

Data Structures and Algorithms

Session 13. March 4, 2009

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Announcements

- * Homework 3 is out. Due 3/9
- * Sample midterm problems on Courseworks
- * Midterm review March 9th
- * Midterm Exam March 11th

Review

- * Solving the Young Tableaux Recurrences
- * buildHeap description and analysis
- * HW2 solutions

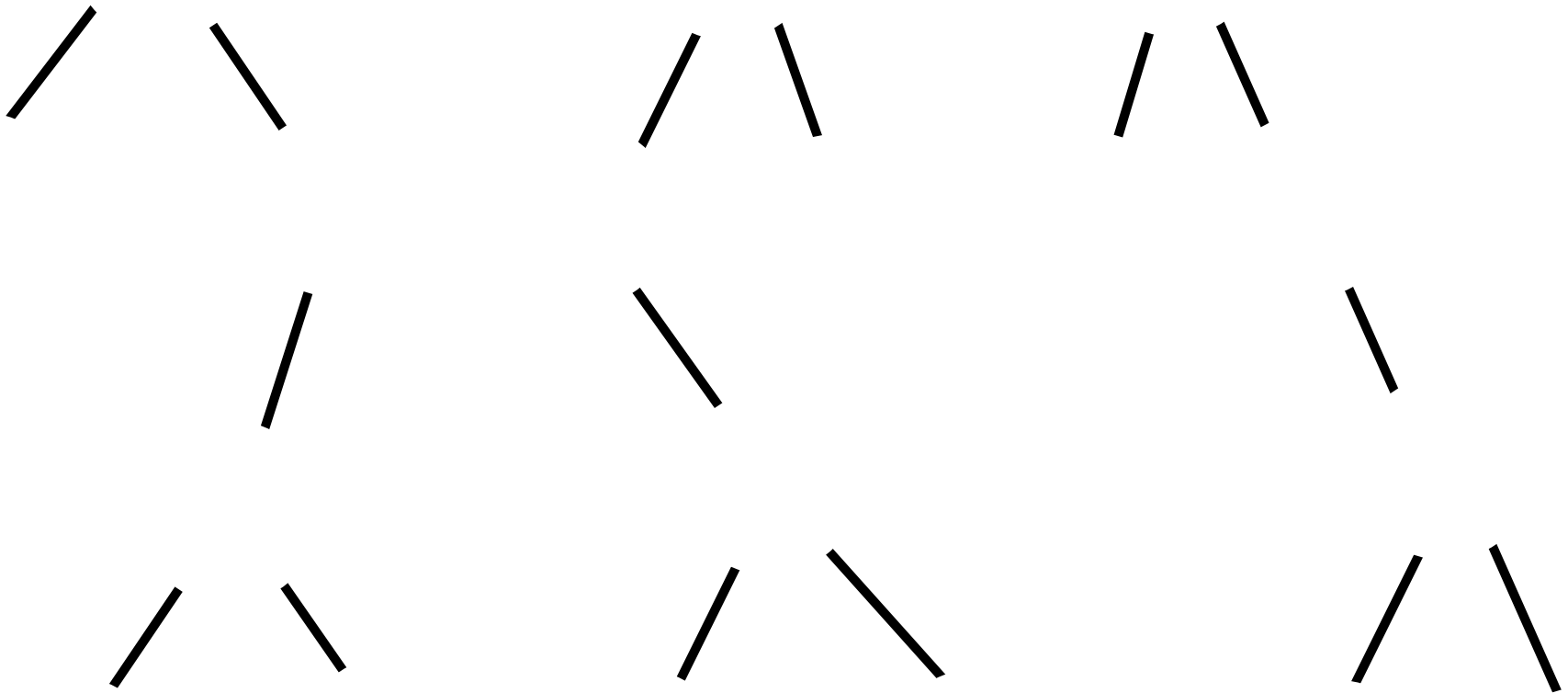
Today's Plan

- * Clarification about isomorphism
- * buildHeap example
- * HeapSort and HeapSelect

