
printf("%3.2f\n",d);
return(0);
7
7
0 0 7
5 . 1 0

```

\section*{Output:}

\section*{Review : printf}

Escape sequences
In newline
It tab
\(\begin{array}{ll}\text { Iv } & \text { vertical tab } \\ \text { If } & \text { new page }\end{array}\)
lb backspace
Ir
carriage return

\section*{Binary Logic}
- In binary logic, variables can have only 2 values:
- True ( commonly associated with 1 )
- False ( commonly associated with 0 )
- Binary Operations are defined through TRUTH TABLES
\begin{tabular}{|c|c|c|c|}
\hline AND & X & Y & v \\
\hline \(v=x \& y\) & 0 & 0 & 0 \\
\hline & 0 & 1 & 0 \\
\hline & 1 & 0 & 0 \\
\hline & 1 & 1 & 1 \\
\hline \[
\mathrm{OR}
\] & X & Y & v \\
\hline \multirow[t]{4}{*}{\(v=x \mid y\)} & 0 & 0 & 0 \\
\hline & 0 & 1 & 1 \\
\hline & 1 & 0 & 1 \\
\hline & 1 & 1 & 1 \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline NOT & \(x\) & \(v\) \\
\hline\(v=!x\) & 0 & 1 \\
\hline & 1 & 0 \\
\hline
\end{tabular}
\begin{tabular}{l|c|c|c|}
\hline EXOR & \(x\) & \(y\) & \(v\) \\
\hline\(v=x^{\wedge} y\) & 0 & 0 & 0 \\
\hline 0 & 1 & 1 \\
\hline 1 & 0 & 1 \\
\hline 1 & 1 & 0 \\
\hline
\end{tabular}

\section*{Binary Logic}
- 1 = true, 0 = false
- Decimal to binary conversion
\[
6_{10}=110_{2}
\]

\section*{Binary Logic}
- 1 = true, 0 = false
remainder
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\(\longrightarrow\)\begin{tabular}{l|l}
\multicolumn{1}{c}{\begin{tabular}{l} 
remai \\
\(\mathbf{6}\) \\
\(\mathbf{6}\) \\
\hline
\end{tabular}} \\
\hline 3 & 1 \\
\hline 1 & 1 \\
\hline 0 &
\end{tabular}

\section*{Binary Logic}
- 1 = true, 0 = false

- Decimal to binary conversion

- Binary to decimal conversion
\(11001_{2}=1 \times 2^{0}+0 \times 2^{1}+0 \times 2^{2}+1 \times 2^{3}+1 \times 2^{4}=25\)

\section*{Binary Logic}
- 1 = true, 0 = false
- Decimal to binary conversion


Divide by 2
\(\longrightarrow\)\begin{tabular}{l|l} 
\\
\\
\hline \(\mathbf{6}\) & 0 \\
\hline 3 & 1 \\
\hline 1 & 1 \\
\hline 0 &
\end{tabular}
- Binary to decimal conversion
\[
11001_{2}=1 \times 2^{0}+0 \times 2^{1}+0 \times 2^{2}+1 \times 2^{3}+1 \times 2^{4}=25
\]
- AND \(v=x \& y\)
\begin{tabular}{|c|c|c|}
\hline \(\mathbf{x}\) & \(\mathbf{y}\) & \(\mathbf{v}\) \\
\hline 0 & 0 & 0 \\
\hline 0 & 1 & 0 \\
\hline 1 & 0 & 0 \\
\hline 1 & 1 & 1 \\
\hline
\end{tabular}
- NOT \(\mathrm{v}=\) ! x

- EXOR \(v=x^{\wedge} y\)
\begin{tabular}{|c|c|c|}
\hline \(\mathbf{x}\) & \(\mathbf{y}\) & \(\mathbf{v}\) \\
\hline 0 & 0 & 0 \\
\hline 0 & 1 & 1 \\
\hline 1 & 0 & 1 \\
\hline 1 & 1 & 0 \\
\hline
\end{tabular}

\section*{Review: Operators}
- Assignment
- Arithmetic
- Increment
- Relational
- Logical
- Bitwise
- Comma


\section*{Operators - Bitwise}
- Work on the binary representation of data
- Remember: computers store and see data in binary format!
```

int x, y, z , t, q, s, v;
x = 3; 000000000000000000000000000000011
y = 16; 00000000000000000000000000010000
z = x << 1; equivalent to z=x 利 100000000000000000000000000000110
t = y >> 3; equivalent to t=y\cdot 2-3 00000000000000000000000000000010
q = x \& y; 00000000000000000000000000000000
s = x l y; 00000000000000000000000000010011
v = x
00000000000000000000000000010011

```

\section*{Operators - Arithmetic}
* / \% + -
- Arithmetic operators have a precedence
int \(x\);
\(x=3+5 * 2-4 / 2 ;\)
- We can use parentheses () to impose our precedence order int \(x\);
\(\mathrm{x}=(3+5) *(2-4) / 2\);
- \% returns the module (or the remainder of the division)
int \(x\);
\(x=5 \% 3 ; \quad / / x=2\)
- We have to be careful with integer vs. float division : remember automatic casting!
```

int x = 3;
float y;
y=x / 2; // y = 1.00

```
```

Possible fixes:

```
Possible fixes:
1)float x = 3;
1)float x = 3;
2)y = (float) x /2;
2)y = (float) x /2;
Then y = 1.50
```

Then y = 1.50

```
```

float y;
y = 1 / 2; // y = 0.00
Possible fix: y = 1.0/2;
Then y = 0.50

## Operators - Increment/Decrement

$$
++\quad-\quad+=\quad-=
$$

int $x=3, y, z$;
$\mathrm{X}++; \longrightarrow \mathrm{x}$ is incremented at the end of statement
$++x ; \longrightarrow x$ is incremented at the beginning of statement

$$
\mathrm{y}=++\mathrm{x}+3 ; / / \mathrm{x}=\mathrm{x}+1 ; \mathrm{y}=\mathrm{x}+3
$$

$$
\mathrm{z}=\mathrm{x}+++3 ; / / \mathrm{z}=\mathrm{x}+3 ; \mathrm{x}=\mathrm{x}+1
$$

$$
x-=2 ; \quad / / x=x-2
$$

## Operators - Relational

$$
\ll=\gg===\text { != }
$$

- Return 0 if statement is false, 1 if statement is true

$$
\begin{array}{ll}
\text { int } \mathrm{x}=3, \mathrm{y}=2, \mathrm{z}, \mathrm{k}, \mathrm{t} \\
\mathrm{z}=\mathrm{x}>\mathrm{y} ; & / / \mathrm{z}=1 \\
\mathrm{k}=\mathrm{x}<=\mathrm{y} ; & / / \mathrm{k}=0 \\
\mathrm{t}=\mathrm{x}!=\mathrm{y} ; & / / \mathrm{t}=1
\end{array}
$$

## Operators - Logical \&\& || !

- A variable with value 0 is false, a variable with value !=0 is true

$$
\begin{aligned}
& \text { int } x=3, y=0, z, k, t, q=-3 \text {; } \\
& z=x \text { \&\& } y ; \quad / / z=0 ; \quad x \text { is true but } y \text { is false } \\
& \mathrm{k}=\mathrm{x} \text { || } \mathrm{y} ; \quad \text { // } \mathrm{k}=1 ; \quad \text { x is true } \\
& \mathrm{t}=\text { ! } \mathrm{q} ; \quad \quad / / \mathrm{t}=0 ; \quad \text { is true }
\end{aligned}
$$

## Arrays

- "A set of consecutive memory locations used to store data" [PCP, Ch 5]
int $\mathrm{X}[4] ; \mathrm{/}$ a vector containing 4 integers

- Indexing starts at 0 !

$$
\begin{aligned}
& \mathrm{X}[0]=3 ; \\
& \mathrm{X}[2]=7 ;
\end{aligned}
$$

- Be careful not to access uninitialized elements!

$$
\text { int } c=x[7] ;
$$

gcc will not complain about this, but the value of $x$ is going to be random!

## Arrays

- Multidimensional arrays

```
int arr[4][3]; // a matrix containing 4x3 = 12 integers
```

| $\operatorname{arr}[0][0]$ | $\operatorname{arr}[0][1]$ | $\operatorname{arr}[0][2]$ |
| :---: | :---: | :---: |
| $\operatorname{arr}[1][0]$ | $\operatorname{arr}[1][1]$ | $\operatorname{arr}[1][2]$ |
| $\operatorname{arr}[2][0]$ | $\operatorname{arr}[2][1]$ | $\operatorname{arr}[2][2]$ |
| $\operatorname{arr}[3][0]$ | $\operatorname{arr}[3][1]$ | $\operatorname{arr}[3][2]$ |

- Indexing starts at 0 !

$$
\begin{aligned}
& \operatorname{arr}[0][0]=1 ; \\
& \operatorname{arr}[3][1]=7 ;
\end{aligned}
$$

- Initialize arrays
int $\mathrm{X}[4]=\{3,6,7,89\}$;
int $Y[2][4]=\{\{19,2,6,99\},\{55,5,555,0\}\}$;
int $\operatorname{Arr}[]=\{3,6,77\}$; This automatically allocates memory for an array of 3 integers


## Strings

- Strings are arrays of char
- $\backslash 0$ ' is a special character that indicates the end of a string

We need 6 characters because there is ' $\backslash 0$ '

| ' H ' | 'e' | '1' | ' 1 | 'o' | ' 10 ' |
| :---: | :---: | :---: | :---: | :---: | :---: |

char s[10] = "Hello";

|  | ' ${ }^{\prime}$ ' | 'e' | '1 | '1' | 'o' | ' 10 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

$$
\begin{aligned}
& \text { char } s[6] ; \\
& s[0]=~ ' H^{\prime} ; \\
& s[1]=~ ' e^{\prime} ; \\
& s[2]=' l^{\prime} ; \\
& s[3]=' l^{\prime} ; \\
& s[4]=' o^{\prime} ; \\
& s[5]=' 10^{\prime} ;
\end{aligned}
$$

- Difference between string and char



## Strings functions

String specific functions are included in the library string.h

```
#include <string.h>
```

char s[6];
s = "Hello";

Illegal! String assignment can be done only at declaration!

- $\operatorname{strcpy}()$ : copy a string to another

```
strcpy( string1 , string2 ); Copy string2 to string1
```

char s[6];
strcpy(s, "Hello");

## String functions

String specific functions are included in the library string.h

- $\operatorname{strcmp}()$ : compare two strings

```
strcmp( string1 , string2 );
    Returns:
    0 if string1 and string2 are the same
    value != 0 otherwise
char s1[] = "Hi";
char s2[] = "Him";
char s3[3];
strcpy( s3, s1 );
int x = strcmp( s1, s2 ); // x != 0
int y = strcmp( s1, s3 ); // y = 0
```


## Strings functions

String specific functions are included in the library string.h

- strcat() : concatenate two strings

```
strcat( string1 , string2);
char s1[] = "Hello ";
char s2[] = "World!";
strcat(s1, s2);
```

| ' H ' | 'e' | ' 1 | ' 1 | 'o' | "، | 'W' | 'o' | 'r' | ' 1 | 'd' | ' $\backslash 0$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

- strlen() : returns the length of a string (does not count ' $\backslash 0$ ')

```
strlen( string );
char s1[] = "Hello";
int x = strlen(s1);
```

    \(/ / x=5\)
    
## Reading Strings

## Use functions from library stdio.h

- fgets() : get string from standard input (command line)

char s1[100];
fgets( s1, sizeof(s1), stdin);

Reads a maximum of sizeof(name) characters of a string from stdin and saves them into string name

NOTE: fgets() reads the newline character ' n ', so we should substitute it with ' $\backslash 0$ ';
name[strlen(name)-1] = ' $0^{\prime}$ ';

| 'H' | 'e' | '1' | '1' | 'o' | ' $\mathrm{n}^{\prime}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 'H' | 'e' | '1' | ' 1 ' | 'o' | ' $0^{\prime}$ |

- sizeof() : returns the size (number of bytes occupied in memory) of a variable (for strings it counts the number of elements, including ' $\backslash 0$ ')


## Reading numbers - Option 1

- First, read a string
- Then, convert string to number
- sscanf() : get string from standard input (command line)

```
sscanf( string, "format", &var1, ..., &varN);
```

char s1[100];
int $x, y ;$
printf("Please enter two numbers separated by a space\n")
fgets( s1, sizeof(s1), stdin);
User enters: 318
sscanf( s1, "\%d \%d", \&x, \&y );
$/ / x=3 ; y=18 ;$

## Reading numbers - Option 2

- Read directly the number
- $\operatorname{scanf()~:~get~string~from~standard~input~(command~line)~and~}$ automatically convert into a number

```
scanf( "format", &var1, ..., &varN);
int x, Y;
printf("Please enter two numbers separated by a space\n")
```

User enters: 318
scanf( "\%d \%d", \&x, \&y );
// $x=3 ; y=18 ;$

## Strings functions - recap

char s1[] = "Hello"; char s2[] = "He"; int x; char c;

- $\quad$ strcmp( $s 1, s 2)$
- $\operatorname{strcpy}(\mathrm{s} 1, \mathrm{~s} 2)$
- strcat( s1, s2)
- strlen( s )
- sizeof( s )
- fgets( $s$, sizeof(s1), stdin) fgets( s1, sizeof(s1), stdin); User enters "7R"
- sscanf( s, "\%d", \&var)

$$
\begin{aligned}
& \mathrm{x}=\operatorname{strcmp}(\mathrm{s} 1, \mathrm{~s} 2) / / \mathrm{x} \text { != } \\
& \text { strcpy( s2, s1 ); // s2 = "Hello" }
\end{aligned}
$$

strcat( s2, s1 ); //s2 = "HelloHello"

$$
x=\operatorname{strlen}(s 1) ; \quad / / x=5 ;
$$

$$
x=\text { sizeof(s1); } \quad / / x=6 ;
$$

// $x=7$; $c=~ ' R ' ;$

Read PCP Ch 6

## Homework 1 review

## HOW TO COMPRESS/UNCOMPRESS folders in UNIX

- Compress folder ~/COMS1003/HW1 to HW1.tar.gz tar -zcvf HW1.tar.gz ~/COMS1003/HW1
- Uncompress HW1.tar.gz to folder ~/COMS1003/HW1new tar -zxvf HW1.tar.gz -C ~/COMS1003/HW1new (note: ~/COMS1003/HW1new must exist already)


## COMsW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

Lecture 6
Spring 2011
Instructor: Michele Merler

## Announcements

## Homework 1 is due next Monday

## Exercise 2 is out

## Today

- Strings
- Control Flow
- Loops (if time permits)


## Review - arrays

- Multidimensional arrays
int $\mathrm{X}[4][3] ; / /$ a matrix containing $4 \times 3=12$ integers

| $X[0][0]$ | $X[0][1]$ | $X[0][2]$ |
| :---: | :---: | :---: |
| $X[1][0]$ | $X[1][1]$ | $X[1][2]$ |
| $X[2][0]$ | $X[2][1]$ | $X[2][2]$ |
| $X[3][0]$ | $X[3][1]$ | $X[3][2]$ |

- Indexing starts at 0 !
$\mathrm{X}[0][0]=1 ;$
$\mathrm{X}[3][1]=7$.
$\mathrm{X}[3][1]=7$;
- Initialize says
int $\operatorname{arr}[4]=\{3,6,7,89\}$;
int $\operatorname{arr} 2[2][4]=\{\{19,2,6,99\},\{55,5,555,0\}\}$;
int $\operatorname{arr}[]=\{3,6,77\}$; This automatically allocates memory for an array of 3 integers


## Strings

- Strings are arrays of char
- $\backslash 0$ ' is a special character that indicates the end of a string

We need 6 characters because there is ' $\backslash 0$ '

| ' H ' | 'e' | '1' | ' 1 | 'o' | ' 10 ' |
| :---: | :---: | :---: | :---: | :---: | :---: |

char s[10] = "Hello";

|  | ' ${ }^{\prime}$ ' | 'e' | '1 | '1' | 'o' | ' 10 |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

$$
\begin{aligned}
& \text { char } s[6] ; \\
& s[0]=~ ' H^{\prime} ; \\
& s[1]=~ ' e^{\prime} ; \\
& s[2]=' l^{\prime} ; \\
& s[3]=' l^{\prime} ; \\
& s[4]=' o^{\prime} ; \\
& s[5]=' 10^{\prime} ;
\end{aligned}
$$

- Difference between string and char



## Strings functions

String specific functions are included in the library string.h

```
#include <string.h>
```

char s[6];
s = "Hello";

Illegal! String assignment can be done only at declaration!

