

Fundamentals of Computer Systems

Thinking Digitally

Stephen A. Edwards and Martha Kim

Columbia University

Fall 2012

The Subject of this Class

0

The Subjects of this Class

0

1

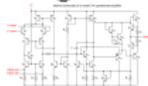
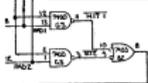
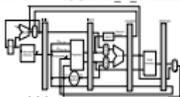
But let your communication be, Yea, yea; Nay, nay: for whatsoever is more than these cometh of evil.

— Matthew 5:37

Engineering Works Because of Abstraction



```
;; voice 1 wave select
ld a, (#CHI_W_NUM)
and a, (#CHI_W_SEL)
ld a, (#CHI_E_TABLE0)
jr nz, #00b4
ld a, (#CHI_E_TABLE0)
```



Application Software

Operating Systems

Architecture

Micro-Architecture

Logic

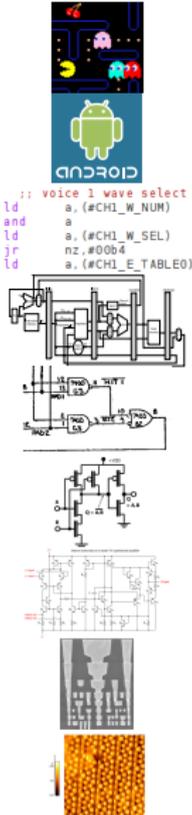
Digital Circuits

Analog Circuits

Devices

Physics

Engineering Works Because of Abstraction



Application Software	COMS 3157, 4156, et al.
Operating Systems	COMS W4118
Architecture	Second Half of 3827
Micro-Architecture	Second Half of 3827
Logic	First Half of 3827
Digital Circuits	First Half of 3827
Analog Circuits	ELEN 3331
Devices	ELEN 3106
Physics	ELEN 3106 et al.

Boring Stuff

Mailing list: csee3827-staff@lists.cs.columbia.edu

<http://www.cs.columbia.edu/~sedwards/classes/2012/3827-spring/>

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Lectures 10:10–11:25 AM, Tue, Thur, 207 Mathematics
Sep 4–Dec 6
Holidays: Nov 6 (Election Day), Nov 22 (Thanksgiving)

Assignments and Grading

Weight	What	When
40%	Six homeworks	See Webpage
30%	Midterm exam	October 23rd
30%	Final exam	During Finals Week (Dec 14–21)

Homework is due at the beginning of lecture.

We will drop the lowest of your six homework scores;

you can { skip
omit
forget
ignore
blow off
screw up
feed to dog
flake out on
sleep through } one with no penalty.

Rules and Regulations

You may collaborate with classmates on homework.

Each assignment turned in must be unique; work must ultimately be your own.

List your collaborators on your homework.

Don't cheat: if you're stupid enough to try, we're smart enough to catch you.

Tests will be closed-book with a one-page “cheat sheet” of your own devising.

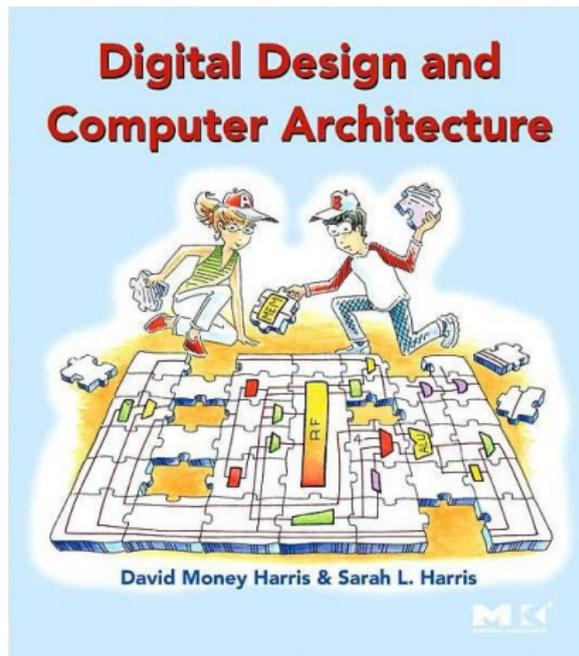
The Text

David Harris and Sarah Harris.