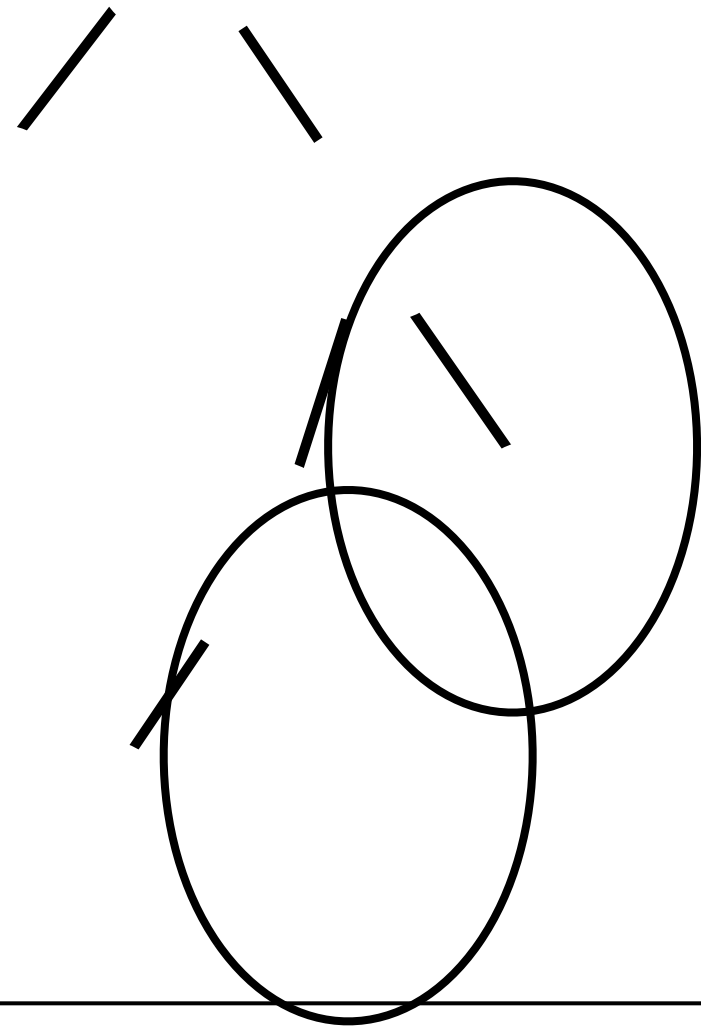
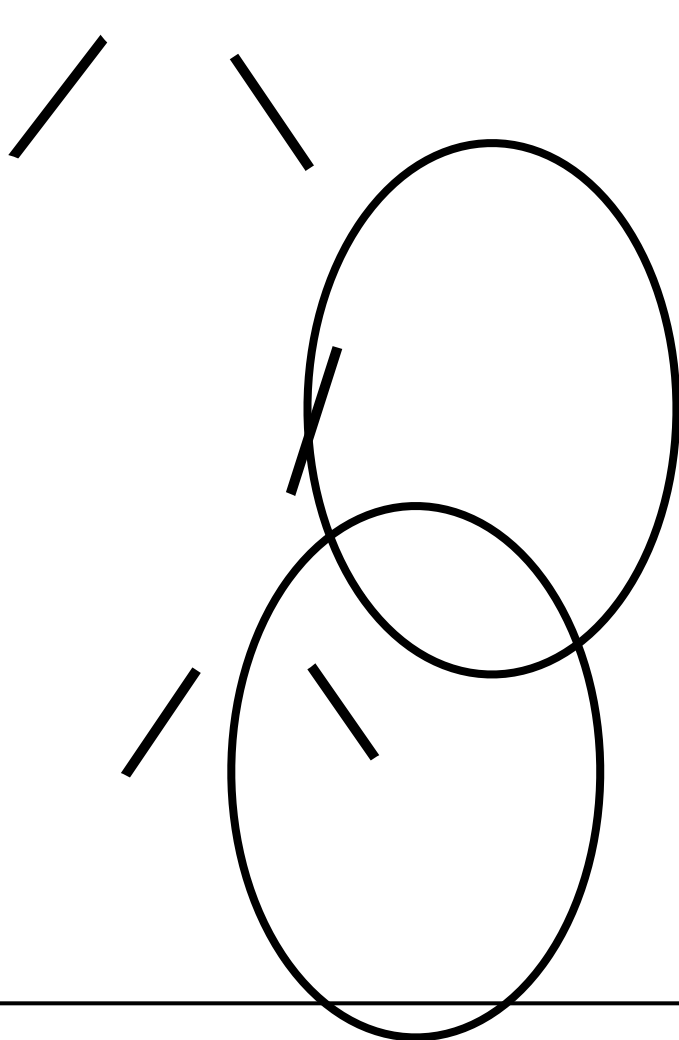
Isomorphism in Trees

- * *iso* means equal, *morph* means shape
- * Isomorphic means “having the same shape”
- * Ignore left and right, element data
- * Consider only *structural* properties:
 - * How many children? How many siblings?
How many cousins? How many grandchildren?