- $\operatorname{strcpy}()$ : copy a string to another

```
strcpy( string1 , string2 ); Copy string2 to string1
```

char s[6];
strcpy(s, "Hello");

## String functions

String specific functions are included in the library string.h

- $\operatorname{strcmp}()$ : compare two strings

```
strcmp( string1 , string2 );
    Returns:
    0 if string1 and string2 are the same
    value != 0 otherwise
char s1[] = "Hi";
char s2[] = "Him";
char s3[3];
strcpy( s3, s1 );
int x = strcmp( s1, s2 ); // x != 0
int y = strcmp( s1, s3 ); // y = 0
```


## Strings functions

String specific functions are included in the library string.h

- strcat() : concatenate two strings

```
strcat( string1 , string2);
char s1[] = "Hello ";
char s2[] = "World!";
strcat(s1, s2);
```

| ' H ' | 'e' | ' 1 | ' 1 | 'o' | "، | 'W' | 'o' | 'r' | ' 1 | 'd' | ' $\backslash 0$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |

- strlen() : returns the length of a string (does not count ' $\backslash 0$ ')

```
strlen( string );
char s1[] = "Hello";
int x = strlen(s1);
```

    \(/ / x=5\)
    
## Reading Strings

## Use functions from library stdio.h

- fgets() : get string from standard input (command line)

char s1[100];
fgets( s1, sizeof(s1), stdin);

Reads a maximum of sizeof(name) characters of a string from stdin and saves them into string name

NOTE: fgets() reads the newline character ' n ', so we should substitute it with ' $\backslash 0$ ';
s1[strlen(s1)-1] = ' $\mathbf{0}^{\prime}$;

| 'H' | 'e' | ' 1 ' | ' 1 ' | 'o' | ' $\mathrm{n}^{\prime}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 'H' | 'e' | '1' | '1' | 'o' | ' $0^{\prime}$ |

- sizeof() : returns the size (number of bytes occupied in memory) of a variable (for strings it counts the number of elements, including ' $\backslash 0$ ')


## Reading numbers - Option 1

- First, read a string
- Then, convert string to number
- sscanf() : get string from standard input (command line)

```
sscanf( string, "format", &var1, ..., &varN);
```

char s1[100];
int $x, y ;$
printf("Please enter two numbers separated by a space\n")
fgets( s1, sizeof(s1), stdin);
User enters: 318
sscanf( s1, "\%d \%d", \&x, \&y );
$/ / x=3 ; y=18 ;$

## Reading numbers - Option 2

- Read directly the number
- $\operatorname{scanf()~:~get~string~from~standard~input~(command~line)~and~}$ automatically convert into a number

```
scanf( "format", &var1, ..., &varN);
int x, Y;
printf("Please enter two numbers separated by a space\n")
```

User enters: 318
scanf( "\%d \%d", \&x, \&y );
// $x=3 ; y=18 ;$

## Strings functions - recap

char s1[] = "Hello"; char s2[] = "He"; int x; char c;

- $\quad$ strcmp( $s 1, s 2)$
- $\operatorname{strcpy}(\mathrm{s} 1, \mathrm{~s} 2)$
- strcat( s1, s2)
- strlen( s )
- sizeof( s )
- fgets( $s$, sizeof(s1), stdin) fgets( s1, sizeof(s1), stdin);
- sscanf( s, "\%d", \&var)
strcat( s2, s1 ); //s2 = "HelloHello"

$$
x=\operatorname{strlen}(s 1) ; \quad / / x=5 ;
$$

$$
x=\text { sizeof(s1); } \quad / / x=6 ;
$$ User enters "7R"

$$
\begin{aligned}
& \mathrm{x}=\operatorname{strcmp}(\mathrm{s} 1, \mathrm{~s} 2) / / \mathrm{x} \text { != } \\
& \operatorname{strcpy}(\mathrm{s} 2, \mathrm{~s} 1) ; / / \mathrm{s} 2=\text { "Hello" }
\end{aligned}
$$

sscanf( s1, "\%d\%c", \&x, \&c); // $x=7 ; c=$ 'R';

## Example - sumNums.c

## Control Flow

- So far we have seen linear programs, statements are executed in the order in which they are written
- What if we want to skip some instructions, or execute them only under certain conditions?
- Solution: control flow


## Control flow - General syntax

\(\left.\begin{array}{|cc|}keyword ( condition )\{ <br>
body statement 1; <br>
\vdots <br>
body statement n; <br>

\}\end{array}\right\}\)| The body is executed |
| :--- |
| only if the condition |
| is true! |

If the body of the control flow has only one statement, we can optionally not use the \{ \}

```
keyword ( condition )
    body statement 1;
```


## Control flow - if

- To execute a particular body of statements only if a particular condition is satisfied



## Example

```
int x = 3, y;
if ( x > 2 ) {
    x++;
\[
y=x
\]
```

\}
printf("y = \%d\n",y);

## Control flow - else

- To execute a particular body of statements only if a particular condition is not satisfied

```
if ( condition ) {
    body statement 1;
        .
    body statement n;
}
else {
    body statement 1;
    body statement m;
}
```


## Example

```
int \(x=3, y\);
if ( \(x\) > 2 ) \{
    x++;
    \(y=x ;\)
\}
else \{
    \(y=2 * x ;\)
\}
printf("y = \%d\n",y);

\section*{Control Flow - if/else example}
\[
\begin{aligned}
& \text { int } x=3, y=1 ; \\
& \text { if( } x>2) \\
& \text { if( } x=4) \\
& y=x \\
& \text { else } y \\
& y=2 * x \\
& \text { printf("y } \left.=\% d \backslash n^{\prime \prime}, y\right) ;
\end{aligned}
\]

\section*{Control Flow - if/else example}
```

int }x=3,y=1
if( x > 2 )
else
y = 2 * x;
printf("y = %d\n",y);

```

\section*{Control Flow - if/else example}
```

int x = 3, y = 1;
if( x > 2 ) {
if( x == 4) {
y = x;
}
else {
y = 2 * x;
}
}
printf("Y = %d\n",y);

```

This is why we need brackets and indentation!

\section*{Control Flow - if/else example}
```

int x = 3, y = 1;
if( x > 2 ) {
if( x == 4) {
y = x;
}
}
else {
y = 2 * x;
}
printf("y = %d\n",y);

```

\section*{Control flow - Switch}

Equivalent to a series of if/else statements

```

int i,j;
switch( i ) {
case 1:
j = i + 1;
break;
case 10:
j = i - 1;
default:
j = 1;
}

```

\section*{Control flow - Switch}

Equivalent to a series of if/else statements
\begin{tabular}{|c|c|}
\hline\(i\) & \(j\) \\
\hline 1 & 2 \\
\hline 10 & 1 \\
\hline \begin{tabular}{c} 
Any other \\
number
\end{tabular} & 1 \\
\hline
\end{tabular}
int i,j;
int i,j;
switch ( i ) \{
switch ( i ) \{
    case 1:
    case 1:
        \(j=i+1 ;\)
        \(j=i+1 ;\)
        break;
        break;
    case 10:
    case 10:
        \(j=i-1 ;\)
        \(j=i-1 ;\)
    default:
    default:
        \(j=1 ;\)
        \(j=1 ;\)
\}
\}

\section*{Control flow - Switch}

Equivalent to a series of if/else statements


If variable has value different
```

int i,j;
switch( i ) {
case 1:
j = i + 1;
break;
case 10:
j = i - 1;
default:
j = 1;
After last case I can
avoid using break

``` from all other cases

\section*{Switch}

Equivalent to a series of if/else statements


\section*{Control Flow - Loops}
- What if we want to perform the same operation multiple times?
- Example: we want to initialize all elements in a 100 dimensional array of integers to the value 7
```

int arr[100];
arr[0] = 7;
arr[1] = 7;
arr[2] = 7;
arr[3] = 7;
.
arr[99] = 7;

```

\section*{Loops - while}
- To execute a particular body of statements only until a particular condition is satisfied

```

Example
int $i=0$;
int arr[100];
while( i < 100 ) \{
$\operatorname{arr}[i]=7$;
i++;
\}

```

\section*{Loops - do/while}
- First execute body statements, then check if condition is satisfied
\begin{tabular}{cc|}
\hline do \(\{\) \\
body statement \(1 ;\) \\
\(\vdots\) \\
body statement \(\mathrm{n} ;\) \\
\} while ( condition ); \\
\end{tabular}
\begin{tabular}{|c|}
\hline Example \\
\hline \[
\begin{aligned}
& \text { int } i=10 \\
& \text { int } j=0 ;
\end{aligned}
\] \\
\hline ```
while( i < 10 )
{
``` \\
\hline \[
\begin{aligned}
& j++; \\
& i++;
\end{aligned}
\] \\
\hline \} \\
\hline
\end{tabular}

Example
int \(i=10 ;\)
int \(j=0 ;\)
do
\(\left\{\begin{array}{r}j++; \\ i++;\end{array}\right.\)
\(\}\) while( \(i<10) ;\)

\section*{Loops - do/while}
- First execute body of statements, then check if condition is satisfied

\begin{tabular}{|c|}
\hline Example \\
\hline \[
\begin{aligned}
& \text { int } i=10 \\
& \text { int } j=0 ;
\end{aligned}
\] \\
\hline ```
while( i < 10 )
{
``` \\
\hline \[
\begin{aligned}
& j++; \\
& i++;
\end{aligned}
\] \\
\hline \} \\
\hline
\end{tabular}
\[
j=0
\]

Example
\[
\begin{aligned}
& \text { int } i=10 ; \\
& \text { int } j=0 ; \\
& \text { do } \\
& \left\{\begin{array}{r}
j++; \\
i++;
\end{array}\right. \\
& \} \text { while }(i<10) ;
\end{aligned}
\]
\[
j=1
\]

\section*{Loops - break}
- To interrupt a loop once a certain condition different from the one in the loop declaration

```

Example
int i = 0;
char s[10] = "hi";
while( i < 10 )
{
if(s[i]=='\0')
break;
printf("%c",s[i]);
i++;

```

\section*{Loops - continue}
- To ignore the following instructions in a loop


\section*{Example}
```

int i = 0, sum = 0;
int s[3] = {7, 5, 9};
while( i < 3 )
{
if(s[i] < 6)
continue;
sum += s[i];
}

```

\section*{break vs. continue}


\section*{break vs. continue}
\begin{tabular}{|c|c|}
\hline \[
\begin{aligned}
& \text { int } x=0, y=0 \text {; } \\
& \text { while }(x<10) ~
\end{aligned}
\] & \[
\begin{aligned}
& \text { int } x=0, y=0 \\
& \text { while }(x<10)
\end{aligned}
\] \\
\hline x++; & x++; \\
\hline if ( \(\mathrm{x}=3\) ) \{ & if \((x==3) ~\{\) \\
\hline continue & break; \\
\hline \} & \} \\
\hline y++; & y++; \\
\hline \(y=9\) & \(y=2\) \\
\hline
\end{tabular}

\section*{Loops - for}
```

for (initial state ; condition ; state change ) {
body statement 1;
.
.
body statement n;
}

```

\section*{Example}
int i;
int arr[100];
for ( i = 0; i < 100 ; i++ ) \{
arr[i] = 7;
```

while( i < 100 ) \{
int i = 0;
int arr[100];
arr[i] = 7;
i++;

```

\section*{Homework 1 review}

\section*{HOW TO COMPRESS/UNCOMPRESS folders in UNIX}
- Compress folder ~/COMS1003/HW1 to HW1.tar.gz tar -zcvf HW1.tar.gz ~/COMS1003/HW1
- Uncompress HW1.tar.gz to folder ~/COMS1003/HW1new tar -zxvf HW1.tar.gz -C ~/COMS1003/HW1new (note: ~/COMS1003/HW1new must exist already)

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

Lecture 7
Spring2011
Instructor: Michele Merler

\section*{Today}
- Loops (from Lec6)
- Scope of variables
- Functions

\section*{Scope of Variables}
- Scope is the portion of program in which a variable is valid
- Depends on where the variable is declared
- Variables can be
- Global : valid everywhere
- Local : valid in a specific portion of the program included in \{ \}

\section*{Scope of Variables}
- Scope is the portion of program in which a variable is valid
- Depends on where the variable is declared
- Variables can be = Global : valid everywhere
- Local : valid in a specific portion of the program included in \{ \}
```

\#include <stdio.h>
double x = 3; /* global variable */
int main() {
double y = 7.2;
if( x > 2){
double z = x / 2;
}
return(0);
}

```


\section*{Scope of variables}
```

\#include <stdio.h>
double z = 1;
int main() {
printf("z1 = %lf\n", z); // z1 = 1.0000000
double z = 7;
if( z > 2){
double z = 0.5;
printf("z2 = %lf\n", z); // z2 = 0.5000000
}
printf("z3 = %lf\n", z); // z3 = 7.0000000
{
double z = 11;
printf("z4 = %lf\n",z); // z4 = 11.0000000
}
printf("z5 = %lf\n",z); // z5 = 7.0000000
return(0);
}

```

\section*{Scope of variables}
```

\#include <stdio.h>
double z = 1;
int main() {
printf("z1 = %lf\n", z); // z1 = 1.0000000
double z = 7;
if( z > 2){
double z = 0.5;
printf("z2 = %lf\n", z); // z2 = 0.5000000
}
printf("z3 = %lf\n", z); // z3 = 7.0000000
{
double z = 11;
printf("z4 = %lf\n",z); // z4 = 11.0000000
}
printf("z5 = %lf\n",z); // z5 = 7.0000000
return(0);
}

```

\section*{Class of Variables}
- A variable can be either
- Temporary : allocated in stack at beginning of block (if too many local variables allocated, stack overflow)
- Permanent : allocated before the program starts
- Global variables are always permanent
- Local variables are temporary unless they are declared static


\section*{Variables - Scope and Class}
\begin{tabular}{|l|l|l|l|}
\hline Declared & Scope & Class & initialized \\
\hline Outside all blocks & Global & Permanent & Once \\
\hline Static outside all blocks & Global & Permanent & Once \\
\hline Inside a block & Local & Temporary & Each time block is entered \\
\hline Static inside a block & Local & Permanent & Once \\
\hline
\end{tabular}
```

\#include <stdio.h>
int z = 0;
static int b;
int main() {
int g=0;
while( z < 3){
int y = 0;
static int }\textrm{x}=0\mathrm{ ;
Y++;
x++;
z++;
printf("x = %d, y = %d, z = %d\n", x, y, z);
}
return(0);

```
                \(x=1, y=1, z=1\)
\(x=2, y=1, z=2\)
\(x=3, y=1, z=3\)
                            y is initialized every time

\section*{Functions}
- Functions allow to write and reuse pieces of code that accomplish a task
- Help keeping large codes ordered


\section*{Functions - Example}

The function sumTwoNumbers takes two numbers as input and returns their sum.
double sumTwoNumbers( double n1, double n2 ) \{ double s;
\(s=n 1+n 2 ;\)
return (s);
\}

\section*{Functions - Example}

The function sumTwoNumbers takes two numbers as input and returns their sum.

Returned type must be consistent!


\section*{Functions - Example}
```

