

COMSW 1003-1

Introduction to Computer Programming in

Lecture 22

Spring 2011

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Today

- Quicksort
- Pointers to functions, implementation of `qsort()`
- HW4 solution

Review – Bubble Sort

1. Start with the first two elements
2. If first element > second element
 - Swap
3. Iterate for all following pairs
4. Repeat steps 1 to 3 until no swaps are necessary

Complexity = $O(n^2)$

Count number of comparisons and swaps

Review - Selection Sort

- Smarter algorithm, but same complexity (worst case)
 1. Find smallest unsorted element
 2. Swap with first unsorted element
 3. Repeat steps 1 and 2 until no more unsorted elements

Complexity = $O(n^2)$

Review - Merge Sort

- One of the fastest algorithms, divide and conquer principle
 - Uses recursion
 - Sorting small sets is faster than sorting large sets
 - Merging 2 sets into a sorted union is faster if the sets are already sorted
1. If set H has 1 element, stop
 2. else
 - Split set into 2 halves H1 and H2 of (approximately) same size
 - Sort H1 and H2 with merge sort recursion
 - Merge the sorted H1 and H2 into a sorted set

Complexity = $O(n \log(n))$

Review - Counting sort

- Intuition: exploit range k of values in set
 - Efficient if k is not much larger than n
1. Find biggest and smallest values in the set
($k = \text{maxVal} - \text{minVal} + 1$)
 2. Create an array C of k elements
 3. Count occurrences $C(i)$ of each value i in the set
 4. Fill ordered set by inserting $C(i)$ elements of value i , for each value in range k

Complexity = $O(n + k)$

Quicksort

- Divide and conquer idea (similar to merge sort)
- In real world cases, on average it is as fast or faster than $O(n \log(n))$ algorithms

Complexity = $O(n^2)$

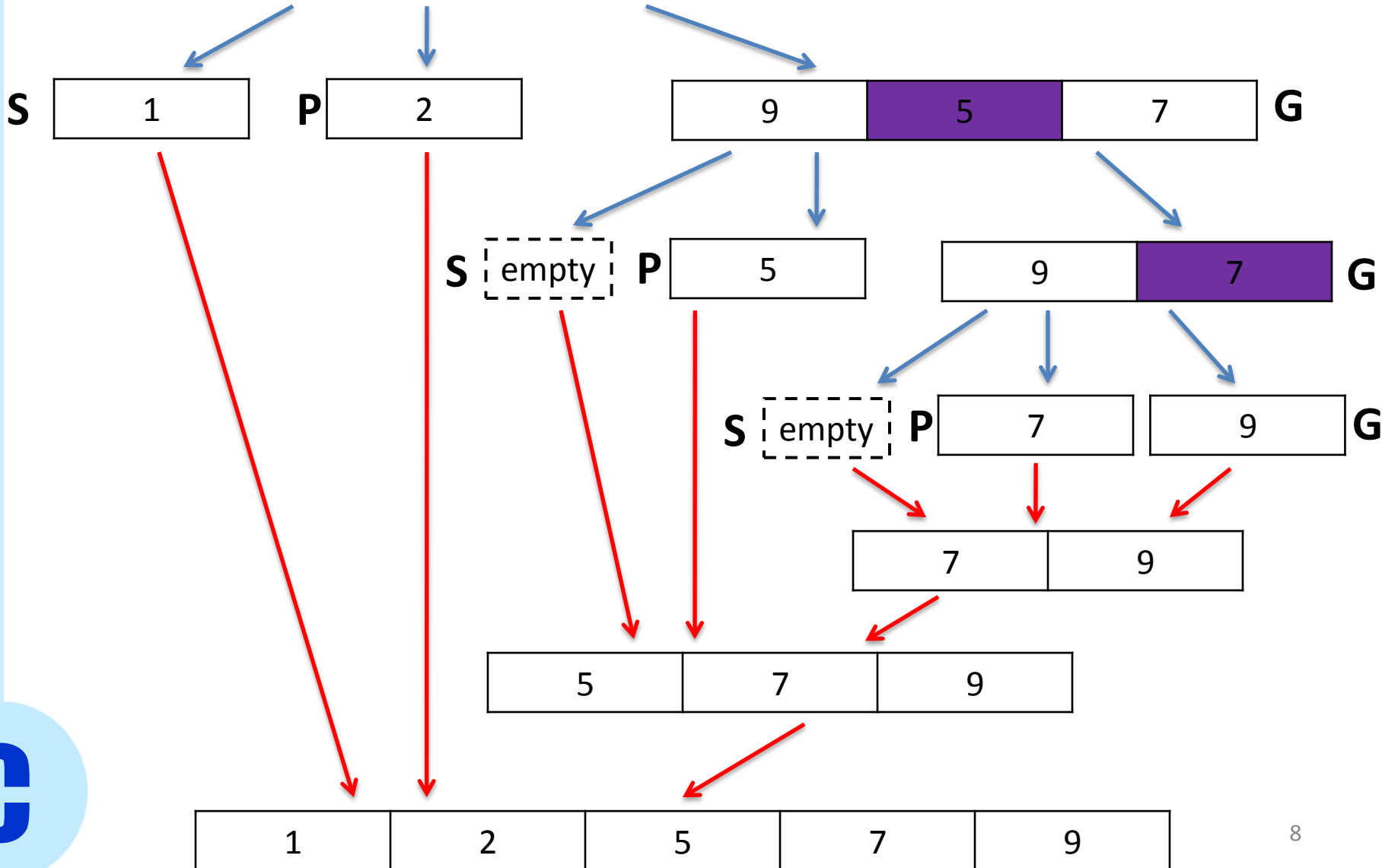
1. Choose an element in the array called **pivot P** and remove it from the array (common choice is median of first, middle and last element)
2. For each element x in the array (minus pivot)
if($x < \text{pivot}$)
 insert x in set S of elements smaller than pivot
else
 insert x in set G of elements greater than pivot
3. return(concatenate(quicksort(S), P , quicksort(G)));

C

Recursive call!

Quicksort

Pivot = median(9,1,2) = 2



Pointers to functions

Pointers to functions

- It is occasionally useful to use pointers to functions
- Since functions are stored in memory, we can reason about their addresses too
- This allows us to say, “run the function at address N on these arguments”
- Useful for being truly general, e.g. `stdlib qsort`

Pointers to functions

```
int (*f_ptr)(); // pointer to function that returns an int
```

Parentheses are important! Without parentheses, **f_ptr** looks like it returns a pointer to an int.

```
int (*f_ptr)(int, int); // pointer to a function
```

```
int greater_than(int a, int b); // function declaration
```

```
f_ptr = greater_than;
```

Pointers to functions

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int (*f_ptr)(int, int);  
  
int greater_than(int a, int b);  
  
f_ptr = greater_than;
```

```
int *ptr;  
  
int x[2];  
  
ptr = x;
```

qsort

- `qsort()` is a general sorting function, defined in `stdlib.h`
- Sort an array of any type, using any comparison criterion
- Define that comparison as a function pointer

```
void qsort(void *base, size_t n, size_t size,  
           int (*cmp)(const void *, const void *));
```

Depending on what function `cmp` points to, `qsort` uses a different criterion to sort the data

qsort

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The compare function should take two entries x and y , and return

+1 if $x > y$

-1 if $x < y$

0 if $x == y$