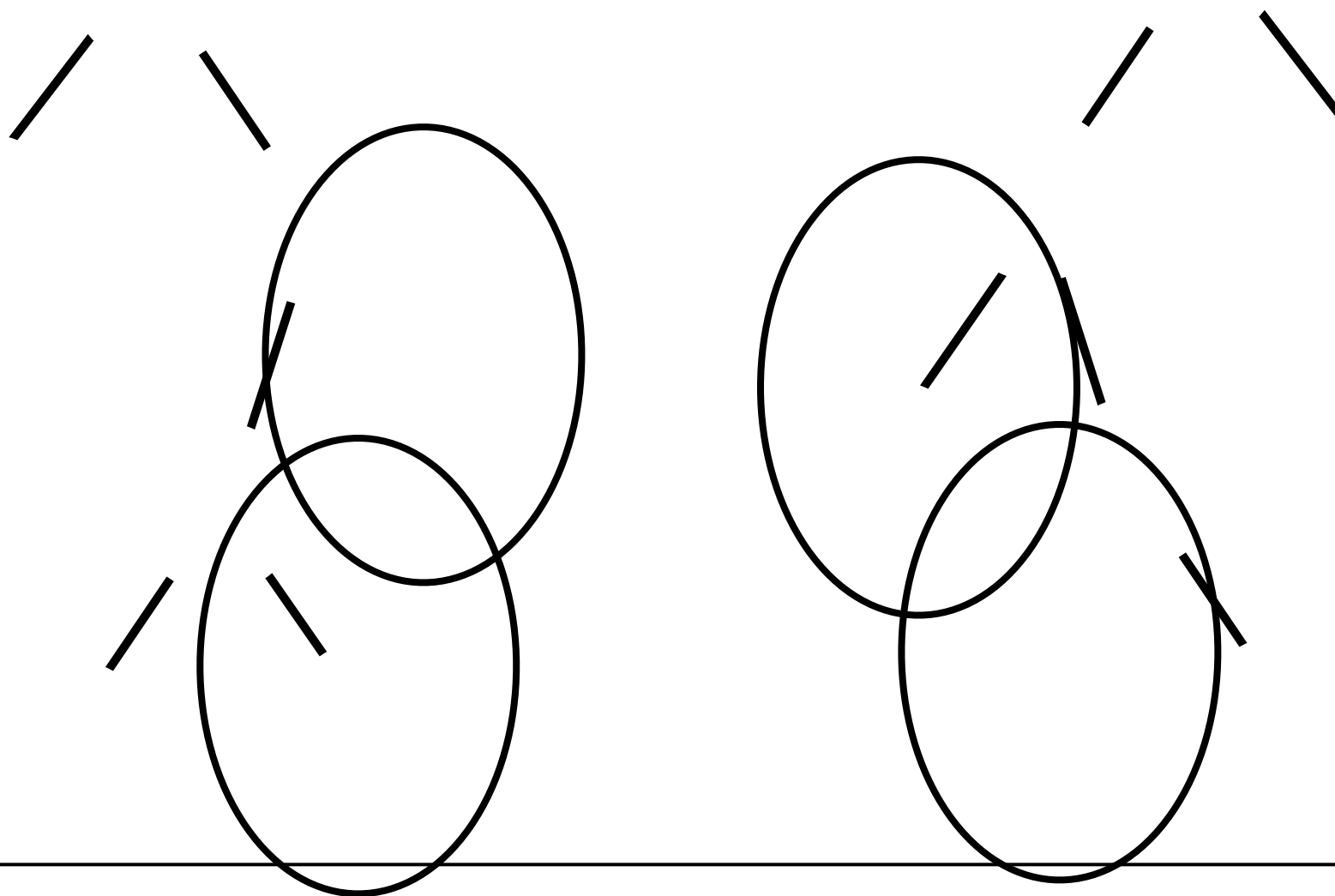
Isomorphic Binary Trees



Not Isomorphic Trees



Not Isomorphic Trees



buildHeap

- * Start at deepest non-leaf node
 - * in array, this is node $N/2$
- * **percolateDown** on all nodes in reverse level-order
 - * for $i = N/2$ to 1
 percolateDown(i)

| | | | | | | | | | | | | | | |
|---|---|----|---|----|---|---|----|---|---|---|----|---|----|----|
| 6 | 3 | 11 | 7 | 14 | 8 | 5 | 15 | 1 | 2 | 4 | 13 | 9 | 10 | 12 |
|---|---|----|---|----|---|---|----|---|---|---|----|---|----|----|

Selection

- * Recall the selection problem: $\text{findKth}(k, A[])$
 - * find the k th smallest element in A
- * Old method:
 - * Sort the array, then return the k th element
 - * Sort the first k elements, insert the remaining $(N-k)$ elements in the proper place