\#include <stdio.h>
double sumTwoNumbers( double n1, double n2 ){
double s;
n1++;
s = n1 + n2;
return(s);
}
int main() {
double x, y, z;
x = 2;
y = 2;
z = sumTwoNumbers(x, y);
printf("%f + %f = %f\n", x, y, z);
return(0);

```

\section*{Functions - void}
- If a function does not take any input
- If a function does not return any value
```

/* function to print an arrow to command line */
void printArrow(void){
/* function body */
return;
}
/* function to print multiple arrows to command line */
void printMultipleArrows(int nTimes){
int i;
for(i = 0; i < nTimes; i++){
printArrow();
}
return;
}
int main() {
int x = 3;
printMultipleArrows(x);
return(0);

## Functions - void

- If a function does not take any input
- If a function does not return any value



## Functions - void

- If a function does not take any input
- If a function does not return any value

Function is declared before being used

```
* function to print an arrow to comman
void printArrow(void){
/* function body */
    return;
    }
    /* function to print multiple arrows to command line */
    void printMultipleArrows(int nTimes){
        int i;
        for(i = 0; i < nTimes; i++){
} printArrow();
    return;
}
int main() {
    int x = 3;
    printMultipleArrows(x);
    return(0);
```


## Functions - Passing Arrays

```
/* function to compute the length of a string*/
int length( char s[] ){
    int size = 0;
    while(s[size] != '\0'){
        size++;
        }
    return size;
}
/* function to copy a string*/
    char[] copyString( char s[] ){
            char s2[100];
            strcpy(s2, s);
            return s2;

\section*{Functions - Passing Arrays}
```

/* function to compute the length of a string*/
int length( char s[] ){
int size = 0;
while(s[size] != '\0'){
size++;
}
return size;
}
function to copy a string*/
copyString( char s[] ){
char s2[100];
strcpy(s2, s);
return s2;

## Functions - exit()

exit () is used to exit (=terminate) the program
Different from return, which simply exits the function
Exit() is defined inside the library stdlib.h
\#include <stdlib.h>
int length( char s[] )\{

```
int size = 0;
while(s[size] != '\0'){
    if(s[size] == 'm')
        exit(-1);
    size++;
}
return size;
```


## COMsW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

## Announcements

Homework 1 correction out this afternoon

Homework 2 is out

- Due Monday, February 28th
- Start early (especially Exercise 2)!


## Today

- Functions
- Recursion
- Debugging (if time)


## Infinite Loops

- Loops where the condition is always TRUE
- Will stop only with:
- break
- modification of the condition variables

```
while ( 1 ){
    /* body modifies x */
    if( x!= 0 ) {
        break;
    }
}
```


## Infinite Loops

- Loops where the condition is always TRUE
- Will stop only with:
- break
- modification of the condition variables



## Operators - Logical

- A variable with value 0 is false, a variable with value !=0 is true

$$
\begin{aligned}
& \text { int } x=3, y=0, z, k, t, q=-3 ; \\
& \mathrm{z}=\mathrm{X} \& \& \mathrm{y} ; \quad / / \mathrm{z}=0 ; \quad \mathrm{x} \text { is true but } \mathrm{y} \text { is false } \\
& \mathrm{k}=\mathrm{x}| | \mathrm{y} ; \quad / / \mathrm{k}=1 ; \quad x \text { is true } \\
& \mathrm{t}=\text { ! } \mathrm{q} ; \quad / / \mathrm{t}=0 ; \quad \text { a is true }
\end{aligned}
$$

## Infinite Loops

- Loops where the condition is always TRUE
- Will stop only with:
- break
- modification of the condition variables

```
int cond = 7;
```

while ( cond ) $\longrightarrow \longrightarrow$ while ( cond is true )
$/ * \operatorname{body} * /$
if( $x[3][5] \quad!=7$ ) $\{$
cond $=0$;
\}
\}


## Functions Example

- Simple calculator
- Program that computes one basic arithmetic operation between 2 numbers


## Functions - Recursion

- What if a function calls itself? Recursion




## Functions - Recursion

- A recursive function must have two properties:
- Ending point (i.e. a terminating condition)
- Simplify the problem (every call is to a simpler input)


## Example: Fibonacci sequence

In mathematics, famous numbers following the sequence 01123581321345589 ...

Given $F_{0}=0, F_{1}=1$ can be computed with recurrence $F_{n}=F_{n-1}+F_{n-2}$
Code to compute the first 100 Fibonacci numbers:

```
int i = 0;
int fib[100];
fib[0] = 0;
fib[1] = 1;
for( i = 2; i < 100 ; i++ ) {
    fib[i] = fib[i-1] + fib[i-2];
\[
\begin{aligned}
& \text { for( } i=2 ; i<100 ; i++ \text { ) \{ } \\
& \quad \text { fib[i] }=\text { fib[i-1] + fib[i-2]; }
\end{aligned}
\]
```

$\}$


## Functions - Recursion

- What if a function calls itself? Recursion
- What is the value of the number at position num in the Fibonacci sequence?

```
/* Fibonacci value of a given position in the sequence */
int fib ( int num ) {
switch(num) {
    case 0:
        return(0);
    case 1:
        return(1);
    default: /* Including recursive calls */
        return(fib(num - 1) + fib(num - 2));
```

    \}
    \}

## Functions - Recursion

- What if a function calls itself? Recursion
- What is the value of the number at position num in the Fibonacci sequence?

```
/* Fibonacci value of a given position in the sequence */
int fib ( int num ) {
switch(num) {
    case 0:
        return(0);
```



``` Ending Points
    case 1:
        return(1);
    default: /* Including recursive calls */
        return(fib(num - 1) + fib(num - 2));
```



```
Simplify problem

\section*{Debugging}

\section*{Debugging}
- Debugging consists basically in finding and correcting run-time errors in your program
- Multiple ways of doing it
- Manual runs (for small programs)
- Insert printf() in key lines
- There also exist INTERACTIVE debugging tools
- We will now see a basic one for UNIX: gdb

\section*{gdb}
1. In order to use gdb on a program, we must use the -g option when compiling it
```

gcc -g program.c -Wall -o nameOfExecutable

```
2. Then, we can use the gdb command to start the interactive debugging environment
gdb nameOfExecutable
1.
```

~0
\$ gcc -g test.c -o test
\$ gdb test
GNU gdb 5.3
Copyright 2002 Free Software Foundation, Inc.
GDB is free software, covered by the GNU General Public Licens
e, and you are
welcome to change it and/or distribute copies of it under cert
ain conditions.
Type "show copying" to see the conditions.
There is absolutely no warranty for GDB. Type "show warranty"
for details.
This GDB was configured as "sparc-sun-solaris2.9"...
(gdb)
(gdb)

```

\section*{gdb commands}
- run : run executable (program)currently watched.
(gdb) run
- kill : kill current execution of program
(gdb) kill
- list : show program source code (gdb) list 2,8 : shows lines 2 to 8 from source program
- print : print value of a variable or expression at the current point (gdb) print buf

\section*{gdb commands}
- break : insert breakpoint in program. Debugging run will stop at the breakpoint
```

(gd.b) break nameSource.c : lineNumber
(gdb) break test.c: 12

```
- next : step to the next line (execute current line)
(gdb) next
- continue : continue with execution until next breakpoint or end of program (gdb) continue
- Quit : exit gdb
(gdb) quit

\section*{Graphical GDB}
- gdb can be run from Emacs
- Press M-x (in Windows Esc-x)
- Insert gdib
- Insert executableName
- Visual debugger


Can enable breakpoints with a click

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

\section*{Are Computers Smarter than Humans?}

IBM's Watson on 'Jeopardy': Computer takes big lead over humans in Round 2

February \(15,2011 \mid 9: 20 \mathrm{pm}\)
\[
\text { (o) } f \text { (o) Comments (o) }
\]


Link

On Tuesday night's "Jeopardy" episode, Watson, the IBM supercomputer, steamrollered to a commanding lead over his human competitors.
http://latimesblogs.latimes.com/technology/2011/02/ibms-watson-on-jeopardy-computer-takes-big-lead-over-humans-in-round-2.html

\section*{Today}
- Homework 1 Correction
- Debugging (from Lecture 8 )
- C Preprocessor

\section*{Conditional Assignment}
- Another way of embedding if-else in a single statement
- Uses the ? : operators
\[
\begin{array}{r}
\text { variable }=\left(\begin{array}{c}
\text { condition }) ~ ? ~ v a l 1 ~: ~ v a l 2 ~ ; ~ \\
\downarrow
\end{array} \begin{array}{l}
\text { If condition is true, we } \\
\text { assign val1 to variable }
\end{array} \quad \begin{array}{l}
\text { If condition is false, we } \\
\text { assign val2 to variable }
\end{array}\right.
\end{array}
\]
\[
\begin{aligned}
& \text { int } x=7, y \text {; } \\
& y=(x>5) ? x: 5 ; \\
& y=7 \\
& \text { int } x=7, y \text {; } \\
& \text { if( } x>5 \text { ) \{ } \\
& y=x ; \\
& \text { else\{ } \\
& y=5 \text {; }
\end{aligned}
\]

\section*{The comma operator}
- In C statements can also be separated by , not only ;
```

int x = 2;
int y;
x++, y = x/3, y += 2;
x++;
y = x/3;
y = y+2;

```
Be careful with declarations!
\[
\begin{aligned}
& \text { int } x=2, \text { char } c=\text { 'm'; Different types, NO } \\
& \text { int } x=2, y ; \quad V \text { Same type, OK }
\end{aligned}
\]

\section*{The comma operator}

\section*{Special case, the for loop statement}

Example: the palindrome word checking. Check if a word is the same when read right to left
int i, flag = 1;
char word[100] = "radar";
for ( \(i=0, j=\operatorname{strlen}(w o r d)-1 ; i<\operatorname{strlen(word)/2;i++,j--1)}\{\)
if( word[i] != word[j] ) \{
flag \(=0\);
break;
\}
\}

\section*{The comma operator}

\section*{Special case, the for loop statement}

Example: the palindrome word checking


\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

const double PI = 3.14;
double r = 5, circ;
circ = 2 * PI * r;
PI = 7;

```

\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

const double PI = 3.14;
double r = 5, circ;
circ = 2 * PI * r;
Once it's initialized, a const variable cannot change value

```

\section*{C Preprocessor}

\section*{C Preprocessor}

Preprocessor is a facility to handle
- Header files
- Macros

Independent from C itself, it's basically a text editor that modifies your code before compiling

Preprocessor statements begin with \# and do not end with;

\section*{C Preprocessor}



\section*{C Preprocessor}
myFile.c (program)

myFile.c
(preprocessor code)

myFile (executable)
\begin{tabular}{c}
0100101010021 \\
0101001010000 \\
\(11110011 \ldots\) \\
\(\vdots\) \\
\(\ldots 010010100001\) \\
1110001110101 \\
\hline
\end{tabular}

\section*{View Preprocessor Code}

\section*{gcc has a special option that allows to run only the preprocessor gcc -E myFile.c}

We can send output to a file using the UNIX > operator
gcc -E myFile.c > outFile.txt
Saves gcc's output to outFile.txt

\section*{Header files}
- Header files are fundamentally libraries
- Their extension is .h
- They contain function definitions, variables declarations, macros
- In order to use them, the preprocessor uses the following code
\#include <nameOfHeader.h>
\#include "nameOfHeader.h" \(\longrightarrow\) For standard C libraries
- So far, we have used predefined \(C\) header files, but we can create our own! (more on this in upcoming Lectures)

\section*{Header files}


\section*{Macros}
- A macro is a piece of code \(\boldsymbol{c}\) which has been given a name \(\boldsymbol{n}\)
- Every time we use that \(\boldsymbol{n}\) in our program, it gets replaced with c
- The preprocessor allows you to declare them with \#define
- Two types:
- Object-like macros
- Function-like macros

\section*{Object like macros}
- Constants, usually defined on top of programs
```

\#define name text_to_substitute

```
\#define SIZE 10
\#define FOR_ALL for ( i=0; i< SIZE; i++ )

\section*{Object like macros}
```

From now on, every time we write SIZE inside our program it is going to be replaced by 10

```

\section*{Object like macros}
- Some compilers do not allow you to declare arrays with a variable as size
```

int size1 = 10;
int arr1[size1]; /* should always cause error */
const int size2=10;
int arr2[size2]; /* causes errors in many compilers */
\#define SIZE 10
int arr3[SIZE]; /* OK in any C compiler */

```

\section*{Function-like macros}
- Macros that can take parameters like functions
```

\#define SQR(x) ((x) * (x))
\#define MAX(x,y) ((x) > (y) ? (x) : (y))

```
- Parameters MUST be included in parentheses in the macro name, without spaces
- It is a good habit to include parameters in parentheses also in the text to be substituted

\section*{Conditional Compilation}
- Allows to use or not certain parts of a program based on definitions of macros
```

\#ifdef var if var is defined, consider the following code
\#ifndef var if var is not defined, consider the following code
\#else
\#endif close if(n)def
\#undef var undefinevar (opposite of \#define)

```

\section*{Conditional Compilation}
```

\#define DEBUG
•
•
\#ifdef DEBUG
printf("The value of x is %d\n", x);
\#endif

```

If DEBUG was defined earlier in the program, then the statement printf(...); is considered, otherwise the preprocessor does not copy it to the file to be compiled

\section*{COMSW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

\author{
Lecture 10
}

Spring 2011
Instructor: Michele Merler

\section*{Announcements}

Change in Office Hours this week

1 hour Wednesday, Feb 23 \({ }^{\text {rd }}, 12 \mathrm{pm}-1 \mathrm{pm}\)
1 hour Saturday, Feb 26 th, \(11 a m-12 p m\)

\section*{Today}
- Preprocessor (from Lecture 9)
- Advanced C Types

\section*{Advanced Types - Struct}
- Arrays group variables of the same type
- Structs group variables of different types

Struct definition
```

struct structName {
fieldType fieldNameval1;
fieldType fieldNameval2;
fieldType fieldNamevalN;
};

```

Once we define the struct, we can use structName as if were a type, to create variables!

\section*{Advanced Types - Struct}

Example: we want to build a database with the name, age and grade of the students in the class
\begin{tabular}{|l|}
\hline Student 1 \\
Name: \\
Age: \\
Grade: \\
\hline
\end{tabular}
\begin{tabular}{|l|}
\hline Student 2 \\
Name: \\
Age: \\
Grade: \\
\hline
\end{tabular}
\begin{tabular}{|l|}
\hline Student N \\
Name: \\
Age: \\
Grade: \\
\hline
\end{tabular}
```

    struct student {
    char name[100];
    int age;
    double grade;
    };
    struct student st1;

```
    st1 is a variable of
    type struct!

\section*{Advanced Types - Struct}

In order to access struct fields, we need to use the
. operator
```

struct student {
char name[100];
int age;
double grade;
};
struct student st1, st2;
st1.age = 3;
st2.age = st1.age - 10;

```

\section*{Advanced Types - Struct}

We can initialize a struct variable at declaration time, just like with arrays
```

struct student {
char name[100];
int age;
double grade;
};
struct student st1 = {"mike", 22, 77.4};

```

\section*{Advanced types - Typedef}
typedef is used to define a new type
typedef type nameOfNewType;
typedef int myInt;
myInt \(\mathrm{c}=3\);
C is of type my Int, which is equivalent to int
typedef int myIntArray[7];
myIntArray arr; \(\longrightarrow\) arr is of type myIntArray, which is equivalent to an array of
for (c=0; c<7; c++) \{ 7 int \(\operatorname{arr}[c]=1 ;\)
\}

\section*{Advanced types - Typedef}
typedef is used to define a new type
```

struct student {
char name[100];
int age;
double grade;
};
struct student st1, st2;
st1.age = 3;
st2.age = st1.age - 10;

```

\section*{Advanced Types - Union}
- Similar to struct, but all fields share same memory
- Same location can be given many different field names
struct value\{
int iVal;
float fVal;
iVal
fVal
```

union value{
int iVal;
float fVal;
};
iVal / fVal

```

We can use the fields of the union only one at a time!

\section*{Advanced Types - Enum}
- Designed for variables containing only a limited set of values
- Defines a set of named integer constants, starting from 0
\[
\text { enum name\{ item1, item2, ... , itemN\}; }
\]
```

| 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |

enum dwarf { BASHFUL, DOC, DOPEY, GRUMPY, HAPPY, SLEEPY, SNEEZY};
enum dwarf myDwarf = SLEEPY;
myDwarf = 1 + HAPPY; // myDwarf = SLEEPY = 5;
int x = GRUMPY + 1; // x = 4;
printf("dwarf %d\n",BASHFUL); // 'dwarf 0'

```

\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

const double PI = 3.14;
double r = 5, circ;
circ = 2 * PI * r;
PI = 7;