Digital Design and Computer Architecture.

Morgan-Kaufmann, 2007.

Almost precisely right for the scope of this class: digital logic and computer architecture



GILDAN
ULTRA
COTTON

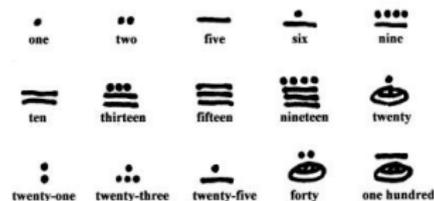
There are only 10 types
of people in the world:
Those who understand binary
and those who don't.

Which Numbering System Should We Use?

Some Older Choices:



Roman: I II III IV V VI VII VIII IX X



Mayan: base 20, Shell = 0

1	𐎀	11	𐎁𐎀	21	𐎁𐎁𐎀	31	𐎁𐎁𐎁𐎀	41	𐎁𐎁𐎁𐎁𐎀	51	𐎁𐎁𐎁𐎁𐎁𐎀
2	𐎀𐎀	12	𐎁𐎀𐎀	22	𐎁𐎁𐎀𐎀	32	𐎁𐎁𐎁𐎀𐎀	42	𐎁𐎁𐎁𐎁𐎀𐎀	52	𐎁𐎁𐎁𐎁𐎁𐎀𐎀
3	𐎀𐎀𐎀	13	𐎁𐎀𐎀𐎀	23	𐎁𐎁𐎀𐎀𐎀	33	𐎁𐎁𐎁𐎀𐎀𐎀	43	𐎁𐎁𐎁𐎁𐎀𐎀𐎀	53	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀
4	𐎀𐎀𐎀𐎀	14	𐎁𐎀𐎀𐎀𐎀	24	𐎁𐎁𐎀𐎀𐎀𐎀	34	𐎁𐎁𐎁𐎀𐎀𐎀𐎀	44	𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀	54	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀
5	𐎀𐎀𐎀𐎀𐎀	15	𐎁𐎀𐎀𐎀𐎀𐎀	25	𐎁𐎁𐎀𐎀𐎀𐎀𐎀	35	𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀	45	𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀	55	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀
6	𐎀𐎀𐎀𐎀𐎀𐎀	16	𐎁𐎀𐎀𐎀𐎀𐎀𐎀	26	𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀	36	𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀	46	𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀	56	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀
7	𐎀𐎀𐎀𐎀𐎀𐎀𐎀	17	𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀	27	𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀	37	𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀	47	𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀	57	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀
8	𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	18	𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	28	𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	38	𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	48	𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	58	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀
9	𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	19	𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	29	𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	39	𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	49	𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	59	𐎁𐎁𐎁𐎁𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀
10	𐎀	20	𐎁	30	𐎁𐎀	40	𐎁𐎀	50	𐎁𐎀		

Babylonian: base 60

The Decimal Positional Numbering System

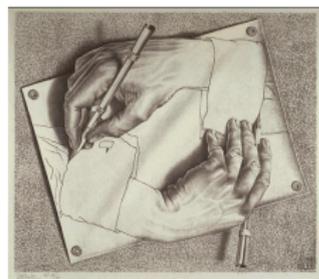


Ten figures: 0 1 2 3 4 5 6 7 8 9

$$7 \times 10^2 + 3 \times 10^1 + 0 \times 10^0 = 730_{10}$$

$$9 \times 10^2 + 9 \times 10^1 + 0 \times 10^0 = 990_{10}$$

Why base ten?



Hexadecimal, Decimal, Octal, and Binary

Hex	Dec	Oct	Bin
0	0	0	0
1	1	1	1
2	2	2	10
3	3	3	11
4	4	4	100
5	5	5	101
6	6	6	110
7	7	7	111
8	8	10	1000
9	9	11	1001
A	10	12	1010
B	11	13	1011
C	12	14	1100
D	13	15	1101
E	14	16	1110
F	15	17	1111

Binary and Octal



DEC PDP-8/I, c. 1968

Oct	Bin
0	0
1	1
2	10
3	11
4	100
5	101
6	110
7	111

$$\begin{aligned} \text{PC} &= 0 \times 2^{11} + 1 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + 1 \times 2^7 + 0 \times 2^6 + \\ & 1 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= 2 \times 8^3 + 6 \times 8^2 + 7 \times 8^1 + 5 \times 8^0 \\ &= 1469_{10} \end{aligned}$$