HeapSelect

- * run buildHeap on the array A
 - * $O(N)$
- * call deleteMin() k times
 - * $k O(\log N) = O(k \log N)$
- * HeapSelect runs in $O(N+k \log N)$

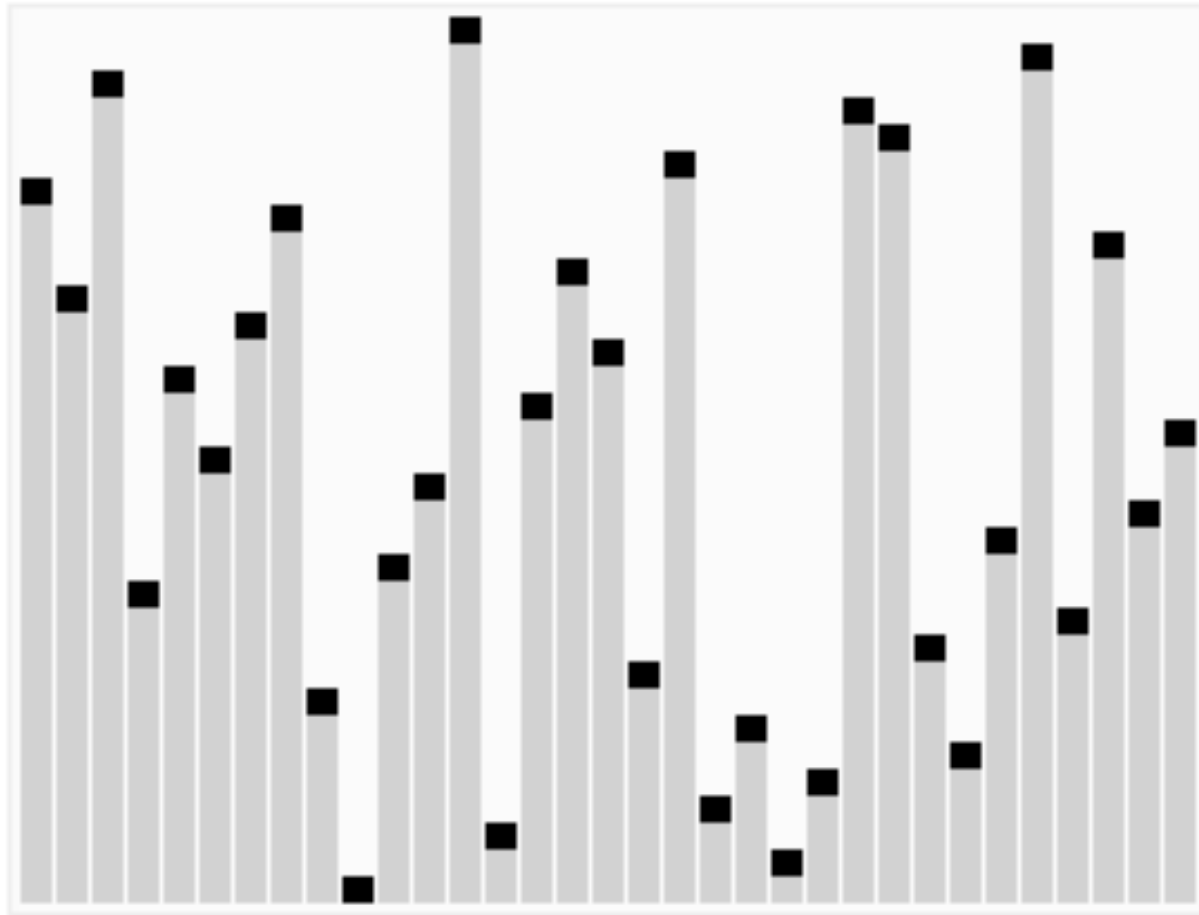
HeapSort

- * Naturally, we can use this idea to sort the array
- * Method 1:
 - * buildHeap on the array
 - * copy output of deleteMin into new array N times
 - * buildHeap costs $O(N)$, deleteMin costs $O(\log N)$
 - * Heapsort 1 costs $O(N+N \log N)=O(N \log N)$

HeapSort in Place

- * We don't need to allocate a new array
- * Instead, use a **max-heap**
 - * Reverse the heap order property: deleteMax
- * After each deleteMax, heap size is 1 less
 - * Stick the extracted max in the freed space

HeapSort Animation



Downloaded from <http://en.wikipedia.org/wiki/Heapsort>

Assignments

- * Continue HW3
- * Weiss 7.5 if you want to read about HeapSort
- * Practice midterm samples