```

\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

const double PI = 3.14;
double r = 5, circ;
circ = 2 * PI * r;
Once it's initialized, a const variable cannot change value

```

\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

double computeCirc( const double r, const double PI){
r++; PI++;
return(2 * r * PI);
}
/* main function */
int main(){
const double PI = 3.14;
double r = 5, circ, circ2;
circ = 2 * PI * r;
circ2 = computeCirc(r, PI);
return 0;
}

```

\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

double computeCirc( double r, const double PI){
r++;
PI++;
return(2 * r * PI);
}
/* main function */
int main(){
const double PI = 3.14;
double r = 5, circ, circ2;
circ = 2 * PI * r;
circ2 = computeCirc(r, PI);
return 0;

```

\section*{Advanced Types - Const}
const defines a variable whose value cannot be changed
```

double computeCirc( double r, double PI){
r+;; V
PT+; \
return(2 * r * PI);
}
/* main function */
int main(){
const double PI = 3.14;
double r = 5, circ, circ2;
circ = 2 * PI * r;
circ2 = computeCirc(r, PI);
return 0;

```

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

Lecture 11
Spring 2011
Instructor: Michele Merler

\section*{Announcements}
- Grades for Homework 1 posted on Coursewors
- Homework 2 is due next Monday at the beginning of class
- Bring the printout to class!

\section*{Pointers}

\section*{Pointers}

Remember what happens when we declare a variable: the computer allocates memory for it.
 computer

\section*{Pointers}

When we assign a value to a variable, the computer stores that value at the address in memory that was previously allocated for that variable.
```

int x
x = 3;

```

\[
\mathrm{x} *=3 ; / / \mathrm{x}=9
\]

Main memory
\(\longrightarrow 000000000 \quad 00000000 \quad 00000000 \quad 00001001\)

\section*{Pointers}

Pointers are variables for memory addresses.
They are declared using the * operator.
They are called pointers because they point to the place in memory where other variables are stored.

How can we know what the address in memory of a variable is? The \& operator.


\section*{Pointers - Syntax}

When we declare a pointer, we must specify the type of variable it will be pointing to
type *ptrName;

If we want to set a pointer to point to a variable, we must use the \& operator
```

ptrName = \&varName;

```
```

int x;
x = 3;
int *y;
y = \&x;

```


\section*{Pointers : operators * and \&}
* dereference operator : gives the value in the memory pointed by a pointer (returns a value)
\& reference operator: gives the address in memory of a variable (returns a pointer)
\[
\begin{aligned}
& \text { int } x=3 ; \\
& \text { int } * p t r ;
\end{aligned}
\]

Make ptr point to the address of \(x\)
```

ptr = \&x; ptr }\longrightarrow000000000 00000000 00000000 00000011

```

Modify the
*ptr \(=5 ; / / \mathrm{x}=5\);
value in address pointed by ptr

\section*{Pointers : operators * and \&}
* dereference operator : gives the value in the memory pointed by a pointer (returns a value)
\& reference operator: gives the address in memory of a variable (returns a pointer)
\[
\begin{aligned}
& \text { int } x=3 ; \\
& \text { int } * \text { ptr }
\end{aligned}
\]

Make ptr point to the address of \(x\)
\[
\text { ptr }=\& x
\]

Modify the value in address pointed by ptr
\begin{tabular}{|l|l|}
\hline Code & Meaning \\
\hline\(x\) & Variable of type int \\
\hline ptr & \begin{tabular}{l} 
Pointer to an element of \\
type int
\end{tabular} \\
\hline \&x & Pointer to \(x\) \\
\hline *ptr & Variable of type int \\
\hline & \\
\hline
\end{tabular}

\section*{Pointers : operators * and \&}
* dereference operator : gives the value in the memory pointed by a pointer (returns a value)
\& reference operator: gives the address in memory of a variable (returns a pointer)
```

int x;
int *ptr;

```

```

\&ptr // pointer to a pointer
*x // x is not a pointer

```

\section*{Pointers : operators * and \&}
* dereference operator : gives the value in the memory pointed by a pointer (returns a value)
\& reference operator: gives the address in memory of a variable (returns a pointer)


\section*{Pointers}

Multiple pointers can point to the same address
```

int x = 3, y = 2;
int *ptr = \&x;
int *ptr2 = ptr;

```


NOTE: first 4 bits omitted to save space

\section*{Pointers}

Multiple pointers can point to the same address
\[
\begin{aligned}
& \text { int } x=3, y=2 ; \\
& \text { int } * p t r=\& x ; \\
& \text { int } * p t r 2=\text { ptr; } \\
& \text { *ptr }=7 ; \quad / / x=7
\end{aligned}
\]


NOTE: first 4 bits omitted to save space

\section*{Pointers}

Multiple pointers can point to the same address
\[
\begin{aligned}
& \text { int } x=3, y=2 ; \\
& \text { int } * \operatorname{ptr}=8 x ; \\
& \text { int } * \text { ptr2 }=\text { ptr; } \\
& \text { *ptr }=7 ; \quad / / x=7 ; \\
& \text { *ptr2 }=* \operatorname{ptr} 2+1 ; / / x=8 ;
\end{aligned}
\]


NOTE: first 4 bits omitted to save space

\section*{Pointers}

Multiple pointers can point to the same address
```

int x = 3, y;
int *ptr = \&x;
int *ptr2 = ptr;
*ptr = 7; // x = 7;

```

```

ptr = \&y;
*ptr2 = 10; // x = 10;
Ptr2 is still pointing to $x$, even if ptr changed

```

\section*{Pointers}

\section*{Be careful when using incremental operators!}
```

int x = 3;
int *ptr = \&x;
*ptr++; // x = ?

```


In this case I am incrementing ptr, NOT the value of the variable pointed by it!

\section*{Pointers}

\section*{Be careful when using incremental operators!}
```

int x = 3;
int *ptr = \&x;
(*ptr)++; // x = 4;

```
ptr \(\longrightarrow\)\begin{tabular}{|c|c|c|c|}
\hline \multicolumn{5}{|c|}{\begin{tabular}{c} 
Main \\
memory
\end{tabular}} \\
\hline 0000 & 0000 & 0000 & 0100 \\
\hline & & & \\
\hline
\end{tabular}

\section*{Pointers and Arrays}
- When set a pointer to an array, the pointer points to the first element in the array
```

float arr[3] = {1, 2, 5};
float *pa;
pa = arr;
pa =\&arr[0]; }

```
- C automatically keeps pointer arithmetic in terms of the size of the variable type being pointed to
```

arr[0] \longleftrightarrow *pa
arr[1] \longleftrightarrow*(pa+1)
arr[2] \longleftrightarrow pa[2]

```

\section*{Pointers and Arrays}
- When set a pointer to an array, the pointer points to the first element in the array
```

float arr[3] = {1, 2, 5};
float *pa;
pa = arr;
pa =\&arr[0]; }

```
- C automatically keeps pointer arithmetic in terms of the size of the variable type being pointed to
 \(\longrightarrow\) Once we have set a pointer to the beginning of one array, we can use it as if it were the array itself!

\section*{Pointers and Arrays}

When set a pointer to an array, the pointer points to the first element in the array
\[
\begin{aligned}
& \text { float } \operatorname{arr}[3]=\{1,2,5\} \text {; } \\
& \text { float } * p=\operatorname{arr} \\
& * p=5 ; / / \operatorname{arr}[0]=5 ;
\end{aligned}
\]


\section*{Pointers and Arrays}

When set a pointer to an array, the pointer points to the first element in the array
\[
\begin{aligned}
& \text { float } \operatorname{arr}[3]=\{1,2,5\} ; \\
& \text { float } * p=\operatorname{arr} ; \\
& * p=5 ; / / \operatorname{arr}[0]=5 ; \\
& p++; \\
& * p=3 ; / / \operatorname{arr}[1]=3 ;
\end{aligned}
\]


\section*{Pointers and Arrays}

When set a pointer to an array, the pointer points to the first element in the array


\section*{Pointers and Arrays}
word[0] word[1] word[2] word[3] word[4] word[5]

```

char *wPtrStart = word;
char *wPtrEnd = wPtrStart + strlen(word)-1;
for( i=0 ; (i < strlen(word)/2) \&\& (flag == 1) ; i++ ){
if( *wPtrStart != *wPtrEnd ){
flag = 0;
}
wPtrStart++;
wPtrEnd--;

```

\title{
Pointers and Arrays
}
word[0] word[1] word[2] word[3] word[4] word[5]

```

char *wPtrStart = word;
char *wPtrEnd = wPtrStart + strlen(word)-1;
for( i=0 ; (i < strlen(word)/2) \&\& (flag == 1) ; i++ ){
if( *wPtrStart != *wPtrEnd ){
flag = 0;
}
wPtrStart++;
wPtrEnd--;
When we increment or decrement, the pointers move by 1 byte (pointers to char)

## Pointers : operators * and \&

Now we know exactly what happens in sscanf!


## Functions

## Passing arguments by value/reference

- Pass by value (what we have seen so far): the value of the variable used at invocation time is copied into a local variable inside the function
- Pass by reference : a pointer to the variable used at invocation time is passed to the function. We can modify the variable's value inside the function


## Functions

## Passing arguments by value/reference

- Pass by value (what we have seen so far): the value of the variable used at invocation time is copied into a local variable inside the function

```
double computeCirc( double rad ){
    rad = 2;
    return(2 * rad * 3.14);
}
int main(){
    double r = 5, circ;
    circ = computeCirc(r);
    return 0;
```


## Functions

## Passing arguments by value/reference

- Pass by value (what we have seen so far): the value of the variable used at invocation time is copied into a local variable inside the function

```
double computeCirc( double rad ){
    rad = 2;
    return(2 * rad * 3.14);
}
int main(){
    double r = 5, circ;
    circ = computeCirc(r);
    return 0;
```

$r$ is not affected by anything we do inside the function

## Functions

## Passing arguments by value/reference

- Pass by reference : a pointer to the variable used at invocation time is passed to the function. We can modify the variable's value inside the function
int main(){
double r = 5, circ;
circ = computeCirc(\&r)
return 0;

```

\section*{Functions}

\section*{Passing arguments by value/reference}
- Pass by reference: a pointer to the variable used at invocation time is passed to the function. We can modify the variable's value inside the function
```

double computeCirc( double *rad ){
*rad = 2;
return(2 * (*rad) * 3.14);
}
int main(){
double r = 5, circ;
circ = computeCirc(\&r);
return 0;

```
                            r has been modified!

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

Lecture 12
Spring 2011
Instructor: Michele Merler

\section*{Announcements}

Homework 3 is out
- Due on Monday, 03/21/11 at the beginning of class, no exceptions

\section*{Midterm}
- In class on Wednesday, 03/09/11
- Will cover everything up to Lecture 13 (included)
- Open books, open notes
- Closed electronic devices

\section*{Today}
- Passing arguments to function by value vs. by reference (from Lec 11)
- Functions returning pointers
- Pointers of pointers

\section*{Functions Returning Pointers}
- Naturally, a function can return a pointer
- This is a way to return an array, but must be careful about what has been allocated in memory
```

returnType * functionName( parmeters )

```

NOTE
NULL is the equivalent of zero for pointers

\section*{Functions Returning Pointers}

Example: using pointers to return a string

Given a string of the type "firstNAme/lastName" We want to split it into two separate entities to print

\section*{Functions Returning Pointers}

POINT 1


POINT 2


POINT 3


POINT 4


\section*{Const pointers}
```

const type *

```

When we try to declare a pointer to be a constant like this, it means that the value at the address in memory it points cannot be modified

This does NOT mean that the pointer is constant, it can be changed!
```

int x = 7, y = 3;
const int *ptr = \&x;
*ptr = 11;

```


\section*{Const pointers}
```

const type *

```

When we try to declare a pointer to be a constant like this, it means that the value at the address in memory it points cannot be modified

This does NOT mean that the pointer is constant, it can be changed!
```

int }x=7,y=3
const int *ptr = \&x;
*ptr = 11;
x = 8; V

```
\begin{tabular}{|c|c|c|c|}
\hline \multicolumn{5}{|c|}{\begin{tabular}{c} 
Main \\
memory
\end{tabular}} \\
\hline 0000 & 0000 & 0000 & 1000 \\
\hline & \(x\) \\
\hline 0000 & 0000 & 0000 & 0011 \\
\hline & \(y\) \\
\hline
\end{tabular}

\section*{Const pointers}
```

const type *

```

When we try to declare a pointer to be a constant like this, it means that the value at the address in memory it points cannot be modified

This does NOT mean that the pointer is constant, it can be changed!
\[
\begin{aligned}
& \text { int } x=7, y=3 ; \\
& \text { const int } * p t r=\& x \\
& * p t r=11 ; \\
& x=8 \\
& \text { ptr }=\& y
\end{aligned}
\]


\section*{Const pointers}
```

const type *

```

When we try to declare a pointer to be a constant like this, it means that the value at the address in memory it points cannot be modified

This does NOT mean that the pointer is constant, it can be changed!
```

int }x=7,y=3
const int *ptr = \&x;
*ptr = 11;
x = 8;
ptr = \&y; V
*ptr = 9;
printf("x = %d, y = %d\n",x,*ptr);

```

\section*{Const pointers}
```

type * const

```

This is the declaration of a constant pointer. In this case, the pointer is fixed, but the value at the address it points to can be modified
```

int x = 7, y = 3;
int * const ptr2 = \&x;
*ptr2 = 9; V
ptr2++;
ptr2 = \&y;
printf("x = %d, x = %d\n", x, *ptr2);

```

\section*{Arrays of strings}
- An array Arr of 3 strings of variable length
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
Arr[2] \longleftrightarrow Arr+2 // "Wondeful"

```
- Arr is an array of \(\mathbf{3}\) elements. Each element in Arr is of type pointer to char.