Hexadecimal Numbers

Base 16: 0 1 2 3 4 5 6 7 8 9 A B C D E F

Instead of groups of 3 bits (octal), Hex uses groups of 4.

$$\begin{aligned} \text{CAF EF00D}_{16} &= 12 \times 16^7 + 10 \times 16^6 + 15 \times 16^5 + 14 \times 16^4 + \\ &\quad 15 \times 16^3 + 0 \times 16^2 + 0 \times 16^1 + 13 \times 16^0 \\ &= 3,405,705,229_{10} \end{aligned}$$

C	A	F	E	F	0	0	D		Hex		
11001010111111101111000000001101									Binary		
3	1	2	7	7	5	7	0	0	1	5	Octal

Computers Rarely Manipulate True Numbers

Infinite memory still very expensive

Finite-precision numbers typical

32-bit processor: naturally manipulates 32-bit numbers

64-bit processor: naturally manipulates 64-bit numbers

How many different numbers can you

represent with 5

binary	
octal	
decimal	digits?
hexadecimal	

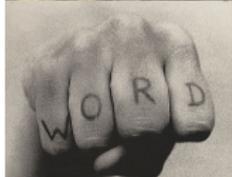
Jargon



Bit Binary digit: 0 or 1



Byte Eight bits



Word Natural number of bits for the processor, e.g., 16, 32, 64



LSB Least Significant Bit ("rightmost")



MSB Most Significant Bit ("leftmost")

Decimal Addition Algorithm

$$\begin{array}{r} 434 \\ +628 \\ \hline \end{array}$$

$$4 + 8 = 12$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1 \\ 434 \\ + 628 \\ \hline 2 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1 \\ 434 \\ + 628 \\ \hline 62 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1\ 1 \\ 434 \\ +628 \\ \hline 062 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1\ 1 \\ 434 \\ +628 \\ \hline 1062 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Binary Addition Algorithm

$$\begin{array}{r} 10011 \\ +11001 \\ \hline \end{array}$$

$$1 + 1 = 10$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 1 \\ 10011 \\ +11001 \\ \hline 0 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 11 \\ 10011 \\ +11001 \\ \hline 00 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 011 \\ 10011 \\ +11001 \\ \hline 100 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \\ 0 + 0 + 1 = 01 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 0011 \\ 10011 \\ +11001 \\ \hline 1100 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \\ 0 + 0 + 1 = 01 \\ 0 + 1 + 1 = 10 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 10011 \\ 10011 \\ +11001 \\ \hline 101100 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \\ 0 + 0 + 1 = 01 \\ 0 + 1 + 1 = 10 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Signed Numbers: Dealing with Negativity

A rectangular image showing a handwritten signature in cursive script. The signature reads "John Hancock" and is written in dark ink on a light-colored, slightly textured paper background. The signature is centered horizontally within the image.

How should both positive and negative numbers be represented?

Signed Magnitude Numbers

You are most familiar with this:
negative numbers have a leading –

In binary, a
leading 1 means
negative:

$$0000_2 = 0$$

$$0010_2 = 2$$

$$1010_2 = -2$$

$$1111_2 = -7$$

$$1000_2 = -0?$$

Can be made to work, but addition is
annoying:

If the signs match, add the magnitudes
and use the same sign.

If the signs differ, subtract the smaller
number from the larger; return the
sign of the larger.

One's Complement Numbers

Like Signed Magnitude, a leading 1 indicates a negative One's Complement number.

To negate a number, complement (flip) each bit.

$$0000_2 = 0$$

$$0010_2 = 2$$

$$1101_2 = -2$$

$$1000_2 = -7$$

$$1111_2 = -0?$$

Addition is nicer: just add the one's complement numbers as if they were normal binary.

Really annoying having a -0 : two numbers are equal if their bits are the same or if one is 0 and the other is -0 .



**NOT ALL
ZEROS
ARE CREATED
EQUAL**

ZERO CALORIES. MAXIMUM PEPSI™ TASTE.