\section*{Arrays of strings}
- An array Arr of 3 strings of variable length
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
Arr[2] \longleftrightarrow Arr+2 // "Wondeful"

```
- An array Arr of 3 strings of maximum length \(=15\)
```

char Arr2[3][15] = { "Hello2", "World2", "Wonderful2" };
Arr2[0] \longleftrightarrow Arr2 // "Hello2"
Arr2[1] \longleftrightarrow Arr2+1 // "World2"

```
```

stringArrays.c

```

\section*{Pointers of pointers}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline \multirow[b]{2}{*}{Arr} & \multirow[b]{2}{*}{0} & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & \multirow[t]{2}{*}{9} \\
\hline & & 'H' & 'e' & T & ' & '0' & 10 & & & & \\
\hline & 1 & 'W' & 'o' & 'r' & \(\uparrow\) & 'd' & '10' & & & & \\
\hline & 2 & 'W' & 'o' & ' n ' & 'd' & 'e' & 'r' & 'f' & 'u' & \(\uparrow\) & \(10^{\prime}\) \\
\hline
\end{tabular}


\section*{Pointers of pointers}
- A pointer can point to another pointer
- In a sense, it's the equivalent of matrices!
```

int x = 3;
int *p = \&x;
int **p2 = \&p;
x = 2; \longleftrightarrow *
char *Arr[3]={ "Hello", "World", "Wonderful" };
char **ptr;
ptr = Arr;

```
```

stringArrays.c

```

\section*{Pointers of pointers}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|}
\hline 0 & 1 & 2 & 3 & 4 & 5 & \multirow[t]{3}{*}{6} & \multirow[t]{3}{*}{7} & \multirow[t]{3}{*}{8} & \multirow[t]{3}{*}{9} \\
\hline Arr \(-{ }^{\text {a }}{ }^{\prime}\) & 'e' & ' & ' & 'o' & 10 & & & & \\
\hline \(\mathrm{ptr}+1 \xrightarrow{\longrightarrow} \mathrm{~W}^{\prime}\) & 'o' & ' r ' & T & 'd' & '10' & & & & \\
\hline \(2{ }^{\prime}{ }^{\prime}\) ' & 'o' & ' n ' & \({ }^{\prime} \mathrm{d}\) ' & 'e' & 'r' & (f) & 'u' & \(\uparrow\) & '10' \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|}
\hline \(2{ }^{\prime}\) ' \({ }^{\prime}\) & 'o' & ' n ' & 'd' & 'e' & 'r' & 'f' & 'u' & I & '2' & 10 \\
\hline 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\
\hline
\end{tabular}
\(11 \quad 12 \quad 13 \quad 14\)

\section*{Pointers of pointers}
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
char **ptr;
ptr = Arr;

```
*(*(ptr+1)+2)
?

\section*{Pointers of pointers}
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
char **ptr;
ptr = Arr;

```
*(*(ptr+1)+2)
1. \(\mathrm{ptr}+1\)

ptr+1 points to the whole line


2. *(ptr+1)
points to the first element of the line


2
 f
\begin{tabular}{l|l|l} 
'u' & I' & 10 '
\end{tabular}

\section*{Pointers of pointers}
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
char **ptr;
ptr = Arr;

```
*(*(ptr+1)+2)
3. \({ }^{*}(\mathrm{ptr}+1)+2\)
*(ptr+1)+2 points to the
 third element of the line

\[
\text { 2. }{ }^{*}(*(p t r+1)+2)
\]

Now we get the value stored at the address we
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|}
\hline & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\
\hline 0 & 'H' & 'e' & T & T & 'o' & '10' & & & & \\
\hline
\end{tabular} pint


\section*{Pointers of pointers}
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
char **ptr;
ptr = Arr;

```
*(*(ptr+1)+2)
Avoid this notation! ptr[1][2] is much better!
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|}
\hline & 0 & 1 & 2 & 3 & 4 & 5 & \multirow[t]{2}{*}{6} & \multirow[t]{2}{*}{7} & \multirow[t]{2}{*}{8} & \multirow[t]{2}{*}{9} \\
\hline 0 & 'H' & 'e' & \(\checkmark\) & \(\checkmark\) & 'o' & '10' & & & & \\
\hline 1 & 'W' & 'o' & 'r' & \(\checkmark\) & 'd' & ' 10 & & & & \\
\hline 2 & 'W' & 'o' & ' n ' & 'd' & 'e' & \({ }^{\prime} \mathrm{r}\) ' & 'f' & 'u' & T & '10' \\
\hline
\end{tabular}

\section*{Pointers vs. Arrays}

\section*{Arrays}

\section*{Pointers}
\begin{tabular}{l|l}
1 D array of 5 int & int \(\mathrm{x}[5] ;\) \\
2D array of 6 int & int \(y[2][3] ;\)
\end{tabular}

int* \(\mathrm{z}[2]=\{\{1,2\},\{2,1\}\} ; \longleftrightarrow\) int \(* *\) zPtr;
char c[]\(=\) "mike";
char *CPtr;
1D array of 5 char string

Space has been
allocated in memory for the arrays

Space has been allocated in memory only for the pointers variables, NOT for the arrays they will point to.
The DIMENSIONS of the arrays are UNKNOWN

\section*{Multidimensional Arrays}
\(2 \times 3\) matrix of double
```

double MO[2][3];
double *M1[2] = M0;
double **M = MO;

```
\begin{tabular}{|l|l|l|l|}
\hline\(M\) & \(M M[0][0]\) & \(M[0][1]\) & \(M[0][2]\) \\
& \(\longrightarrow M[1]\) & \(M[1][0]\) & \(M[1][1]\) \\
\hline & \(M[1][2]\) \\
\hline
\end{tabular}
double **
double *
double

\section*{Multidimensional Arrays}
\(2 \times 3\) matrix of double
```

double MO[2][3];
double *M1[2] = MO;

```

The difference between \(\mathrm{M} 0, \mathrm{M} 1\) and M is that M1 and M can have ANY SIZE!
\begin{tabular}{|c|c|c|c|c|}
\hline M & M[0] & M[0][0] & M[0][1] & M[0][2] \\
\hline & M[1] & M[1][0] & M[1][1] & \(\mathrm{M}[1][2]\) \\
\hline
\end{tabular}
double **
double *
double

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

Lecture 13
Spring 2011
Instructor: Michele Merler

\section*{Today}
- Finish pointers (from Lecture 12 )
- FILE I/O

\section*{Pointers of pointers}
float \(\mathrm{A}[2]=\{1,2\}\);
float \(B[3]=\{7,1,5\}\);


\section*{Pointers of pointers}
float \(\mathrm{A}[2]=\{1,2\}\);
float \(B[3]=\{7,1,5\}\);
float * \(\mathrm{p}=\mathrm{B}\);


\section*{Pointers of pointers}
float \(\mathrm{A}[2]=\{1,2\}\);
float \(B[3]=\{7,1,5\}\);
float * \(\mathrm{p}=\mathrm{B}\);
float *p1[2];


\section*{Pointers of pointers}
float \(\mathrm{A}[2]=\{1,2\}\);
float \(B[3]=\{7,1,5\}\);
float * \(=B\);
float *p1[2];
\(\mathrm{p} 1[0]=\mathrm{A} ; / / \mathrm{p} 1[0]\) is a pointer to float \(\mathrm{p} 1[1]=\mathrm{B}\); // \(\mathrm{p} 1[1]\) is a pointer to float


\section*{Pointers of pointers}
float \(\mathrm{A}[2]=\{1,2\}\);
float \(B[3]=\{7,1,5\}\);
float * \(\mathrm{p}=\mathrm{B}\);
float *p1[2];
p1[0] = A;
p1[1] \(=B ;\)
float **p2 = p1;


\section*{Pointers of pointers}
float \(\mathrm{A}[2]=\{1,2\}\);
float \(B[3]=\{7,1,5\}\);
float * \(\mathrm{p}=\mathrm{B}\);
float *p1[2];
p1[0] = A;
p1[1] \(=B\);
float \({ }^{* *}\) p2 \(=\) p1;
float \(\mathrm{f} 1=\mathrm{p} 2[0][2] ; / / \mathrm{f} 1=\mathrm{A}[2]=\) float \(\mathrm{f} 2=\mathrm{p} 2[1][2] ; / / \mathrm{f} 2=\mathrm{B}[2]=5\)
float \(\mathrm{f} 3=\mathrm{p} 2[0][1] ; / / \mathrm{f} 3=\mathrm{A}[1]=2\)


\section*{Files Input/Output}

\section*{Files I/O}
- So far we have seen functions to read/write to command line (standard input/output)
- The same functions can be used to read/write to files
- (f)printf(), (f)scanf(), fgets()
- All those functions are included in the <stdio. h> library

\section*{Files I/O Pipeline}
- Files have a special type of variable associated with them: FILE *
- In order to read/write to a file, we must first OPEN it
- After we are done, we must CLOSE the file


\section*{Files I/O}
- Files have a special type of variable associated with them: FILE *
- In order to read/write to a file, we must first OPEN it
- After we are done, we must CLOSE the file


\section*{fopen()}

\section*{FILE * fopen( char *fileName, char *mode);}
- fileName is a regular string with the name of the file
- mode determines the type of I/O we want to do
- "r" : read
- "w" : write, fileName is created if it did not exist
- "a" : append, write to existing file, starting at the end
- "b" : file is binary (associated with other modes, for example "wb" means write binary, "rb" read binary, etc.)
- " \(\mathrm{r}+\) " : read and write
- "w+" : read and write, fileName is created if it did not exist
- In case of failure (for example trying to read from a non-existing file) fopen() returns NULL

\section*{fclose()}
```

int fclose( FILE *fVar );

```
- fVar is a file variable (type FILE *)
- fclose() returns
- 0 on success
- non-zero for error

\section*{Stdin, stdout, stderr}
- C provides 3 files (or filestreams) which are always open:
- stdin : standard input, read from command line
- stdout: standard output, write to command line
- stderr: standard error, write to command line
- They are used as default values for various I/O functions

\section*{Read Functions}
- fgetc() : read a single character
```

int fgetc( FILE *fVar )

```

Returns the special flag EOF if it has reached the end of the file
- fgets() : read a string, one line at a time
char* fgets( char* string, size_t size, FILE *fVar )
Returns string if successful, NULL is error or found EOF

\section*{Read Functions}
- fscanf() : read a formatted line
int fscanf( FILE *fVar,"format1 ... formatN", \&var1,...\&varN)

Reads one line from a file

Returns the number of variables successfully converted

\section*{Write Functions}
- fputc() : write a single character
int fputc ( char ch, FILE *fVar )

Returns ch if successful, the special flag EOF if there is an error
- fputs() : write a string
int fputs ( const char *string, FILE *fVar )
Returns a nonzero number if successful, EOF if there is an error

\section*{Write Functions}
- fprintf() : print to file a formatted line
int fprintf( FILE *fVar,"format1 ... formatN", var1,..., varN)

Prints one line to a file
Returns the number of variables successfully converted

\section*{Read/Write to Files}
- C has an internal pointer to the current position in the opened file
- After each read/write operation the pointer is updated

FILE *inFile = fopen("data.txt","r");
this is a file to read \(\backslash n\) can we do it? \(\backslash n\) 2 * \(3 \backslash \mathrm{n}\)
int \(c h=\) fgetc(inFile);


\section*{feof()}
- feof() checks if we reached the end of a file, without having to use fget(), fscanf() etc.
```

int feof( FILE *fVar )

```

Returns a value different from zero if reached end of file , zero otherwise
FILE *inFile = fopen( "data.txt","r");
```

while(1) {
int ch = fgetc(inFile);
if( ch = EOF ){
break;
}
}

```
```

while( !feof(inFile) ) {
int ch = fgetc(inFile);
}

```

\section*{Summary of Functions}
\begin{tabular}{|l|l|l}
\hline Name & Input & Output \\
\hline fprintf() & formatted text + args & file \\
\hline printf() & formatted text + args & stdout \\
\hline sprintf() & formatted text + args & string \\
\hline fputc(), fputs() & char,string & file \\
\hline fscanf() & file & formatted text + args \\
\hline scanf() & stdin & formatted text + args \\
\hline sscanf() & string & formatted text + args \\
\hline fgetc(), fgets() & file & (char) int, string \\
\hline
\end{tabular}

\section*{Buffered Output}
- The OS does not write directly to a file stream
- For efficiency, it first prints to a buffer (= local placeholder in main memory)
- When the buffer is full, it prints it all to the file stream
- If we want to write in a specific moment, without buffering, we can us the function fflush()
```

int fflush( FILE *fVar )

```

Returns 0 if successful, EOF in the case of error

\section*{Buffered Output}
```

printf("starting\n");
do_step1();
printf("done with 1\n");
do_step2();
printf("done with 2\n");
do_step3();
printf("done with 3\n");

```
```

printf("starting\n");
fflush(stdout);
do_step1();
printf("done with 1\n");
fflush(stdout);
do_step2();
printf("done with 2\n");
fflush(stdout);
do_step3();
printf("done with 3\n");
fflush(stdout);

```

Prints to buffer, after last printf() prints to stdout

After each printf() prints to stdout

\section*{File Formatting}
- It is a good habit to create data files with HEADERS, especially when dealing with large amount of data
- HEADERS are one or two lines at the beginning of a file specifying the size of the data and some other info
- With headers, a program knows how to properly read a file
\begin{tabular}{|lllllll|}
\hline \multicolumn{3}{|l}{ VectorTable } \\
cols & 7 & & & & & \\
rows & 3 & & & & & \\
0 & 2 & 5 & 7 & 8 & 22 & 16 \\
10 & 66 & 52 & 7 & 8 & 82 & 6 \\
99 & 1 & 34 & 34 & 87 & 22 & 97
\end{tabular}

\section*{File Formatting}
- It is a good habit to create data files with HEADERS, especially when dealing with large amount of data
- HEADERS are one or two lines at the beginning of a file specifying the size of the data and some other info
- With headers, a program knows how to properly read a file
\begin{tabular}{|lllllll}
\hline \multicolumn{8}{|l|}{\begin{tabular}{l} 
VectorTable \\
cols
\end{tabular}} & 7 \\
rows & 3 & & & & & \\
\hline 0 & 2 & 5 & 7 & 8 & 22 & 16 \\
10 & 66 & 52 & 7 & 8 & 82 & 6 \\
99 & 1 & 34 & 34 & 87 & 22 & 97 \\
\hline
\end{tabular}

\section*{File Formatting}
- Ideally, format should be readable by humans and by computer programs
- Computer programs are not very robust, so must be specific (i.e. tab versus spaces)
- When you have huge amounts of data, you can give up on human-readability and use BINARY format for efficiency
- Example: color_histogram table

\section*{Binary Files}

In order to read/write to binary files, we must use the "rb" / "wb" flags in the option of fopen()
size_t fread(void *ptr, size_t s, size_t n, FILE *f);
size_t fwrite(const void *ptr, size_t s, size_t n, FILE *f);
- ptr = (pointer) array where we want to store the data we read/ we want to write
- \(s=\) size of each element in the array ptr
- \(\mathrm{n}=\) number of elements in the array ptr
- \(f=\) file to read from/write to
size_t is a C type to indicate the size (in bytes) of an element . You can think of it as a special integer.
For example, sizeof() returns a variable of type size_t

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

Lecture 14
Spring 2011
Instructor: Michele Merler

\section*{Announcements}

Homework 4 out on Wednesday, due o n Monday April 11th

Homework 3 solution out later today

\section*{Today}
- Midterm Solution
- Finish FILE I/O (from Lecture 13)
- C standard libraries

\section*{Midterm Solution}

Midterm Solution uploaded to Shared Files in Courseworks

Midterm Statistics
- Average grade: 72
- Standard deviation: 17

\section*{C Standard Libraries}
- C provides a series of useful functions already implemented in standard libraries
- We have already seen some (stdio.h, string.h)
- In order to use the functions in a library, we must include the library header
\#include <libraryName.h>

\section*{C Standard Libraries}

\section*{C Standard Libraries}
- stdio.h : input/output
- string.h : functions on strings
- stdlib.h : utility functions
- math.h : mathematical functions
- ctype.h : character class test
- assert.h : diagnostics
- limits.h and float.h : implementation-defined limits
- time.h : date and time functions
- A few more

\section*{C Standard Libraries}
- stdio.h : input/output
- string.h : functions on strings
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- math.h : mathematical functions
- ctype.h : character class test
- assert.h : diagnostics
- limits.h and float.h : implementation-defined limits
- time.h : date and time functions
- A few more

\section*{stdio.h}
- Standard input and output
- Input/output from command line (keyborad)
- fprintf(), fgets(), sscanf()
- Input/output from files
- FILE, fopen(), fclose()

\section*{string.h}

\section*{Operations involving strings}
```

string s1, s2;

```
char c;
- int \(\mathrm{n}=\operatorname{strcmp}(\mathrm{s} 1, \mathrm{~s} 2):\) compare s 1 and s 2 , \(\mathrm{if}(\mathrm{s} 1==\mathrm{s} 2)->\mathrm{n}=0\)
- int len \(=\) strlen (s1) : return length of s1
- char \(* \mathrm{pc}=\operatorname{strchr}(\mathrm{s} 1, \mathrm{c}):\) return pointer to first occurrence of c in s 1
- char \(* \mathrm{ps}=\operatorname{strstr}(\mathrm{s} 1, \mathrm{~s} 2):\) return pointer to first occurrence of string
\(s 2\) in s1, or NULL if not present
- char \({ }^{\text {strcpy }}(\mathrm{s} 1, \mathrm{~s} 2):\) copy string s2 into s1, return s1
- char *strcat(s1, s2) : append s2 to s1 (concatenate), return s1
- char *strtok(s1, s2) : split long strings into pieces, or tokens

\section*{stdlib.h}

\section*{Number conversions}
- float \(\mathrm{nf}=\) atof(const char \(* \mathrm{~s}\) ) : converts string \(s\) to float
- int \(\mathrm{n}=\) atoi(const char \(*\) s) : convert string s to int

\section*{Memory allocation}
malloc(), free() : memory management

\section*{Other utilities}
- int \(\mathrm{n}=\) rand () : returns a (pseudo) random int between 0 and constant RAND_MAX
- void srand(unsigned int n) : seeds rand generator
- system(string s) : runssin OS

\section*{math.h}
- Mathematical functions
- Often needs to be specially linked when compiling because takes advantage of specialized math hardware in processor
gcc -Im -Wall -o myProgram myProgram.c
```

double functionName( double c )

```
- \(\sin (x), \cos (x), \tan (x)\)
- exp \((x), \log (x), \log 10(x): e^{x}\), natural and base-10 logarithm
- pow \((x, y) \quad: X^{y}\)
- sqrt(x) : square root
- ceil(x), floor(x) : closest int above or below
- \(y=f a b s(x) \quad:\) absolute value , if \(x=-3.2, y\) will be 3.2

\section*{ctype.h}

Utility functions to check for types of char
int functionName( unsigned char c )
- isalpha(c) :check if \(c\) is an alphabet character ' \(a\) '-' \(z^{\prime}\) ', ' \(A^{\prime}\) '- \(Z\) '
- isdigit(c) : check if c is digit ' 0 ' \(\mathbf{\prime}^{\prime} 9\) '
- isalnum(c) :isalpha(c) or isdigit(c)
- iscntrl(c) : control char (i.e. \(\backslash n, \backslash t, \backslash b)\)
- islower(c) , isupper(c) :lowercase/uppercase
\[
\text { Return value is } 0 \text { if false , !=0 if true }
\]

\section*{ctype.h}

Utility functions to convert from lower case to upper case
```

char functionName(char c )

```
- \(d=\) tolower (c) :if \(c\) is ' \(T\) ', \(d\) will be ' \(t\) '
- \(d=\) toupper ( \(c) \quad:\) if \(c\) is ' \(m\) ', \(d\) will be ' \(M\) '

\section*{limits.h and float.h}

Contain various important constants such as the minimum and maximum possible values for certain types, sizes of types, etc.
- CHAR_BIT (bits in a char)
- INT_MAX, CHAR_MAX, LONG_MAX
(maximum value of int, char, long int)
- INT_MIN, CHAR_MIN, LONG_MIN
- FLT_DIG (decimal digits of precision)
- FLT_MIN, FLT_MAX (min. and max. value of float)
- DBL_MIN, DBL_MAX (and of double precision float)

\section*{time.h}

\section*{Provides new type to represent time, time_t}
- time_t time(NULL) : returns current time
- time_t clock() : returns processor time used by program since beginning of execution
- strftime(A, sizeof(A), "formatted text", time struct) :
format text with placeholders:
\%a weekday
\%b month
\%c date and time
\%d day of month
\%H hour

\section*{assert.h}
- Provides a macro to check if critical conditions are met during your program
- Nice way to test programs
```

assert( expression )

```

If the expression is false, the program will print to command line:
Assertion failed: expression , file filename, line lineNumber

\section*{More}
- stdarg.h : allows you to create functions with variable argument lists
- signal.h - provides constants and utilities for standardized error codes for when things go wrong

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

\section*{Announcements}

Homework 4 out, due April \(11^{\text {th }}\) at the beginning of class

Read CPL Chapter 5

\section*{Today}
- Finish C Standard Libraries
- Pointers to void
- Begin Dynamic Memory Allocation

\section*{Review : operators * and \&}
* dereference operator : gives the value in the memory pointed by a pointer (returns a value)
\& reference operator: gives the address in memory of a variable (returns a pointer)
\[
\begin{aligned}
& \text { int } x=3 ; \\
& \text { int } * p t r ;
\end{aligned}
\]

Make ptr point to the address of \(x\)
```

ptr = \&x; ptr }\longrightarrow000000000 00000000 00000000 00000011

```

Modify the
*ptr \(=5 ; / / \mathrm{x}=5\);
value in address
pointed by ptr

\section*{Review : Pointers of pointers}
- A pointer can point to another pointer
- In a sense, it's the equivalent of matrices!
```

int x = 3;
int *p = \&x;
int **p2 = \&p;
x = 2; \longleftrightarrow *
char *Arr[3]={ "Hello", "World", "Wonderful" };
char **ptr;
ptr = Arr;

```

\section*{Review: Pointers vs. Arrays}

\section*{Arrays \\ Pointers}

1D array of 5 int \(\mid\) int \(x[5]\);
2D array of 6 int int \(y[2][3]\); 2x3 matrix

2 D array of 4 int \(2 \times 2\) matrix

1D array of 5 char string
char \(c[]=\) "mike";

char *CPtr;

Space has been
allocated in memory for the arrays

Space has been allocated in memory only for the pointers variables, NOT for the arrays they will point to.
The DIMENSIONS of the arrays are UNKNOWN

\section*{Multidimensional Arrays}
\(2 \times 3\) matrix of double
```

double MO[2][3];
double *M1[2] = M0;
double **M = MO;

```
\begin{tabular}{|l|l|l|l|}
\hline\(M\) & \(M M[0][0]\) & \(M[0][1]\) & \(M[0][2]\) \\
& \(\longrightarrow M[1]\) & \(M[1][0]\) & \(M[1][1]\) \\
\hline & \(M[1][2]\) \\
\hline
\end{tabular}
double **
double *
double

\section*{Multidimensional Arrays}
\(2 \times 3\) matrix of double
```

double MO[2][3];
double *M1[2] = MO;

```

The difference between \(\mathrm{M} 0, \mathrm{M} 1\) and M is that M1 and M can have ANY SIZE!
\begin{tabular}{|c|c|c|c|}
\hline \(\mathrm{M} \longrightarrow \mathrm{M}[0]\) \\
& \(\mathrm{M}[0][0]\) & \(\mathrm{M}[0][1]\) & \(\mathrm{M}[0][2]\) \\
\hline \(\mathrm{M}[1]\) & \(\mathrm{M}[1][0]\) & \(\mathrm{M}[1][1]\) & \(\mathrm{M}[1][2]\) \\
\hline
\end{tabular}
double **
double *
double

\title{
Review : Pointers and Arrays
}
word[0] word[1] word[2] word[3] word[4] word[5]

```

char word[8] = "RADAR";
char *wPtrStart = word;

```

\section*{char* is a string}

\section*{Pointers vs. Arrays}
- Arrays represent actual memory allocated space

int myArr[10];
- Pointers point to a place in memory
int *myPtr;
\[
\text { myPtr } \longrightarrow
\]
myArr

\section*{Pointers vs. Arrays}
- Arrays represent actual memory allocated space
myArr
\(\square \square \square \square \square \square \square\)
int myArr[10];
- Pointers point to a place in memory
\[
\begin{aligned}
& \text { int *myPtr; } \\
& \text { myPtr }=\text { myArr; }
\end{aligned}
\]


\section*{sizeof()}
- So far, we have been using sizeof() to determine the length of a string (including ' \(\backslash 0\) ')
- sizeof() is a more general function, that returns the size, measured in bytes, of a variable or a type
size_t sizeof( var )
- size_t can be used (implicitly casted) as an integer

\section*{Void *}
void * means a pointer of ANY type Sometimes functions can use void * as argument and return type.

This allows the programmer to specify the type of pointer to use at invocation time

This is a form of function overloading (popular in C++)
void *function_name( void *arg1, ... , void *argN )
int i;
double d; int *pi; double *pd;
void *pr;

\section*{Void *}
\[
\begin{array}{ll}
\text { pi }=\& d ; & \text { // Compiler warning } \\
\text { pd }=\& i ; & \text { // Compiler warning }
\end{array}
\]
\[
\text { jv = \&i; } \quad / / \text { OK }
\]
printf("\%d\n", *pr); // Compiler error
printf("\%d\n", *(int *)pr); // OK
\[
\mathrm{pv}=\& \mathrm{~d}
\]
// OK
printf("\%f\n", *pr); // Compiler error
printf("\%f\n", *(double *)pr); // OK
\[
\text { pr = \&i; } \quad / / \text { OK }
\]
\[
d=*(d o u b l e ~ *) p v ; ~ / / ~ R u n t i m e ~ e r r o r ~
\]

\section*{Void *}

\section*{Example}
```

void *pointElement( void *A, int ind, int type ){
if( type == 1 ){
return( A + sizeof(int) * ind );
}
}

```
int main() \{
int \(M[3]=\{1,2,3\}\);
int element \(=1\);
int \(* \mathrm{M} 2=(i n t *)\) pointElement \((\mathrm{M}\), element, 1);

\section*{Void *}

\section*{Example}
```

void *pointElement( void *A, int ind, int type ){
if( type == 1 ){
return( A + sizeof(int) * ind );
}
}

```
int main() \{
    Explicit cast
    int \(M[3]=\{1,2,3\}\);
    int element = 1;
    int \(* M 2=(i n t *)\) pointElement \((M\), element, 1\()\);

\section*{Dynamic Memory Allocation}

\section*{Functions related to DMA are in the library stdlib.h}
```

void *malloc( size_t numBytes )

```

Allocates numbytes bytes in memory (specificaly, in a part of memory called heap)

The elements in the allocated memory are not initialized
Returns a pointer to the allocated memory on success, or NULL on failure
```

void *calloc( size_t numElements, size_t size )

```

Allocates size*numelements bytes in memory
All elements in the allocated memory are set to zero
Returns a pointer to the allocated memory on success, or NULL on failure

\section*{Dynamic Memory Allocation}

Example: create an array of 10 integers int myArr[10];
- Malloc()

Example
int \(*_{m y A r r}=(i n t ~ *)\) malloc( 10 * sizeof(int) );
- Calloc()

Example
```

int *myArr = (int *) calloc( 10 , sizeof(int) );

```

\section*{Dynamic Memory Allocation}

\section*{Functions related to DMA are in the library stdlib.h}
```

void *realloc(void *ptr, size_t size)

```

Changes the size of the allocated memory block pointed by ptr to size
Returns a pointer to the allocated memory on success, or NULL on failure
```

void free(void *ptr)

```

De-allocates (frees) the space in memory pointed by ptr

\section*{Dynamic Memory Allocation}

Example: create an array of 10 integers, resize it to 15, then free the space in memory
1) int *myArr \(=(i n t *)\) malloc( 10 * sizeof(int) );
2) myArr = realloc( myArr, 15 * sizeof(int) );

3 ) free ( myArr );

1 )
2 )
3 )


\section*{Dynamic Memory Allocation}

Example: reading an indefinitely long command line

So far we have been reading strings from command line using an array
```

char line[100];
fgets( line, sizeof(line), stdin);

```

What if the user enters a command with 105 characters?

\section*{Dynamic Memory Allocation}

\section*{Multidimensional Arrays}
\(2 \times 3\) matrix of double

double **
double *
double

\section*{Dynamic Memory Allocation}

\section*{Multidimensional Arrays}
\(2 \times 3\) matrix of double
```

double** M = (double**) malloc( 2 * sizeof(double *) );
int i;
for ( i = 0 ; i<2; i++ ){
M[i] = malloc( 3 * sizeof(int) );
}
/* use M as a regular 2-dimensional array */
for ( i = 0 ; i<2; i++ ){
free( M[i] );
}
free( M );

```

\section*{Memory Leaks}

Space in the heap is LIMITED, therefore we must be careful and free memory

There are two cases in whish freeing memory becomes impossible:
- when we move a pointer after allocating memory
```

int N = 40000;
char *str = "Hello";
char *giantString = malloc(N*sizeof(char));
giantString = str;

```

Now we cannot find anymore the location of the block of allocated memory

\section*{Memory Leaks}

Space in the heap is LIMITED, therefore we must be careful and free memory

There are two cases in whish freeing memory becomes impossible:
- if we reallocate memory using the same pointer
```

int N = 40000;
char *giantString = malloc(N*sizeof(char));

```
/* do something */
giantString \(=\) malloc( \(N^{*}\) sizeof(char));
giantString now points to a newly allocated block of memory, the location of the previous one is lost

\section*{COMSW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

\section*{Review - Arrays of strings}
- An array Arr of 3 strings of variable length
```

char *Arr[3]={ "Hello", "World", "Wonderful" };
Arr[2] = Arr+2 // "Wondeful"

```
- An array Arr of 3 strings of maximum length \(=15\)
```

char Arr2[3][15] = { "Hello2", "World2", "Wonderful2" };

```
Arr2[0] = Arr2 // "Hello2"
Arr2[1] = Arr2+1 // "World2"

\section*{Program's Inputs}
- When we run a program, sometimes we want to pass some input arguments to it
- This can be done by writing them in the command line, immediately after the program name
- The program's inputs must be separated by spaces

Example
The program sumTwoNumbers sums two numbers.

We can pass the two input numbers directly when we invoke the program's executable (instead of the usual I/O operations, such as printing to command line the message "please insert two numbers:", followed by fgets() etc.)
./sumTwoNumbers 35

\section*{Command Line Arguments}
- Input parameters of the function main()
- argc, argv
```