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Two's Complement Numbers



Really neat trick: make the most significant bit represent a *negative* number instead of positive:

$$1101_2 = -8 + 4 + 1 = -3$$

$$1111_2 = -8 + 4 + 2 + 1 = -1$$

$$0111_2 = 4 + 2 + 1 = 7$$

$$1000_2 = -8$$

Easy addition: just add in binary and discard any carry.

Negation: complement each bit (as in one's complement) then add 1.

Very good property: no -0

Two's complement numbers are equal if all their bits are the same.

Number Representations Compared

Bits	Binary	Signed Mag.	One's Comp.	Two's Comp.
0000	0	0	0	0
0001	1	1	1	1
⋮				
0111	7	7	7	7
1000	8	-0	-7	-8
1001	9	-1	-6	-7
⋮				
1110	14	-6	-1	-2
1111	15	-7	-0	-1

Smallest number

Largest number

Fixed-point Numbers



How to represent fractional numbers? In decimal, we continue with negative powers of 10:

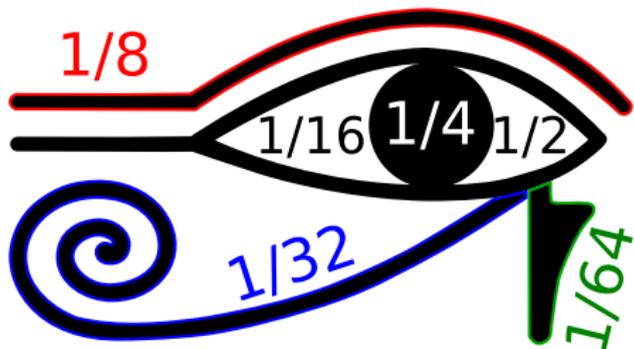
$$31.4159 = 3 \times 10^1 + 1 \times 10^0 + 4 \times 10^{-1} + 1 \times 10^{-2} + 5 \times 10^{-3} + 9 \times 10^{-4}$$

The same trick works in binary:

$$\begin{aligned} 1011.0110_2 &= 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 + \\ &\quad 0 \times 2^{-1} + 1 \times 2^{-2} + 1 \times 2^{-3} + 0 \times 2^{-4} \\ &= 8 + 2 + 1 + 0.25 + 0.125 \\ &= 11.375 \end{aligned}$$

F a
u c
Interesting

The ancient Egyptians used binary fractions:



The Eye of Horus

Binary-Coded Decimal



thinkgeek.com

Humans prefer reading decimal numbers; computers prefer binary.

BCD is a compromise: every four bits represents a decimal digit.

Dec	BCD
0	0000 0000
1	0000 0001
2	0000 0010
⋮	⋮
8	0000 1000
9	0000 1001
10	0001 0000
11	0001 0001
⋮	⋮
18	0001 1000
19	0001 1001
20	0010 0000
⋮	⋮

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 158 \\ +242 \\ \hline \end{array}$$

$$\begin{array}{r} 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \end{array} \text{ First group}$$



BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 158 \\ +242 \\ \hline \end{array}$$

$$\begin{array}{r} 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +0110 \\ \hline \end{array} \begin{array}{l} \text{First group} \\ \text{Correction} \end{array}$$

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 1 \\ 158 \\ +242 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 1 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +\ 0110 \\ \hline 1010\ 0000 \end{array} \begin{array}{l} \text{First group} \\ \text{Correction} \\ \text{Second group} \end{array}$$

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 1 \\ 158 \\ +242 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 1 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +0110 \\ \hline 1010\ 0000 \\ +0110 \\ \hline \hline \hline \end{array}$$

First group
Correction
Second group
Correction

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 11 \\ 158 \\ +242 \\ \hline 00 \end{array}$$

$$\begin{array}{r} 11 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +0110 \\ \hline 1010\ 0000 \\ +0110 \\ \hline 0100\ 0000 \\ \hline \hline \end{array}$$

First group
Correction
Second group
Correction
Third group

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 11 \\ 158 \\ +242 \\ \hline 400 \end{array}$$

$$\begin{array}{r} 11 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ + 0110 \\ \hline 1010\ 0000 \\ + 0110 \\ \hline 0100\ 0000 \\ \hline 0100\ 0000\ 0000 \end{array}$$

First group
Correction
Second group
Correction
Third group
(No correction)
Result