int main( int argc, char* argv[] )

```
argc • Integer
- Specifies the number of arguments on the command line (including the program name)
argv • Array of strings
- Contains the actual arguments on the command line
- First element is the name of the program

\section*{Command line arguments}

It is a good habit, especially when a program takes input arguments, to specify in a header on the top of the main file:
- Program name and purpose
- Program usage: syntax to use to invoke (run) the program with input arguments
- Description of input arguments
- Description of output from the program

> It is common to add a -help option to print the relevant information about program usage and input arguments

\section*{Command line arguments}

\section*{Example}

Program calculator, reads two numbers, the operator, and prints the result

\section*{Linux Wildcard Characters}

Linux has a series of wildcard characters
* ? []
* Represents strings of arbitrary length containing any possible character
* all items (directories and files) - with or without a suffix
r* items beginning with the letter "r"
boot* items beginning with "boot"
*mem* all items contain "mem" anywhere in the name
*.png items having the suffix of ".png" - that end in ".png"

We must be very careful when we use wildcard characters as input, because argc and argv recognize them!

\section*{Linux Wildcard Characters}

\section*{Linux has a series of wildcard characters * ? [ ]}
? Represents one single character which has any possible value
?.txt items starting with only one character and ending in ".txt"
Examples: b.txt and 3.txt
memo?.sxw items beginning with "memo", having a single character after "memo", and having the suffix of ".sxw"
Examples: memo1.sxw and memoh.sxw - not memo23.sxw
memo??.sxw items beginning with "memo", having a two characters (only) after "memo", and having the suffix of ".sxw"
Examples: memo21.sxw and memok9.sxw - not memos.sxw

We must be very careful when we use wildcard characters as input, because argc and argv recognize them!

\section*{Linux Wildcard Characters}

\section*{Linux has a series of wildcard characters \\ * ? []}
[] Represents intervals of characters values
[a-z]* items that begin with any lower case letter and end in any other characters
[A-Z]-list.dat items that begin with any upper case letter and end in "-list.dat"
[a-zA-Z]report.sxc items that begin with any lower case or upper case letter and end in "report.sxc"
[e-t].c items that begin with any lower case letter between ' \(e\) ' and ' \(t\) ' and end in ". \(c\) "

We must be very careful when we use wildcard characters as input, because argc and argv recognize them!

\section*{Homework 3 Solution}

\section*{COMsW 1003-1}

\section*{Introduction to Computer Programming in \(\boldsymbol{E}\)}

\section*{Modular Programming}

\section*{Review - Header files}
- Header files are fundamentally libraries
- Their extension is .h
- They contain function definitions, variables declarations, macros
- In order to use them, the preprocessor uses the following code
\#include <nameOfHeader.h>
\#include "nameOfHeader. \(h\) " \(\longrightarrow\) For standard C libraries
- So far, we have used predefined \(C\) header files, but we can create our own! (more on this next week)

\section*{Modular Programming}
- So far we have seen only small programs, in one single file
- What about bigger programs? Need to keep them organized, especially if multiple people work on the same project
- They are organized in multiple, organized parts: MODULES

\section*{Modules}
- A module is "a collection of functions that perform related tasks" [PCPCh18]
- A module is basically a user defined library
- Two parts:
- Public : tells the user how to use the functions in the module. Contains declaration of data structures and functions
- Private : implements the functions in the module

\section*{Modules}
- Two parts:
- Public: tells the user how to use the functions in the module. Contains definition of data structures and functions
- Private : implements the functions in the module


Public

Private

\section*{Header}
- A header should contain:
- A section describing what the module does
- Common constants
- Common structures
- Public functions declarations
- Extern declarations for public variables

\section*{Function Declaration vs. Definition}
- All identifiers in C need to be declared before they are used, including functions
- Function declaration needs to be done before the first call of the function
- The declaration (or prototype) includes
- return type
- number and type of the arguments
- The function definition is the actual implementation of the function
- Function definition can be used as implicit declaration

\section*{Modules}

\section*{mainProgram.c calculator.h calculator.c}


\section*{Compile modules together}
- We need a way to "glue" the modules together
- We need to compile not only the main program file, but also the user defined modules that the program uses
- Solution : makefile

\section*{Makefile}
- make routine offered in UNIX (but also in other environments)
- make looks at the file named Makefile in the same folder and invokes the compiler according to the rules in Makefile

\section*{Makefile - Single file}

```


# Makefile for UNIX system

# using a GNU C compiler (gcc)

\#---------------------------------------------------------

# this is a comment

oldCalculator: oldCalculator.c
gcc -Wall -o oldCalculator oldCalculator.c

```

\section*{Makefile - Single file}

```


# Makefile for UNIX system

# using a GNU C compiler (gcc)

```

\# this is a comment \(\longrightarrow\) Comments start with a \# sign
oldCalculator: oldCalculator.c
    gcc-Wall-o oldCalculator oldCalculator.c

Rule: gcc command we are used to
The second statement MUST start with a TAB!

\section*{Makefile - Single file}

```


# Makefile for UNIX system

# using a GNU C compiler (gcc)

```

```

CC=gcc
CFLAGS=-Wall
oldCalculator: oldCalculator.c
\$(CC) \$(CFLAGS) - o oldCalculator oldCalculator.c

```

\section*{Makefile - Single file}

\(\mathrm{CC}=\mathrm{gCc}\)
\(\mathrm{CFLAGS}=-\mathrm{Wall}\)\(\longrightarrow\) macros
oldCalculator: oldCalculator.c


Rule: gcc command we are used to
The second statement MUST start with a TAB!

\section*{Makefile}
- Macros
```

name=data
\$ (name) \longrightarrowdata

```

Whenever \$(name) is found, it gets substituted with data
Same as object-type macros for Preprocessor
- Rules


UNIX compiles target from source using command Default command is \$(CC) \$(CFLAGS) -c source

\section*{Makefile - Single file}
\#
```Makefile for UNIX system\#
```

\# using a GNU C compiler (gcc) ..... \#
\# ..... \#
$\mathrm{CC}=\mathrm{gcc}$

```CFLAGS=-Wall
oldCalculator: oldCalculator.c
    $(CC) $(CFLAGS) -o oldCalculator oldCalculator.c
clean:
rm -f oldCalculator
```


## Makefile - Single file



CC=gcc
CFLAGS $=$ - Wall $\longrightarrow$ macros

The second statement MUST start with a TAB!
Rule: Clean up files

## Makefile

- If I have multiple rules, I can use the name of the target to execute only the rule I want
- By default, make executes only the first rule


## Example

\$make clean

## Makefile - Multiple Modules


$\mathrm{CC}=\mathrm{gcc}$
CFLAGS=-Wall
mainCalc : mainProgram.c calculator.o \$(CC) \$(CFLAGS) -o mainCalc mainProgram.c calculator.o
calculator.o : calculator.c calculator.h \$(CC) \$(CFLAGS) - c calculator.c
clean:
rm -f calculator.o mainProgram

## Makefile - Multiple Modules


$\mathrm{CC}=\mathrm{gcc}$
CFLAGS $=-$ Wall
maincalc : mainProgram.c calculator.o \$(CC) \$(CFLAGS) -o mainCalc mainProgram.c calculator.o
calculator.o: calculator.c calculator.h \$ (CC) \$(CFLAGS) - C calculator.c
clean:
rm -f calculator.o maincalc

We must use the -c option to compile a module instead of an executable!

## Makefile

- Rules

```
target: source [source2] [source3] ...
    command
    command2
    command3
        \vdots
```

UNIX compiles target from source using command Default command is \$(CC) \$(CFLAGS) - c source
make is smart: it compiles only modules that need it If target has already been compiled and source did not change, make will skip this rule


This rule instead is ALWAYS executed by the compiler, because source is not specified in the first line

## Extern/Static Variables

- Extern is used to specify that a variable or function is defined outside the current file

When same variable is used by different modules, extern is a way to declare a global variable which can be used in all modules

- Static is used to specify that a variable is local to the current file (for global variables)

Remember the use for local variables (Lec7): local static means permanent

## COMSW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

## Lecture 19

Spring 2011
Instructor: Michele Merler

## Basic Data Structures

## Basic Data Structures

- So far, the only data structures we have seen to store data have been arrays ( and structs )
- There are other (and potentially more useful) data structures that can be used
- Lists
- Trees
- Benefits:
- Dynamically grow and shrink is easy
- Search is faster


## Linked Lists

- A chain of elements
- First element is called HEAD
- Each element (called NODE) points to the next
- The last node does not point to anything
- Like a treasure hunt with clues leading one to another



## Pointers to structs

- Pointers can point to any type, including structs
- There is a particular way of accessing fields in a struct through a pointer: the -> operator

```
struct person {
    int age;
    char *name;
}
struct person p1={15, "Luke" };
struct person *ptr = &p1;
ptr->age=20; // (*ptr).age = 20;
printf("%s\n", ptr->name);
```


## Linked Lists

- Structure declaration for a node of a linked list

```
struct ll_node {
    int value;
    struct ll_node *next;
};
typedef struct ll_node node;
```



## Linked Lists Initialization

```
struct ll_node {
    int value;
    struct ll_node *next;
};
```

```
node *head = (node *) malloc(sizeof(node));
head->value = 0;
head->next = NULL;
```

- First node (HEAD) of the list is just a pointer to the list, it not counted as an actual node in the list
- Value set to 0 (could be any number, maybe a counter)
- The list is still empty, there is only HEAD, so next is NULL (end of the list)


## Linked Lists

## Initialization



```
node *head = (node *) malloc(sizeof(node));
head->value = 0;
head->next = NULL;
```

- First node (HEAD) of the list is just a pointer to the list, it not counted as an actual node in the list
- Value set to 0 (could be any number, maybe a counter)
- The list is still empty, there is only HEAD, so next is NULL (end of the list)


## Linked Lists <br> Insert node in front

```
struct ll_node {
    int value;
    struct ll_node *next;
};
```

```
int addNodeFront( int val, node *head ){
    node *newNode = (node *) malloc(sizeof(node));
    newNode->value = val;
    newNode->next = head->next;
    head->next = newNode;
    return 0;
}
```


## Linked Lists - Insert node in front

 int addNodeFront( int val, node *head )\{1) node *newNode $=($ node $*)$ malloc(sizeof(node));
2) newNode->value = val;
3) newNode->next = head->next;
4) head->next $=$ newNode; return 0;
addNodeFront( 7, head );

3
newNode


head


4


## Linked Lists - Insert node in front

int addNodeFront( int val, node *head ) \{

1) node *newNode = (node *) malloc(sizeof(node));
2) newNode->value = val;
3) newNode->next = head->next;
4) head->next = newNode; return 0;
\}

$$
\begin{aligned}
& \text { addNodeFront( } 7, \text { head ); } \\
& \text { addNodeFront( } 5 \text {, head ); }
\end{aligned}
$$



## Linked Lists <br> Insert node at position N

```
struct ll_node {
    int value;
    struct ll_node *next;
```

int addNode( int val, node *head, int pos ) \{

```
        node *newNode = (node*) malloc( sizeof(node) );
    newNode->value = val;
    int i;
    node *tmp = head;
    for(i=0 ; i<pos; i++)
    tmp = tmp->next;
```

newNode->next = tmp->next;
tmp->next = newNode;
return 0;
\}

## Linked Lists - Insert node at position N

```
int addNode( int val, node *head, int pos ){
```

1) node $*$ newNode $=($ node*) malloc ( sizeof(node) );
newNode->value = val;
2) node *tmp = head;

$$
\begin{aligned}
\text { for }(i=0 & ; i<p o s ; i++) \\
t m p & =\text { tmp->next }
\end{aligned}
$$

3) newNode->next = tmp->next;
4) tmp->next $=$ newNode;

> return 0;

1
newNode


```
```

addNode( 4, head, 2 );

```
```

```
```

addNode( 4, head, 2 );

```
```


## Linked Lists - Insert node at position N

```
int addNode( int val, node *head, int pos ){
```

2) node *tmp = head;

$$
\begin{aligned}
\text { for }(i=0 & ; i<p o s ; i++) \\
t m p & =\text { tmp->next }
\end{aligned}
$$

3) newNode->next = tmp->next;
4) tmp->next $=$ newNode;
return 0;
\}
addNode( 4, head, 2 );


## Linked Lists - Insert node at position N

```
int addNode( int val, node *head, int pos ){
```

$$
\text { node } * \text { tmp }=\text { head; }
$$

2) for(i=0 ; i<pos; i++)

$$
\text { tmp }=\text { tmp->next; }
$$

3) newNode->next $=$ tmp->next;
4) tmp->next $=$ newNode;
```
        return 0;
```

\}

$$
\text { addNode( 4, head, } 2 \text { ); }
$$



## Linked Lists <br> Delete Node

```
struct ll_node {
    int value;
    struct ll_node *next;
};
```

int removeNodePosition( node *head, int pos )\{
int i;
node $*$ tmp $=$ head;
for (i=0 ; i<pos; i++)
tmp $=$ tmp->next;
node* tmp2 = tmp->next;
tmp->next $=$ tmp->next->next;
free(tmp2);
return 0;

## Linked Lists - Delete Node

int removeNodePosition( node *head, int pos )\{ int i;

1) node *tmp = head;

$$
\begin{gathered}
\text { for (i=0 ; i<pos; i++) } \\
\text { tmp }=\text { tmp->next; }
\end{gathered}
$$

2) node* tmp2 = tmp->next;
```
tmp->next = tmp->next->next;
```

3) free(tmp2);
return 0;
\}
```
removeNode( head, 1 );
```

1


## Linked Lists - Delete Node

int removeNode( node *head, int pos )\{

```
int i;
```

1) node *tmp = head;

$$
\begin{gathered}
\text { for (i=0 ; i<pos; i++) } \\
\text { tmp }=\text { tmp->next; }
\end{gathered}
$$

2) node* tmp2 = tmp->next;
```
    tmp->next = tmp->next->next;
```

3) free(tmp2);
return 0;
\}
```
removeNode( head, 1 );
```



## Linked Lists - Delete Node

int removeNode( node *head, int pos )\{

```
int i;
```

1) node *tmp = head;

$$
\begin{gathered}
\text { for (i=0 ; i<pos; i++) } \\
\text { tmp }=\text { tmp->next; }
\end{gathered}
$$

2) node* tmp2 = tmp->next;
```
tmp->next = tmp->next->next;
```

3) free(tmp2);
return 0;
\}
```
removeNode( head, 1 );
```

```
removeNode( head, 1 );
```



## Linked Lists Delete Whole List

```
struct ll_node {
    int value;
    struct ll_node *next;
};
```

int destroyList( node **head )\{
node *tmp;
while( (*head)->next $!=$ NULL ) \{
top $=(*$ head $) ;$
$(*$ head $)=(*$ head $)->n e x t ;$
free(tmp);
\}
return 0;
\}
destroyList( \&head );

## Linked Lists Delete Whole List

```
struct ll_node {
    int value;
    struct ll_node *next;
};
```

int destroyList( node **head ) \{
node *tmp;
while ( $(*$ head $)->$ next $!=$
tmp $=(*$ head $) ;$
$(*$ head $)=(*$ head $)->$ next;
free $(t m p) ;$
I need to pass head by reference, because I am changing it within the function
destroyList( \&head );

## Doubly linked lists

- Pointer to next AND previous node
- Faster backtracking

```
struct dll_node {
    int value;
    struct dll_node *prev;
    struct dll_node *next;
};
```



## Binary Trees

- Like lists, but each node has a pointer to two elements:
- Left has a value < current node
- Right has a value > current node
- First node is called ROOT

```
struct t_node {
    int value;
    struct t_node *left;
    struct t_node *right;
};
```



## Binary Trees

- Left has a value < current node
- Right has a value > current node


## C



## Binary Trees

- Left has a value < current node
- Right has a value > current node



## Binary Trees

- Left has a value < current node
- Right has a value > current node


```
struct t_node {
int value;
struct t_node *left;
struct t_node *right;
};
```

Nodes at the bottom level or without children are called LEAVES


## Binary Trees

Inserting number x into a Binary Tree:

1. Start at root
2. if (current node is NULL)
create new node and set node's value to $\mathbf{x}$
3. else

if ( $x$ >= current node's value )<br>follow right pointer<br>else<br>follow left pointer

Go to 1

## Binary Trees

Example: [ 11262317908 ]


## Binary Trees

Example: [ 11262317908 ]


## Binary Trees

Example: [ 11262317908 ]


## Binary Trees

Example: [ 11262317908 ]


NULL


NULL NULL
NULL NULL
NULL

## Trees Definitions

- Root : node with no parents. Leaf : node with no children
- Depth (of a node) : path from root to node
- Level: set of nodes with same depth
- Height or depth (of a tree) : maximum depth
- Size (of a tree) : total number of nodes
- Balanced binary tree : depth of all the leaves differs by at most 1.



## Read PCP Chapter 17

## COMSW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

## Lecture 20

Instructor: Michele Merler

## Announcements

- HW5 out this Wednesday,
- Due on Wednesday, April 27th before class
- Final on Monday May $9^{\text {th }}$, from 9am to 12 pm, in class
- Same format as Midterm


## Today

- Quick review of linked lists
- Binary Trees
- Complexity Analysis


## Introduction to Complexity Analysis

## Measuring Algorithms

- In Computer Science, we are interested in finding a function that defines the quantity of some resource consumed by a particular algorithm
- This function is often referred to as a complexity of the algorithm
- The resources we usually investigate are
- running time
- memory requirements


## Measuring Algorithms

- We want to express complexity in the most general way possible
- Running time and space typically depend on input size

For varying input sizes, we can write time and space requirements as functions of $\boldsymbol{n}$.

- Algorithms run on different machines

For varying implementation, we use a description independent from constant factors.

## Example

Given an array $X$ of 10 elements of type int

| $X$ | 7 | 1 | 44 | 2 | 34 | 9 | 12 | 7 | 33 | 12 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Complexity analysis

- What is the running time (RT) of an algorithm that sums the elements in the array?
- How much space (SP) in memory is used by that algorithm?

```
int X[10];
int i, sum = X[0];
for(i=1; i<10;i++){
    sum += X[i];
```


## Example

Given an array $X$ of 10 elements of type int

| $X$ | 7 | 1 | 44 | 2 | 34 | 9 | 12 | 7 | 33 | 12 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Complexity analysis

- What is the running time (RT) of an algorithm that sums the elements in the array?
- How much space (SP) in memory is used by that algorithm?

| Machine 1 |
| :---: |
| Addition $\rightarrow 2$ seconds |
| int $\quad \rightarrow 4$ bytes |

$\mathrm{RT}=9 * 2=18$
$\mathrm{SP}=10 * 4+2 * 4=48$
$\mathrm{RT}=9 * 3=27$
$\mathrm{SP}=10 * 8+2 * 8=96$

| Machine 2 |  |
| ---: | :--- |
| Addition | $\rightarrow 3$ seconds |
| int $\quad \rightarrow 8$ bytes |  |

```
int X[10];
int i, sum = X[0];
for(i=1; i<10;i++){
        sum += X[i];
```

\}
$\mathrm{RT}=9 * 2=18$
$\mathrm{SP}=10 * 8+2 * 8=96$

## Example

Given an array of 10 elements of type int

| $X$ | 7 | 1 | 44 | 2 | 34 | 9 | 12 | 7 | 33 | 12 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Complexity analysis

- What is the running time (RT) of an algorithm that sums the elements in the array?
- How much space (SP) in memory is used by that algorithm?

This is not general!
Performance of machines, not of algorithm! What if array has $\boldsymbol{n}$ elements?

```
int X[10];
int i, sum = X[0];
```

for $(i=1 ; \quad i<10 ; i++)\{$

## Example

Given an array of 10 elements of type int

| $X$ | 7 | 1 | 44 | 2 | 34 | 9 | 12 | 7 | 33 | 12 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

Complexity analysis

- What is the running time (RT) of an algorithm that sums the elements in the array?
- How much space (SP) in memory is used by that algorithm?

```
int X[10];
int i, sum = X[0];
for(i=1; i<10;i++){
    sum += X[i];
```

\}

We want to express complexity of algorithm in terms of
$-n$ : number of elements in array (variable)

- c : number of seconds to execute addition (constant)
- b: number of bytes to store elements (constant)

$$
\begin{aligned}
& R T=c(n-1) \\
& S P=b(n+2)
\end{aligned}
$$

## Big - O Notation

GOAL: estimate the order of the function $f(n)$ that represents RT or SP in terms of $\boldsymbol{n}$

$$
f(n)=O(g(n))
$$

$$
\begin{aligned}
& f(n) \underset{n \rightarrow \infty}{=} O(g(n)) \\
& \Leftrightarrow \\
& \exists C>0 \text { and } n_{0}: \\
& |f(n)| \leq C|g(n)| \quad \forall n>n_{0}
\end{aligned}
$$

## Big - O Notation

GOAL: estimate the order of the function $f(n)$ that represents RT or SP in terms of $\boldsymbol{n}$

$$
f(n)=O(g(n))
$$

$f(n) \underset{n \rightarrow \infty}{=} O(g(n)) \quad f(n)$ equals oh of $g(n)$ as $n$ tends to infinity
$\longleftrightarrow$
$\exists C>0$ and $n_{0}:$
if and only if
there exists a positive constant $\mathbf{C}$ and a value $\boldsymbol{n}_{0}$ such that
for all $n$ greater than $n_{0}$, the absolute value of $f(n)$ is smaller than $\mathbf{C}$ times the absolute value of $\mathrm{g}(\mathrm{n})$

## Big - O Notation

GOAL: estimate the order of the function $f(\boldsymbol{n})$ that represents RT or SP in terms of $n$

$$
\begin{aligned}
& f(n)=O(g(n)) \\
& f(n) \underset{n \rightarrow \infty}{=} O(g(n)) \\
& \Leftrightarrow \\
& \exists C>O \text { and } n_{0}: \\
& |f(n)| \leq C \mid g(n) \quad \forall n>n_{0}
\end{aligned}
$$

In other words, big-O means less than some constant scaling When analyzing complexity with big-O notation, we always consider the WORST CASE SCENARIO

## Big-O notation: Examples

- $f(n)=3 n^{4}+7 n^{2}-5 n+8$

$$
\begin{aligned}
\left|3 n^{4}+7 n^{2}-5 n+8\right| & \leq 3 n^{4}+7 n^{2}+|5 n|+8 \\
& \leq 3 n^{4}+7 n^{4}+5 n^{4}+8 n^{4} \\
& \leq 23 n^{4} \\
|\mathrm{f}(\mathrm{n})| & \leq \mathrm{C}|\mathrm{~g}(\mathrm{n})| \\
\mathrm{f}(\mathrm{n}) & =\mathrm{O}\left(\mathrm{n}^{4}\right)
\end{aligned}
$$

- What is the running time (RT) of an algorithm that sums $n$ elements in an array?

$$
C(n-1)=O(n-1)=O(n)
$$

## Big - O : common cases

## O(1) - constant time

- The algorithm requires the same fixed number of steps regardless of the size of the task
- Example: insert an element in front of a linked list
int addNodeFront( int val, node *head )\{

1) node *newNode = malloc(sizeof(node));
2) newNode->value = val;
3) newNode->next = head->next;
4) head->next $=$ newNode;

No matter how long the list is, this operation always requires 4 steps
$\mathrm{O}(4)=\mathrm{O}(1)$

## Big - O : common cases

## O(1) - constant time

$$
\begin{gathered}
\text { if } c \ll n \\
O(c)=O(1)
\end{gathered}
$$

- The algorithm requires the same fixed number of steps regardless of the size of the task
- Example: insert an element in front of a linked list
int addNodeFront( int val, node *head )\{

1) node *newNode $=$ malloc(sizeof(node));
2) newNode->value = val;
3) newNode->next = head->next;
4) head->next = newNode;

No matter how long the list is, this operation always requires 4 steps RT $=\mathrm{O}(4)=\mathrm{O}(1)$

## Big - O : common cases

## $O(n)$ - linear time

- The algorithm requires a number of steps proportional to the size of the task
- Examples:
- Travers a linked list or an array with $\boldsymbol{n}$ elements;
- Find the maximum and minimum element in a list or array

```
for(i=0 ; i < n; i++){
    if(arr[i] < minVal)
        minVal = arr[i];
    if(arr[i] > maxVal)
        RT=O(2n)=O(n)
        SP=O(n+2)=O(n)
    maxVal = arr[i];
```


## Big - O : common cases

## $\mathrm{O}\left(\mathrm{n}^{2}\right)$ - quadratic time

- The number of operations is proportional to the size of the task squared.
- Example: Finding duplicates in an unsorted list of $n$ elements

```
for(i=0 ; i < n; i++){
    for(j=0 ; j < n; j++){
    if( (i!=j) && arr[i] == arr[j] )
        dup[i][j] = 1;
```

    \}
    $$
\mathrm{RT}=\mathrm{O}\left(4 \mathrm{n}^{2}+\mathrm{n}\right)=\mathrm{O}\left(\mathrm{n}^{2}\right)
$$

            increment \(\mathrm{i} \longrightarrow \mathrm{n}\) times
            increment \(j \longrightarrow n^{2}\) times
            check i! \(=j \longrightarrow n^{2}\) times
            check arr[i]==arr[j] \(\longrightarrow n^{2}\) times
        dup [i][j]=1 \(\longrightarrow(n-1)^{*}(n-1)\) times
    
## Big - O : common cases

## $\mathrm{O}(\log (\mathrm{n}))$ - logarithmic time

- Example: Find operation in a balanced binary tree with $\boldsymbol{n}$ nodes

$$
n=15
$$

$$
\begin{aligned}
& n=15 \\
& \text { height of tree }=3=\left\lfloor\log _{2}(n)\right\rfloor
\end{aligned}
$$

## Big - O : common cases

## $O(n \log (n))-" n \log (n) "$ time

- Examples: sorting algorithms (will see in next class)
- quicksort
- mergesort


## Big - O : common cases

## $O\left(a^{n}\right)$ - exponential time

 $a>1$- Example: Recursive Fibonacci implementation

```
int fib( int n ) {
```

```
switch(n) {
    case 0:
        return(0);
    case 1:
        return(1);
    default:
        return(fib(n-1) + fib(n-2));
}
```

How many times is fib() called?
Cost of fib() without return statement = $2=0(1)$

$$
\begin{aligned}
& \mathrm{RT}(\mathrm{n})=\mathrm{RT}(\mathrm{n}-1)+\mathrm{RT}(\mathrm{n}-2)+\mathrm{O}(1) \\
& \mathrm{RT}=\mathrm{O}\left(\mathrm{a}^{\mathrm{n}}\right) \\
& \quad a^{n}=a^{n-1}+a^{n-2} \\
& \quad a^{2}=a+1 \\
& \quad a=\frac{1+\sqrt{5}}{2} \approx 1.6
\end{aligned}
$$

## Big-O : Relationship among common cases

## $O(1)<O(\log n)<O(n)<O(n \log n)<O\left(n^{2}\right)<O\left(n^{3}\right)<O\left(a^{n}\right)$

Example: big-O when a function is the sum of several statements

```
int i=0;
for(i=0 ; i < n; i++){
    for(j=0 ; j < n; j++){
        if( (i!=j) && arr[i] == arr[j]) increment i
        dup[i][j] = 1;
    }
}
    RT=O(4n2+n)=O(n' )
        increment j
    check i!=j
    check arr[i]==arr[j]
    dup[i][j] = 1
```

Longest operation dominates (worst case)

## COMSW 1003-1

## Introduction to Computer Programming in $\boldsymbol{E}$

Lecture 21
Spring 2011
Instructor: Michele Merler

## Big-O : Relationship among common cases

## $O(1)<O(\log n)<O(n)<O(n \log n)<O\left(n^{2}\right)<O\left(n^{3}\right)<O\left(a^{n}\right)$

Example: big-O when a function is the sum of several statements

```
int i=0;
for(i=0 ; i < n; i++){
    for(j=0 ; j < n; j++){
        if( (i!=j) && arr[i] == arr[j]) increment i
        dup[i][j] = 1;
    }
}
    RT=O(4n+ +n)=O(n' )
        increment j
        check i!=j
        check arr[i]==arr[j]
    dup[i][j] = 1
```

        Longest operation dominates (worst case)
    
## Sorting

## Sorting

- Given a set of N elements, put them in order according to some criteria (alphabetical, relevance, date, smallest to largest, etc.)
- One of the most studied problems in Computer Science
- Everybody uses it every day



## Sorting

- Given a set of N elements, put them in order according to some criteria
- Compare pairs of elements
- Many algorithms, some of the most famous are:
- Bubble sort
- Selection sort
- Insertion sort
- Merge sort
- Counting sort
- In following examples, we'll see smallest to biggest sorting


## Bubble Sort

1. Start with the first two elements
2. If first element > second element

- Swap

3. Iterate for all following pairs
4. Repeat steps 1 to 3 until no swaps are necessary

$$
\text { Complexity }=0\left(n^{2}\right)
$$

Count number of comparisons and swaps

## Bubble Sort



## Bubble Sort



## Selection Sort

- Smarter algorithm, but same complexity (worst case)

1. Find smallest unsorted element
2. Swap with first unsorted element
3. Repeat steps 1 and 2 until no more unsorted elements

Complexity $=\mathbf{O}\left(\mathrm{n}^{2}\right)$

First unsorted element

## Selection Sort



## Insertion Sort

- Main idea: keep 2 separate sets (one sorted, one unsorted), and move elements from unsorted to sorted set one at a time
- Better performance in case many elements are already sorted, quadratic in worst case

1) Initialize 2 sets

- One set of sorted elements (contains only first element in the array)
- One set of unsorted elements (all the other elements in the array)

2) A) Take first element in unsorted set and
B) Insert it into sorted set at proper position
3) Repeat steps 2A) and 2B) until unsorted set is empty

Complexity $=\mathbf{O}\left(\mathrm{n}^{2}\right)$

## Insertion sort

| 1) | ded | Unsorted set |  |  |  | 1 check |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 9 | 5 | 3 | 4 | 1 |  |
| 2A) | 9 | 5 | 3 | 4 | 1 |  |
| 2B) | 5 | 9 | 3 | 4 | 1 | 1 swap |
| 2A) | 5 | 9 | 3 | 4 | 1 |  |
|  | 5 | 3 | 9 | 4 | 1 | 2 checks 2 swaps |
|  | 3 | 5 | 9 | 4 | 1 |  |
| 2A) | 3 | 5 | 9 | 4 | 1 | $\mathrm{n}-2$ checks <br> n-3 swaps |
| 2B) | 3 | 5 | 4 | 9 | 1 |  |
|  | 3 | 4 | 5 | 9 | 1 |  |
| 2A) | 3 | 4 | 5 | 9 | 1 |  |
|  |  | 3 | 4 | 5 | 9 |  |

## Merge Sort

- One of the fastest algorithms, divide and conquer principle
- Uses recursion
- Sorting small sets is faster than sorting large sets
- Merging 2 sets into a sorted union is faster if the sets are already sorted

1. If set H has 1 element, stop
2. else

- Split set into 2 halves H 1 and H 2 of (approximately) same size
- Sort H1 and H 2 with merge sort recursion
- Merge the sorted H1 and H2 into a sorted set

Complexity $=\mathbf{O}(n \log (n))$

## Merge Sort



## Merge Sort



## Merge Sort

Similar to trees, we perform $\log _{2}(n)$ splits and merges
Each merge takes $\mathrm{O}(\mathrm{n})$ in the worst case

## Merge routine:

Given H 1 and H 2 of size n 1 and n 2 respectively, create H of length $\mathrm{n}=\mathrm{n} 1+\mathrm{n} 2$

```
int c1=0, c2=0;
for (i=0; i<n; i++){
    if( (c1<n1) && ((H1[c1] < H2[c2]) || (c2==n2)) ){
                        H[i] = H1[c1];
                        c1++;
    }
    else{
        H[i] = H2[c2];
        c2++;
    }
```

\}


## Counting sort

- Intuition: exploit range $k$ of values in set
- Efficient if $k$ is not much larger than $n$

1. Find biggest and smallest values in the set ( $k=$ maxVal - minVal +1 )
2. Create an array $C$ of $k$ elements
3. Count occurrences $C(i)$ of each value $i$ in the set
4. Fill ordered set by inserting $C(i)$ elements of value $i$, for each value in range $k$

$$
\text { Complexity = } 0(n+k)
$$

## Counting sort

Example: range of values in set is $[1,5], k=5$


## Homework 4 Solution

